

# COUNTY PRESERVATION

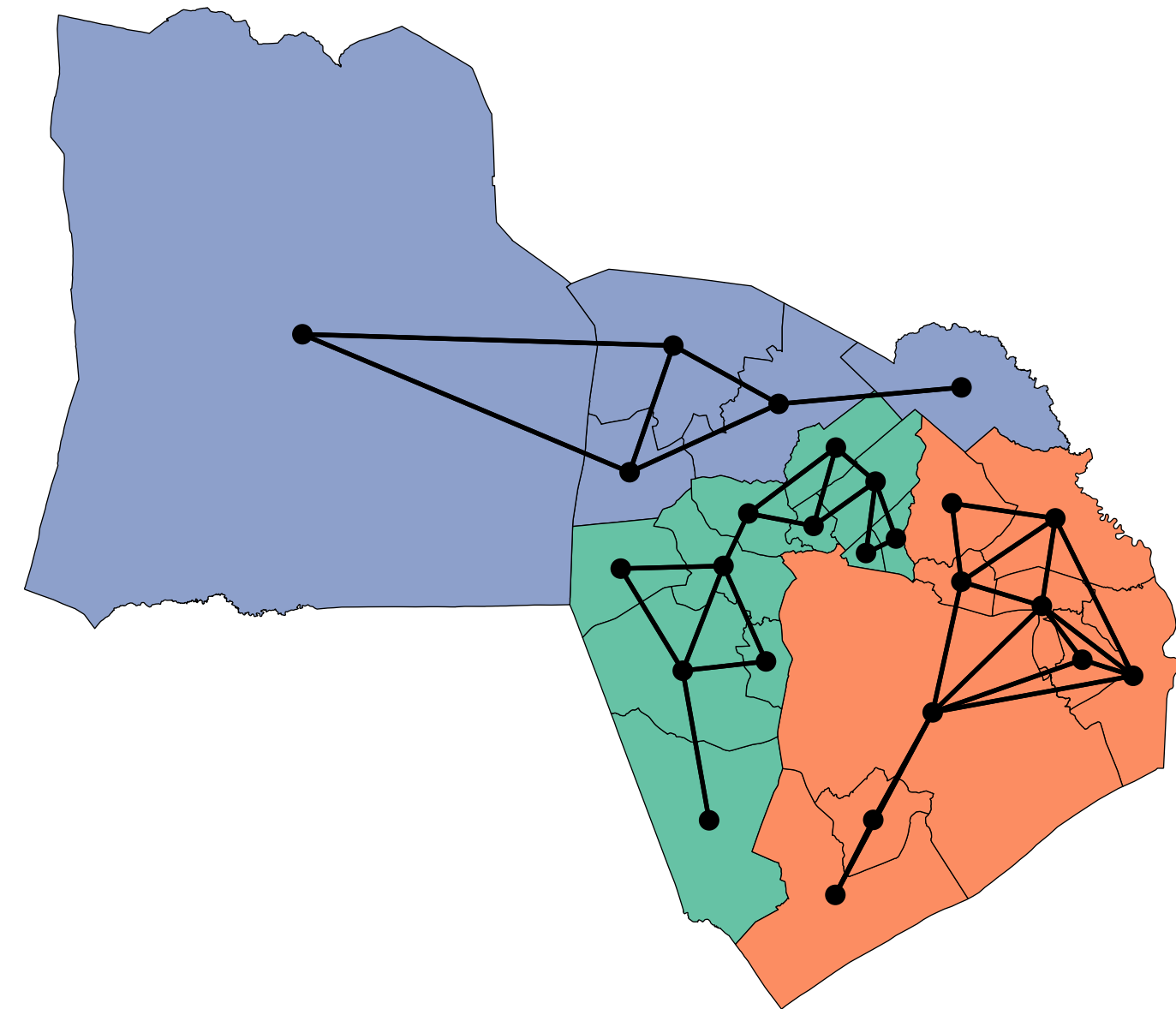
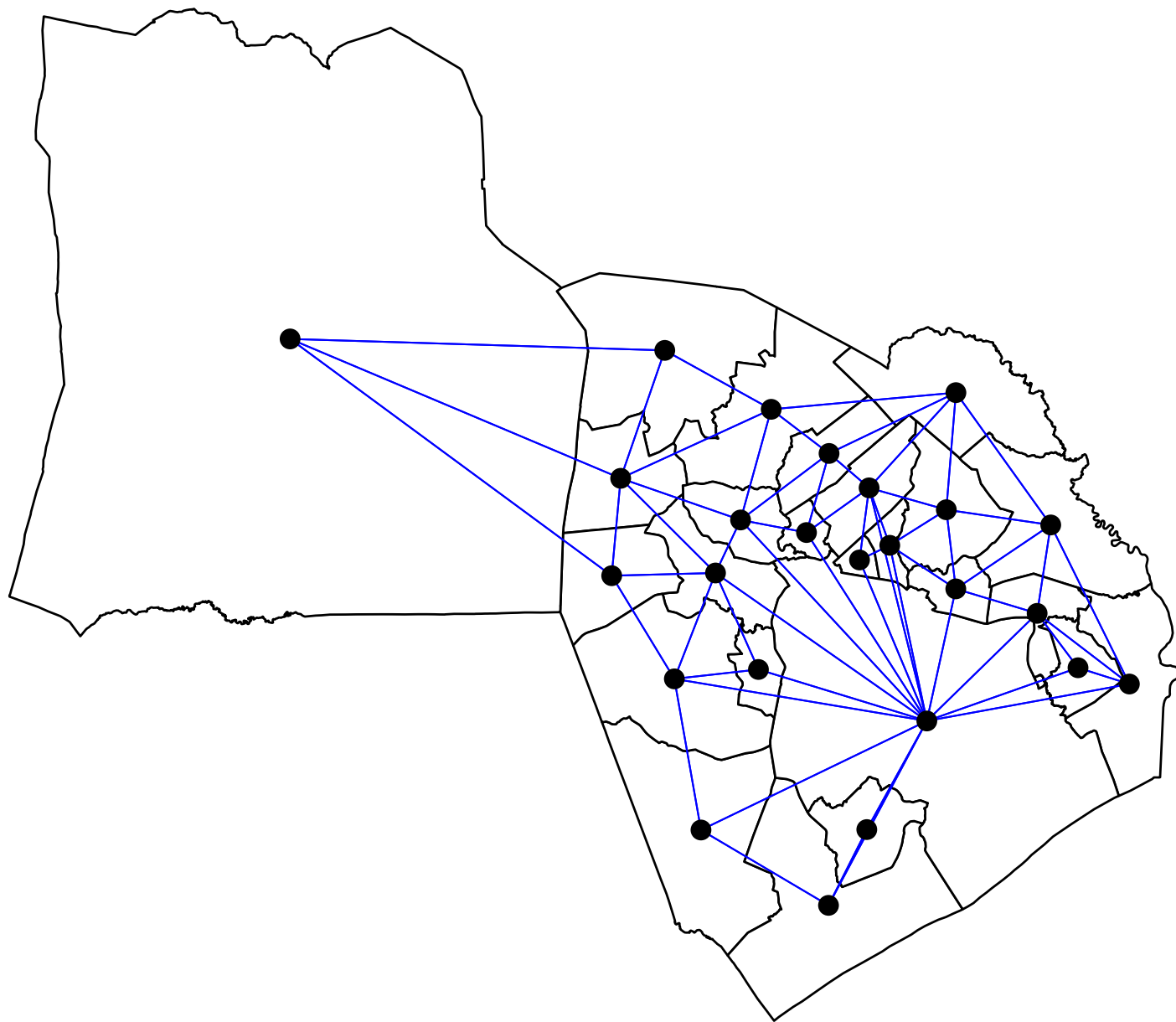
SAMPLING, CLUSTERING, CONSTRAINING

QUANTIFYING GERRYMANDERING  
MARCH 3RD, 2020

GREGORY HERSCHLAG, JONATHAN MATTINGLY,  
DANIEL CARTER, ZACH HUNTER, ERIC AUTRY  
+THE TEAM AT DUKE (AND BEYOND)

**If  $n$  districts:**

$\xi: (\text{precincts}) \longrightarrow \{1, 2, \dots, n\} \iff \xi = (\xi_1, \dots, \xi_n)$  **partition into sub-graphs**

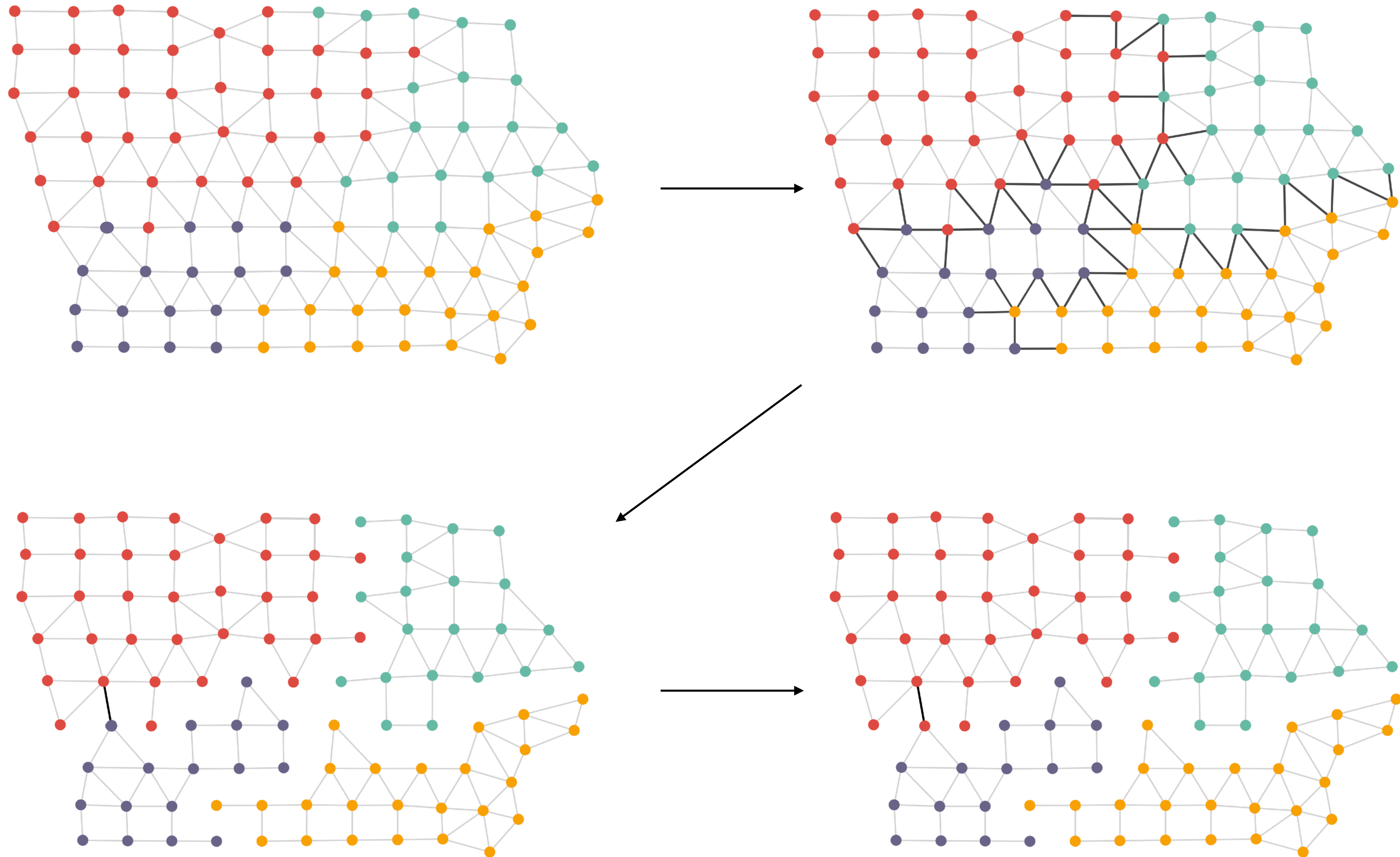


**(Probability of plan)  $\propto e^{-\beta(\text{Plan Score})}$**

$$\mu(\xi) = \frac{1}{Z} e^{-\beta J(\xi)}$$



# Single Node Flip Markov Chains



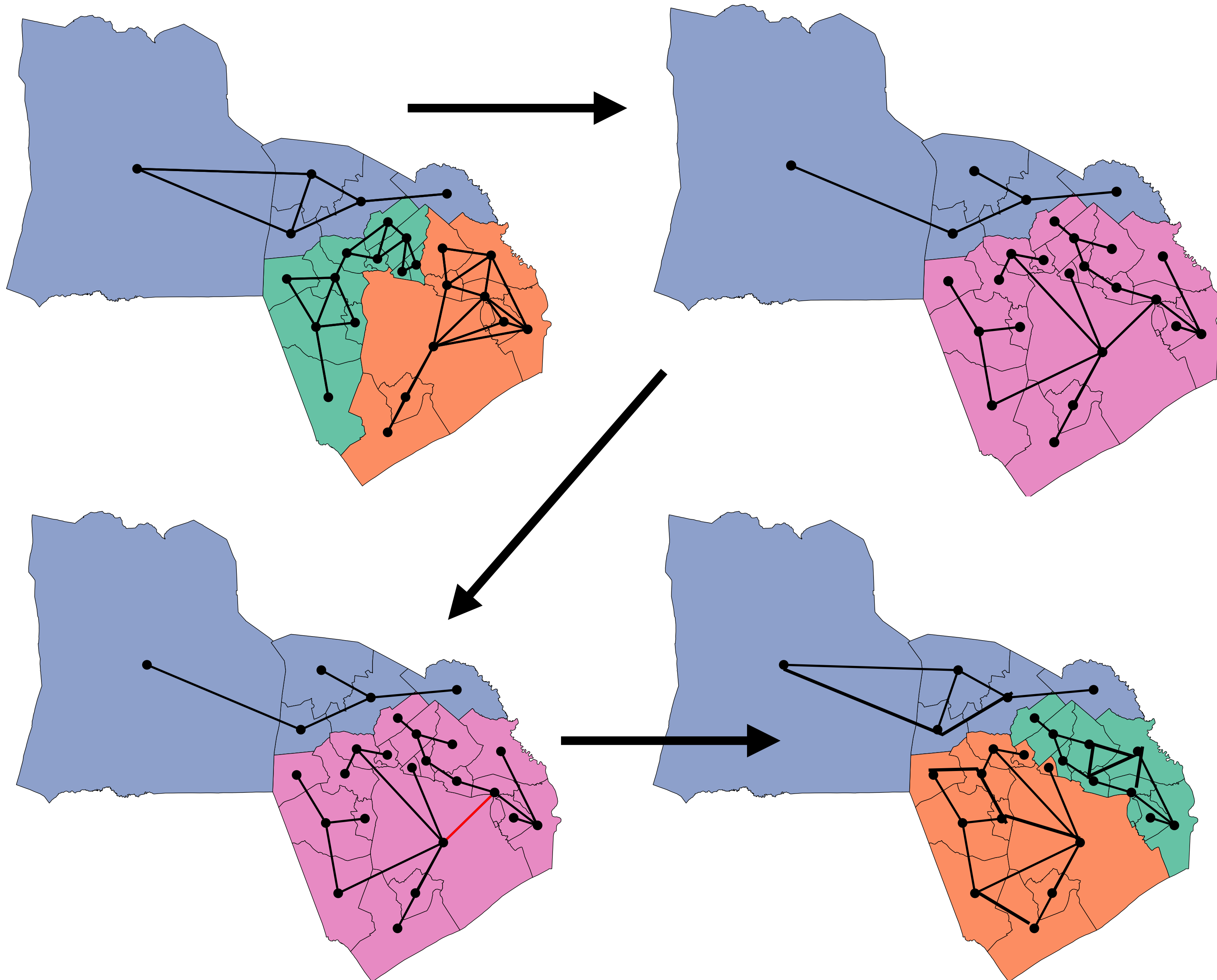
Then accept/reject according to a score function

# One Step of ReCom Markov Chain

DeFord, Duchin,  
Solomon

## ReCom Algorithm

1. Pick adjacent pair of districts to merge
2. Draw Spanning tree on merged graph (Willson's Alg)
3. Find permissible cuts (e.g. within Pop constraint)
4. cut in two, return new subgraphs



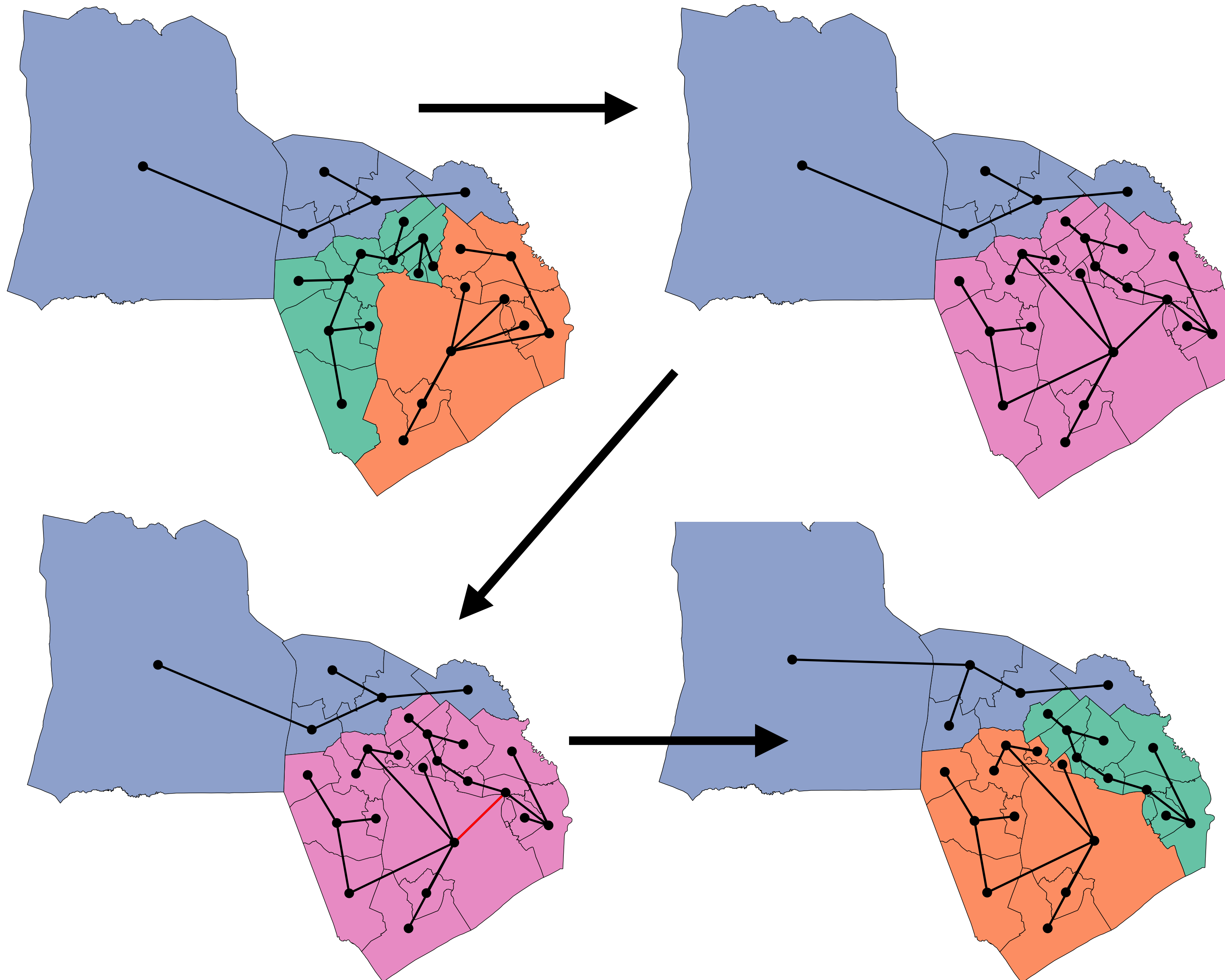


# Or expanded state space in Merge Split

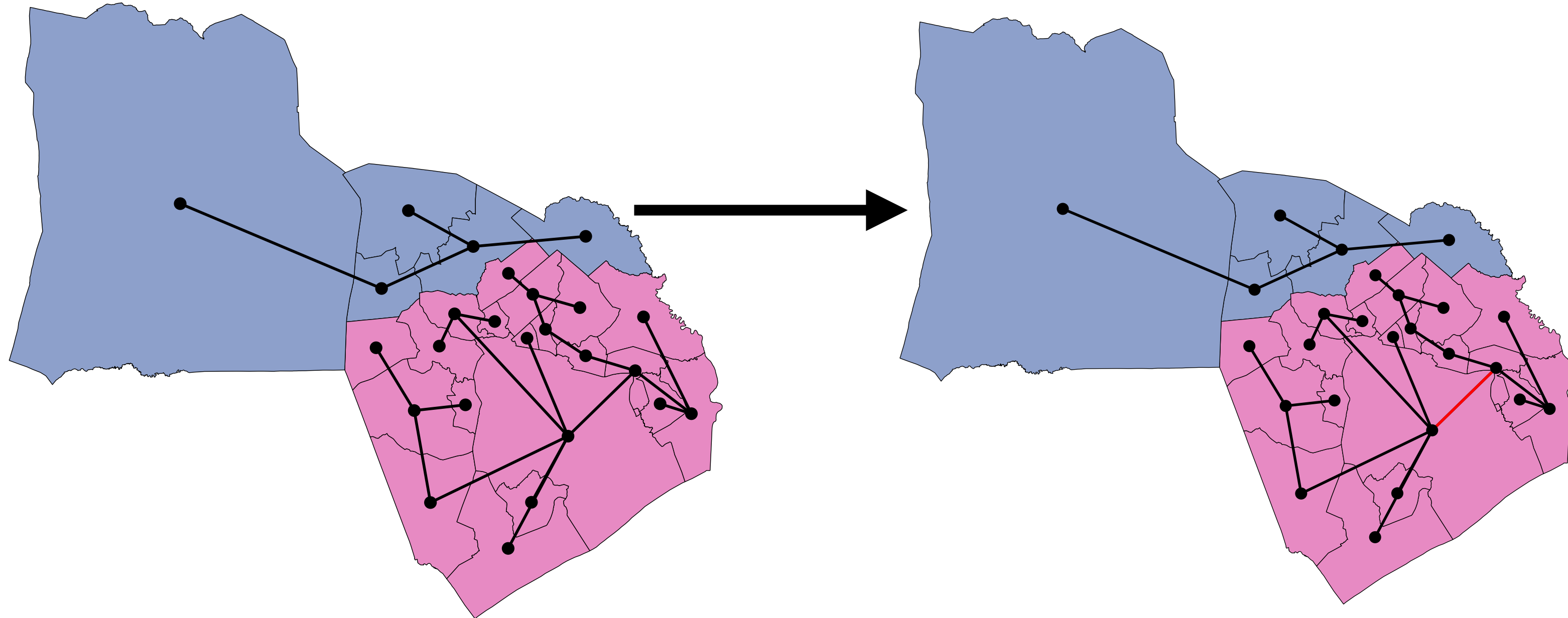
**Carter, GH, Hunter,  
Mattingly**

## Merge-Split Algorithm

1. Pick adjacent pair of districts to merge
2. Draw Spanning tree on merged graph (Willson's Alg)
3. Find permissible cuts (e.g. within Pop constraint)
4. Cut in two, return new subgraphs



# Aside on Recom and improvements



$$P_{tree}(T_A, T_B, e)P_{cut}(e) = \frac{1}{\tau(A \cup B)} \frac{1}{|E_{cut}|}$$

**Needed for Recom  
(computationally intractable)**

$$\sum_{T_A \in ST(A)} \sum_{T_B \in ST(B)} \sum_{e \in A \cup B} P_{tree}(T_A, T_B, e)P_{cut}(e) = \sum_{T_A \in ST(A)} \sum_{T_B \in ST(B)} \sum_{e \in A \cup B} \frac{1}{\tau(A \cup B)} \frac{1}{|E_{cut}|(T_A, T_B, e)}$$

**Merge-Split**

$$\sum_{e \in A \cup B} P_{tree}(T_A, T_B, e)P_{cut}(e) = \sum_{e \in A \cup B} \frac{1}{\tau(A \cup B)} \frac{1}{|E_{cut}|(T_A, T_B, e)}$$

**Reversible Recom**

$$\sum_{T_A \in ST(A)} \sum_{T_B \in ST(B)} \sum_{e \in A \cup B} P_{tree}(T_A, T_B, e)P_{cut}(e) = \sum_{T_A \in ST(A)} \sum_{T_B \in ST(B)} \sum_{e \in A \cup B} \frac{1}{\tau(A \cup B)} \frac{1}{m}$$



# County preservation in North Carolina

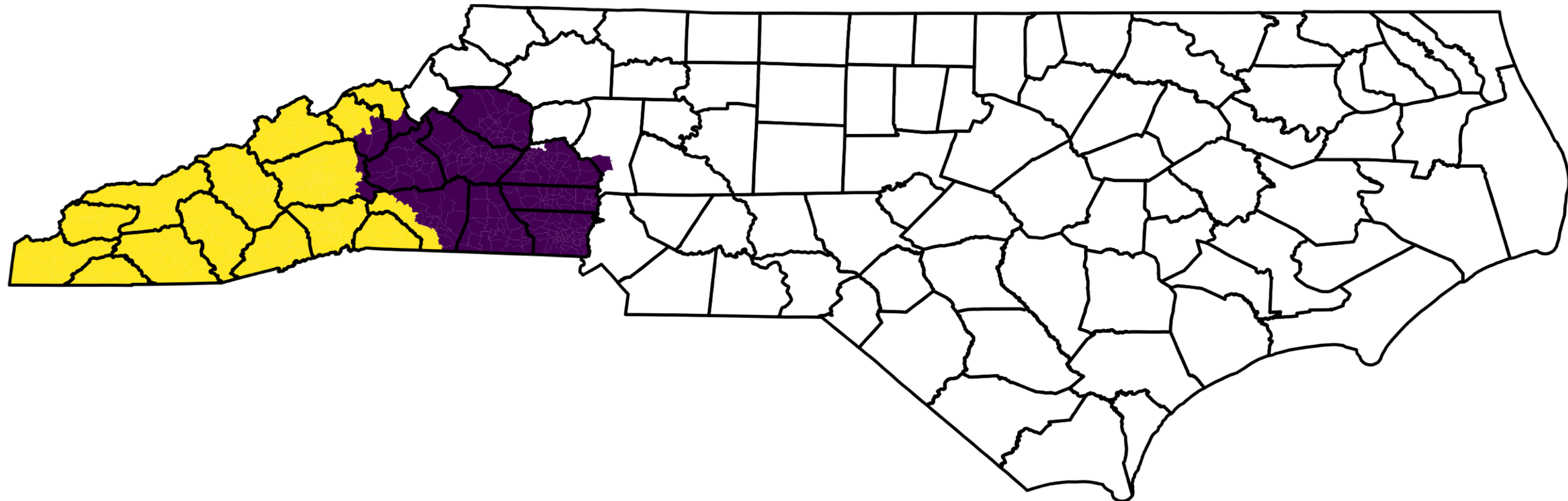
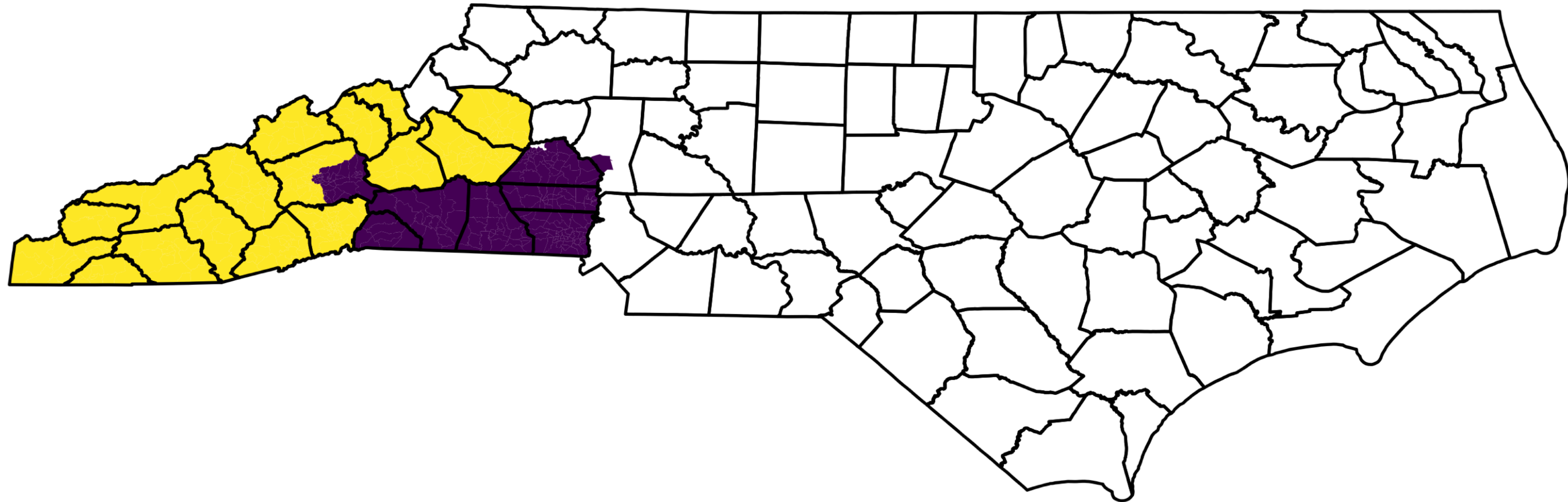
- ***Congressional***

- *2016 Criteria: “Division of counties shall only be made for reasons of equalizing population, consideration of incumbency and political impact. Reasonable efforts shall be made not to divide a county into more than two districts.”*

- ***Legislative***

- *County clusters*
- *Minimization of traversals*
- *When counties are split, keep them together as much as possible*

# Difficulty of county preservation





# What are the minimal number of county splits?

**Theorem:**  $d$  districts must introduce at least  $d-1$  county splits (when we can't evenly partition counties)

**County splits:** the county splits of a county are the number of districts a county intersects minus 1

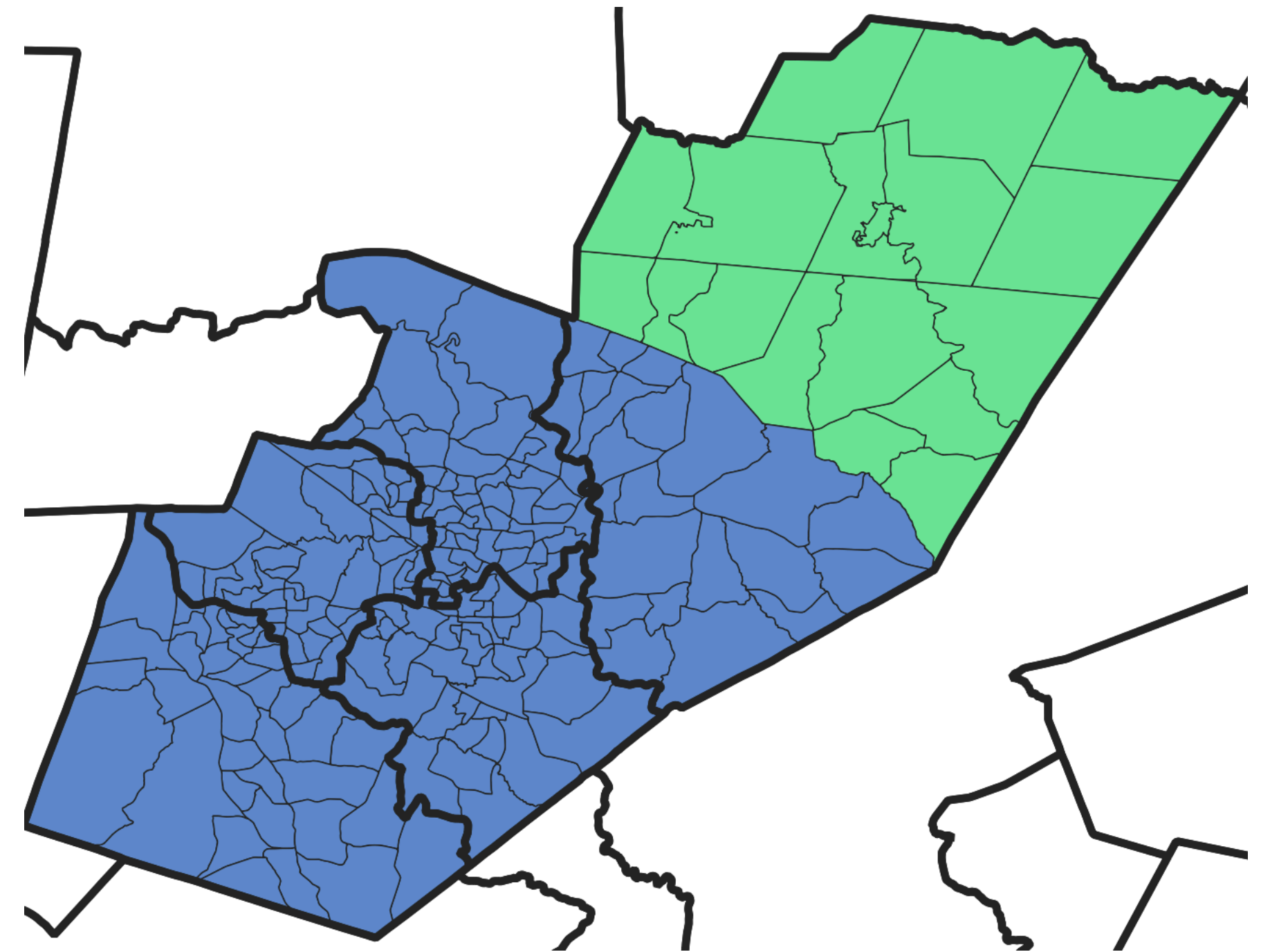
**Conjecture:** In nearly all redistricting problems, the bound from the theorem will be tight

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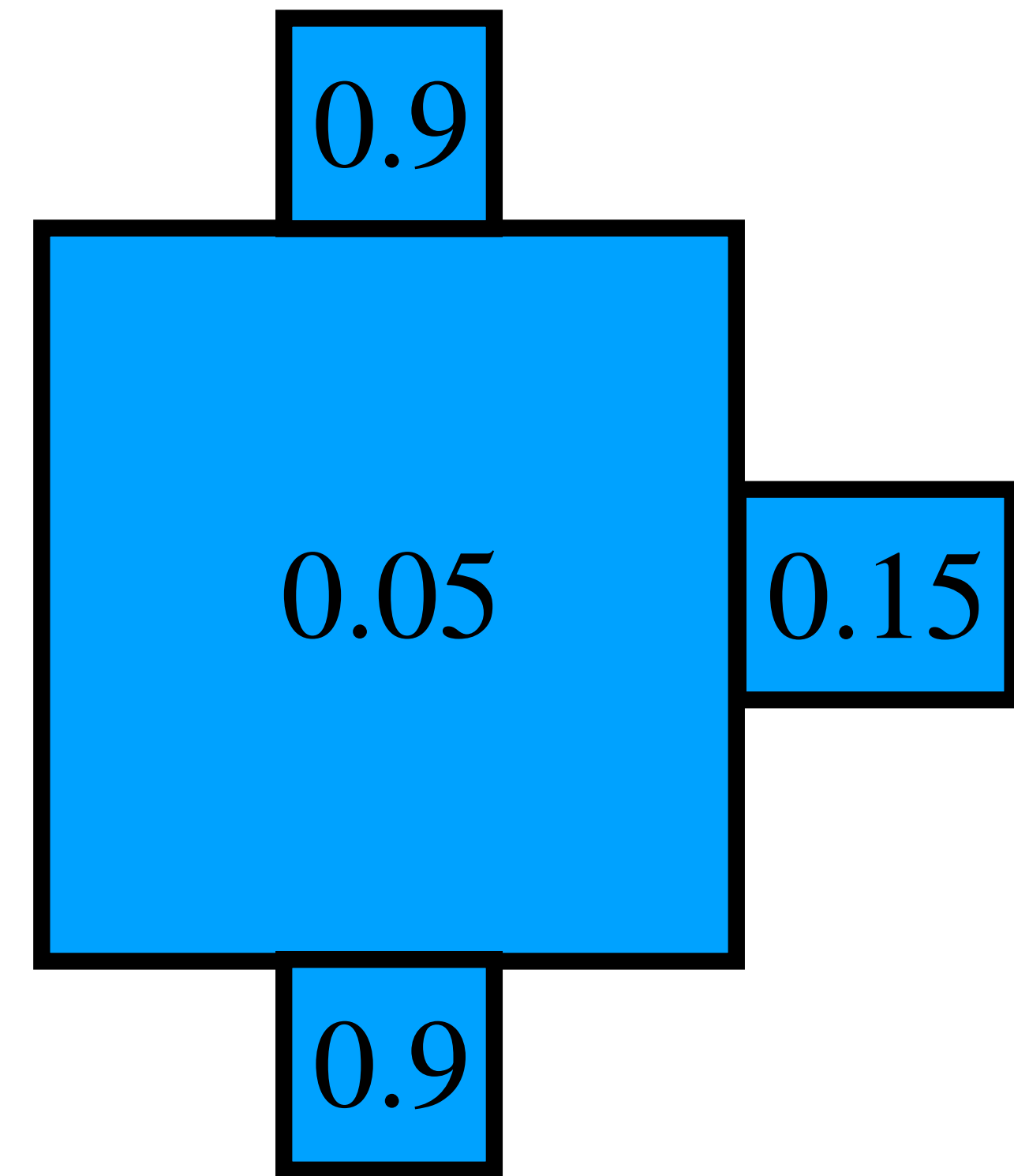


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Fraction of a district population

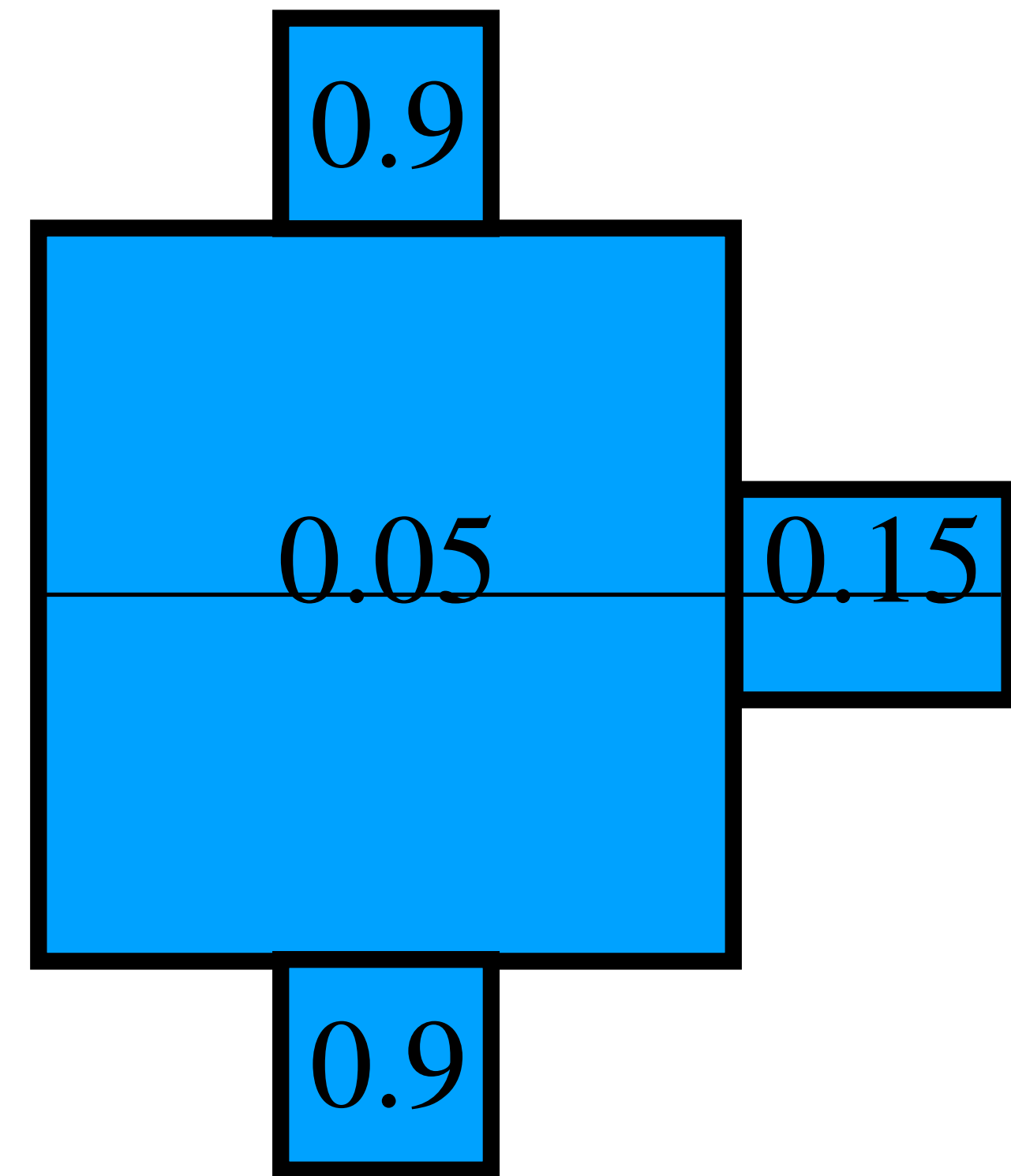
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Carter, GH, Hunter, Mattingly  
<https://arxiv.org/abs/1908.11801>; under review



**Two districts must split the central county and one of the satellites**

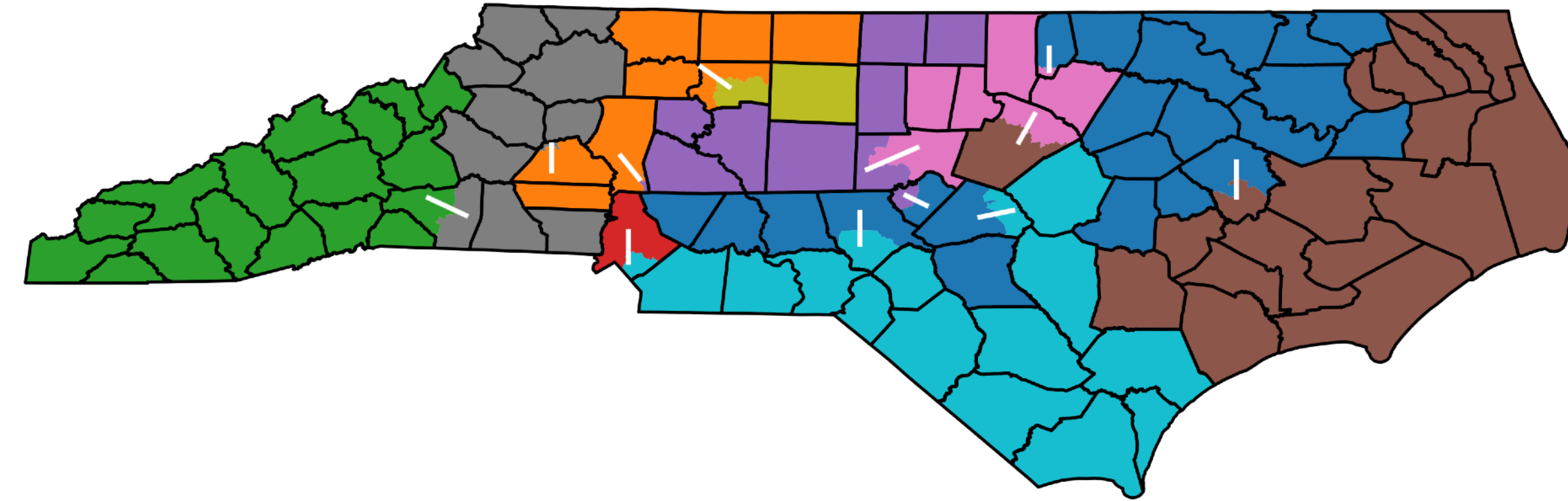


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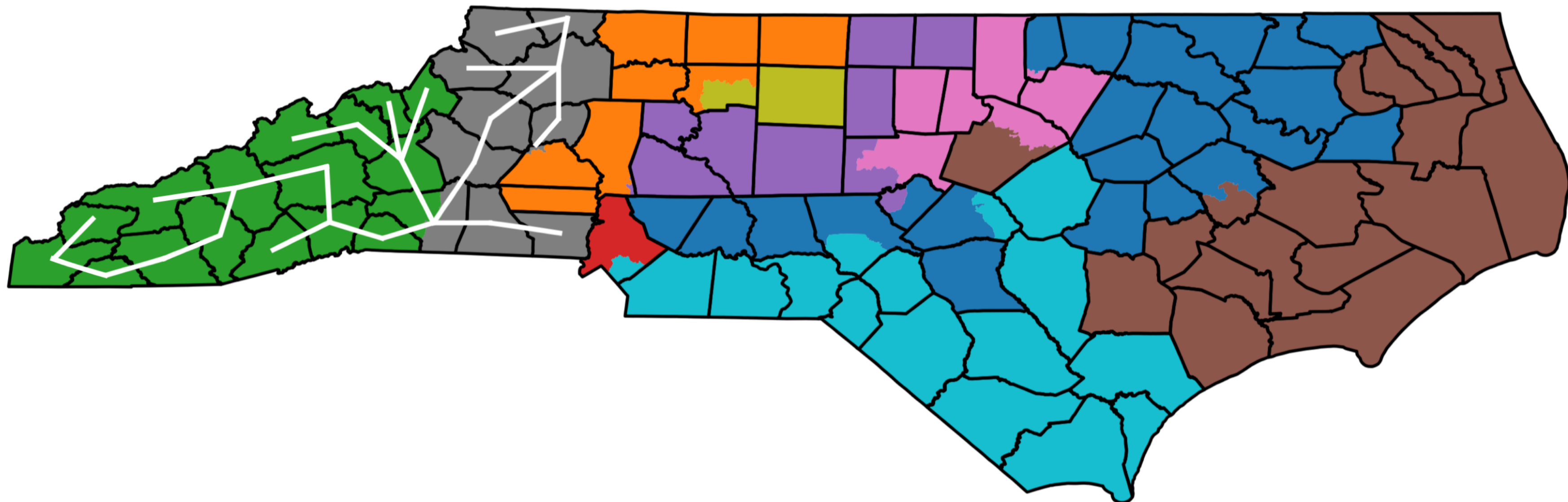
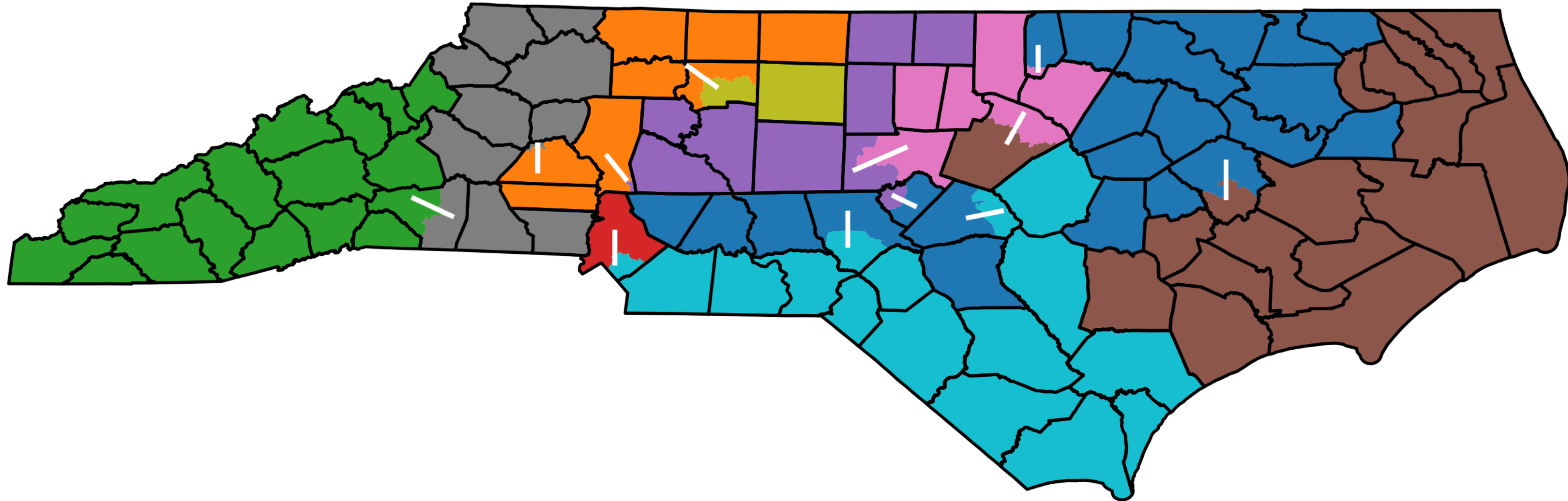
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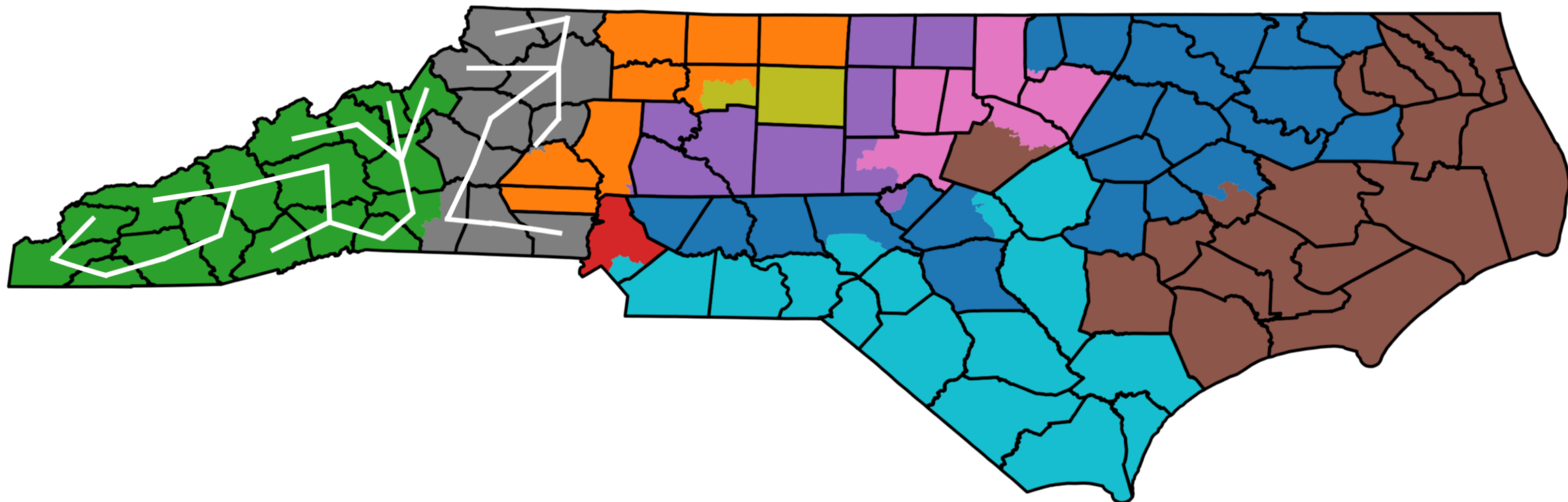
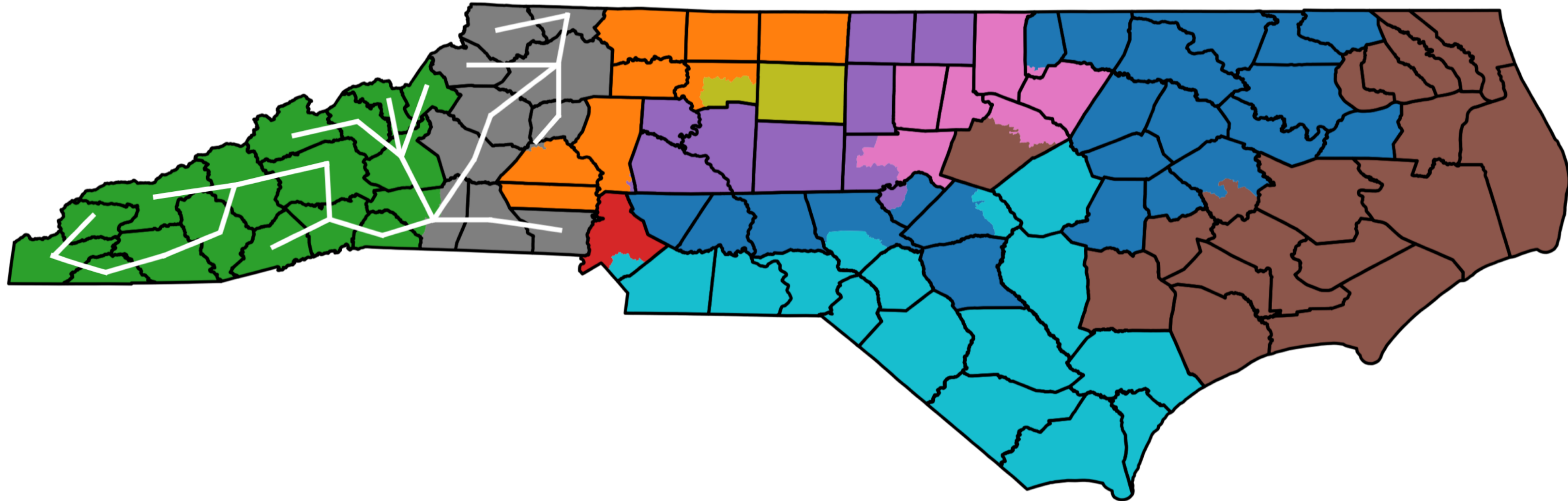
Enacted Plan (2019)

# Draw a merged tree on counties now



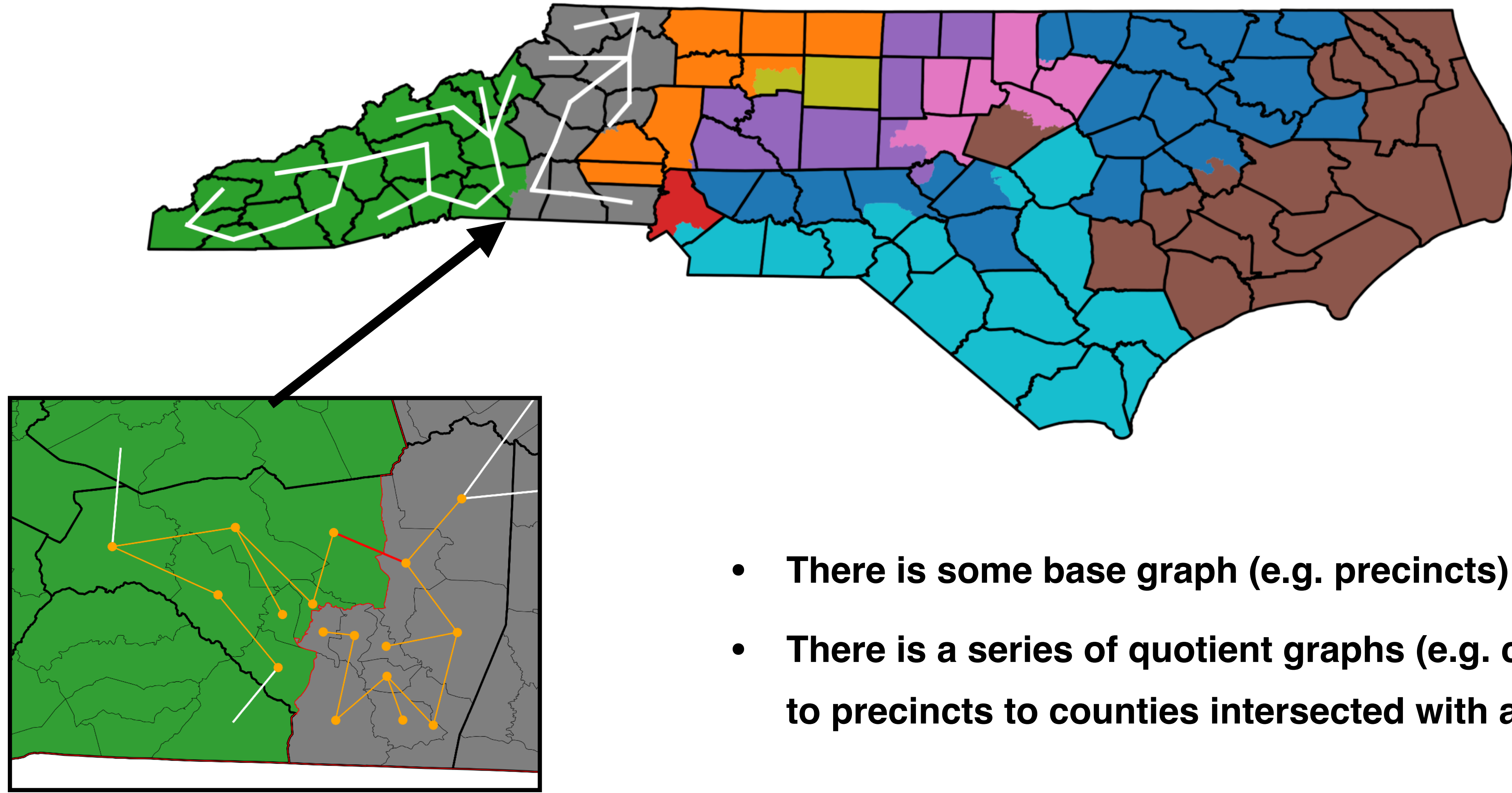


# Splitting on Nodes rather than Edges



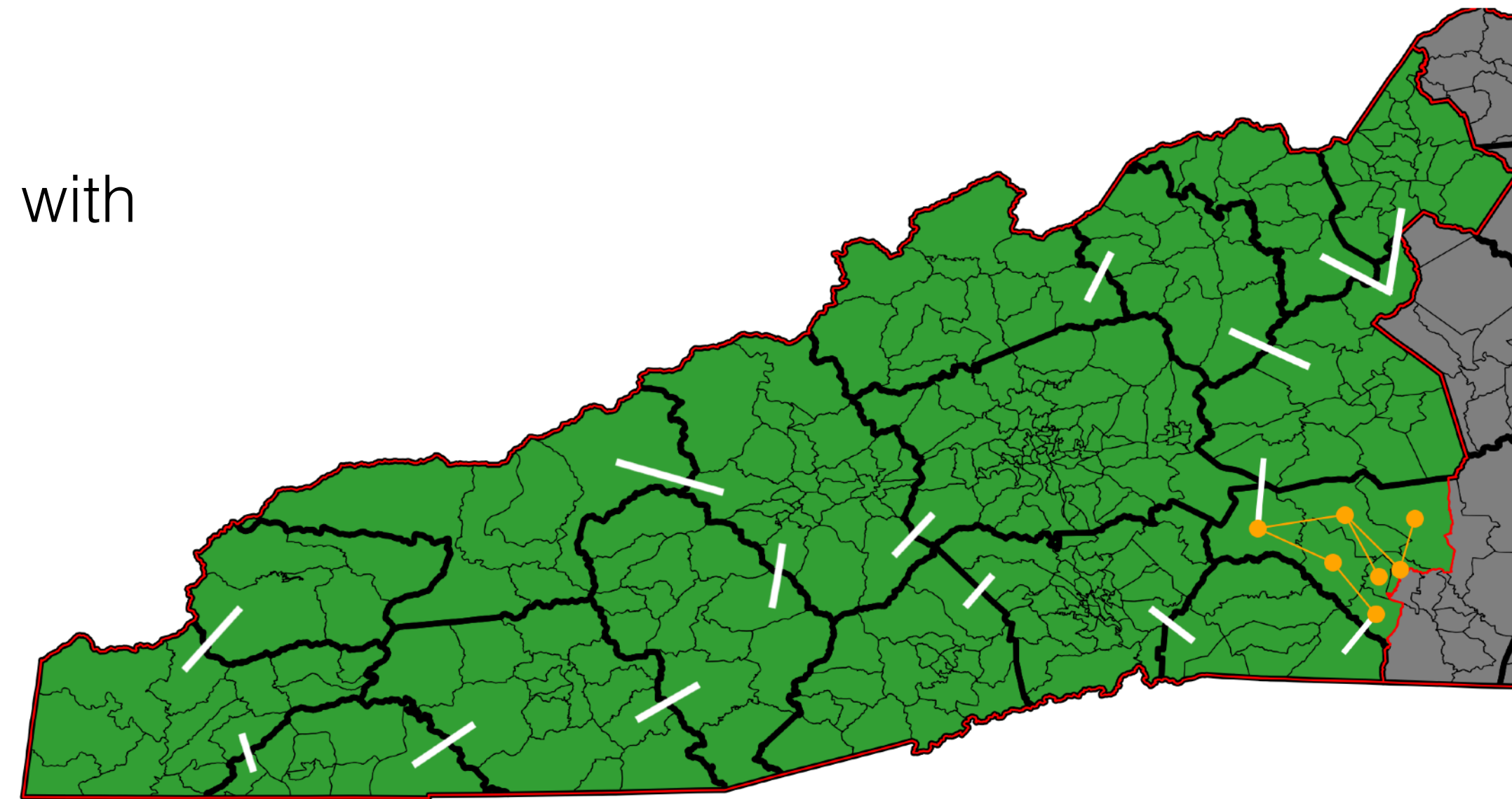
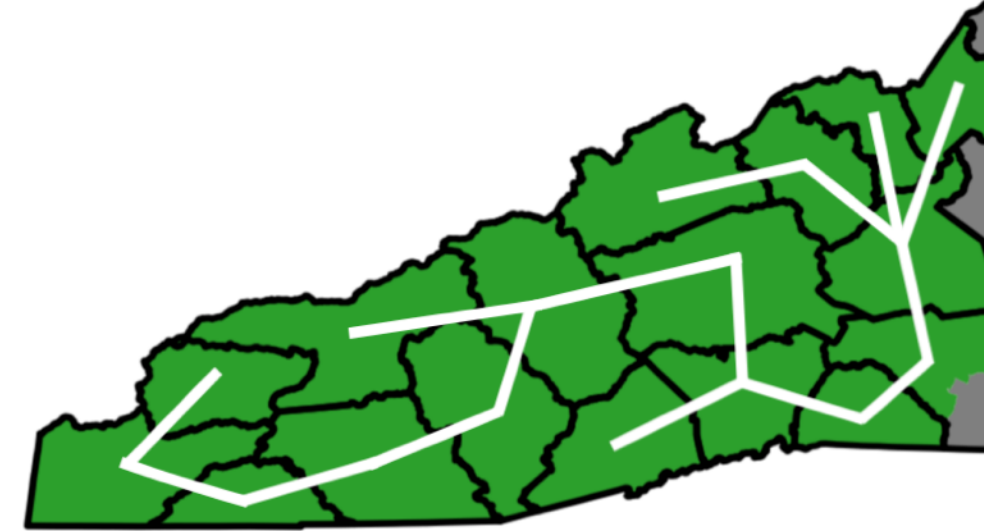


# Nodes as Nested Trees

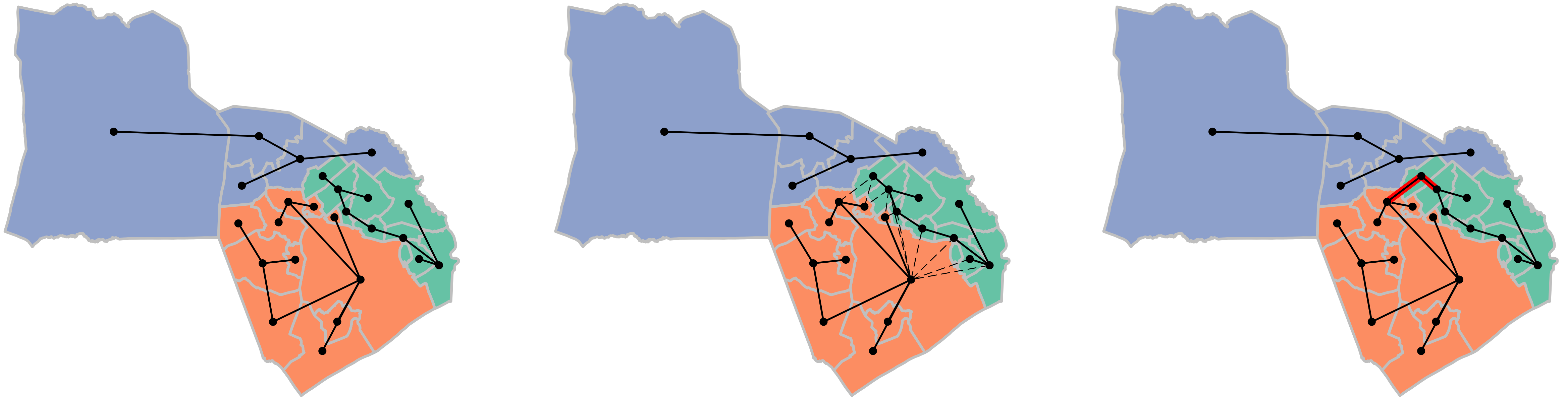


# The State Space

- The state space is a forest at the finest level
- Each district is a tree on the finest scale such that
  - There is a simple tree graph at the county level when intersected with the district
- Every nested node within the county/district graph has a tree with specified edges the top level tree



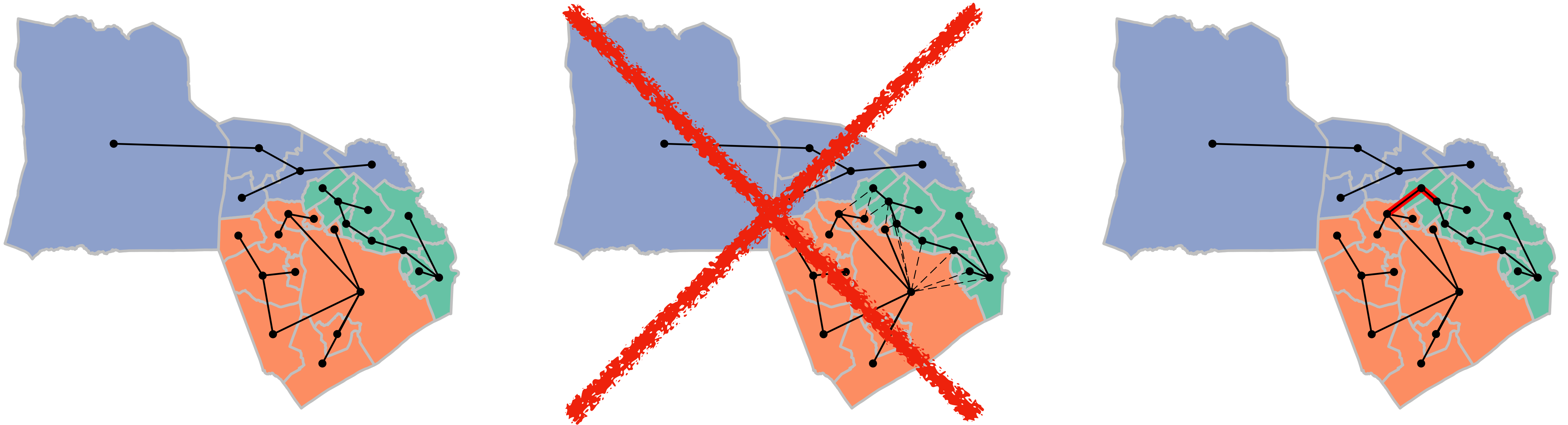
# The State Space



**Reverse Probability in Merge-Split requires reconnecting every edge that could have drawn the merged tree**



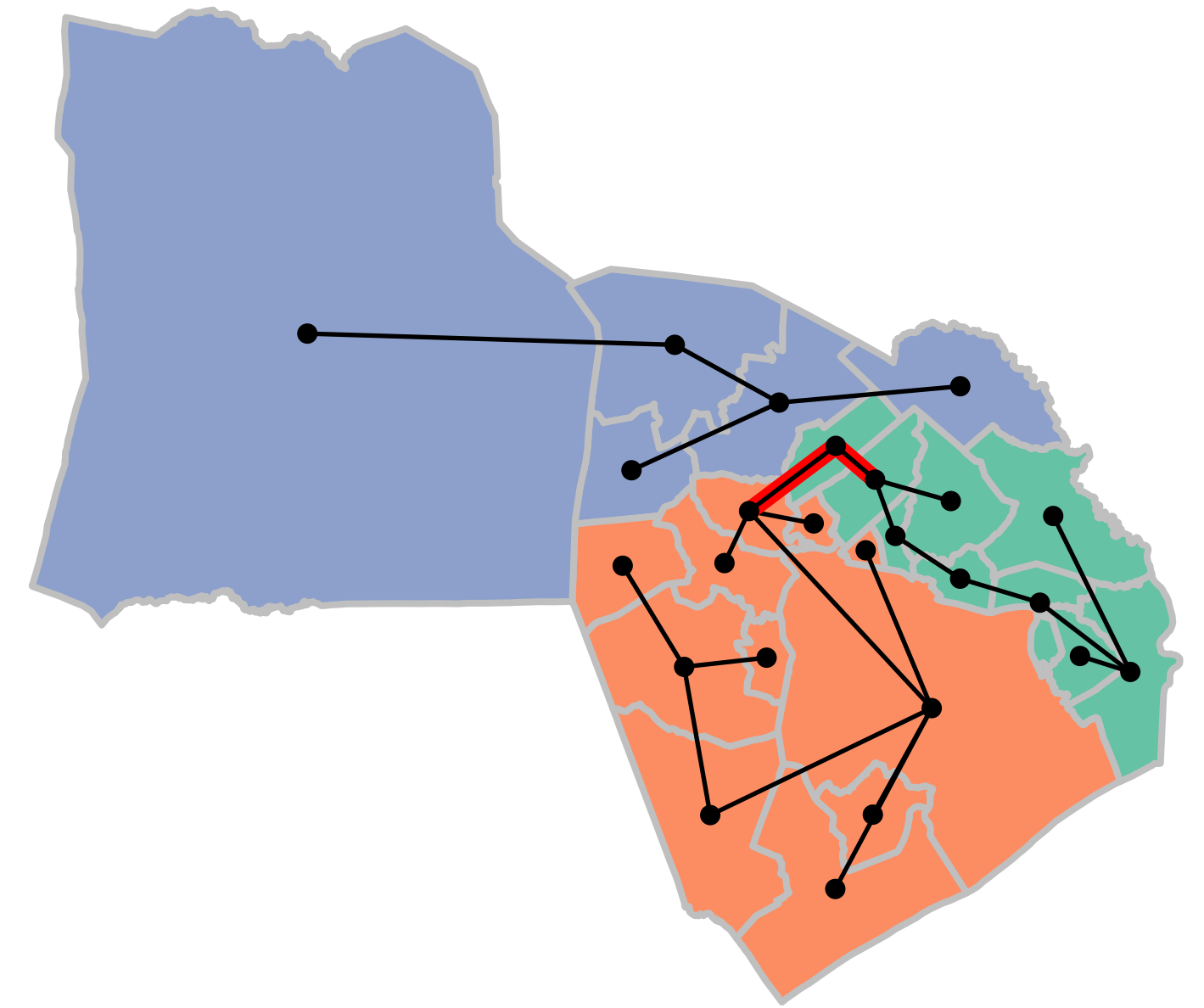
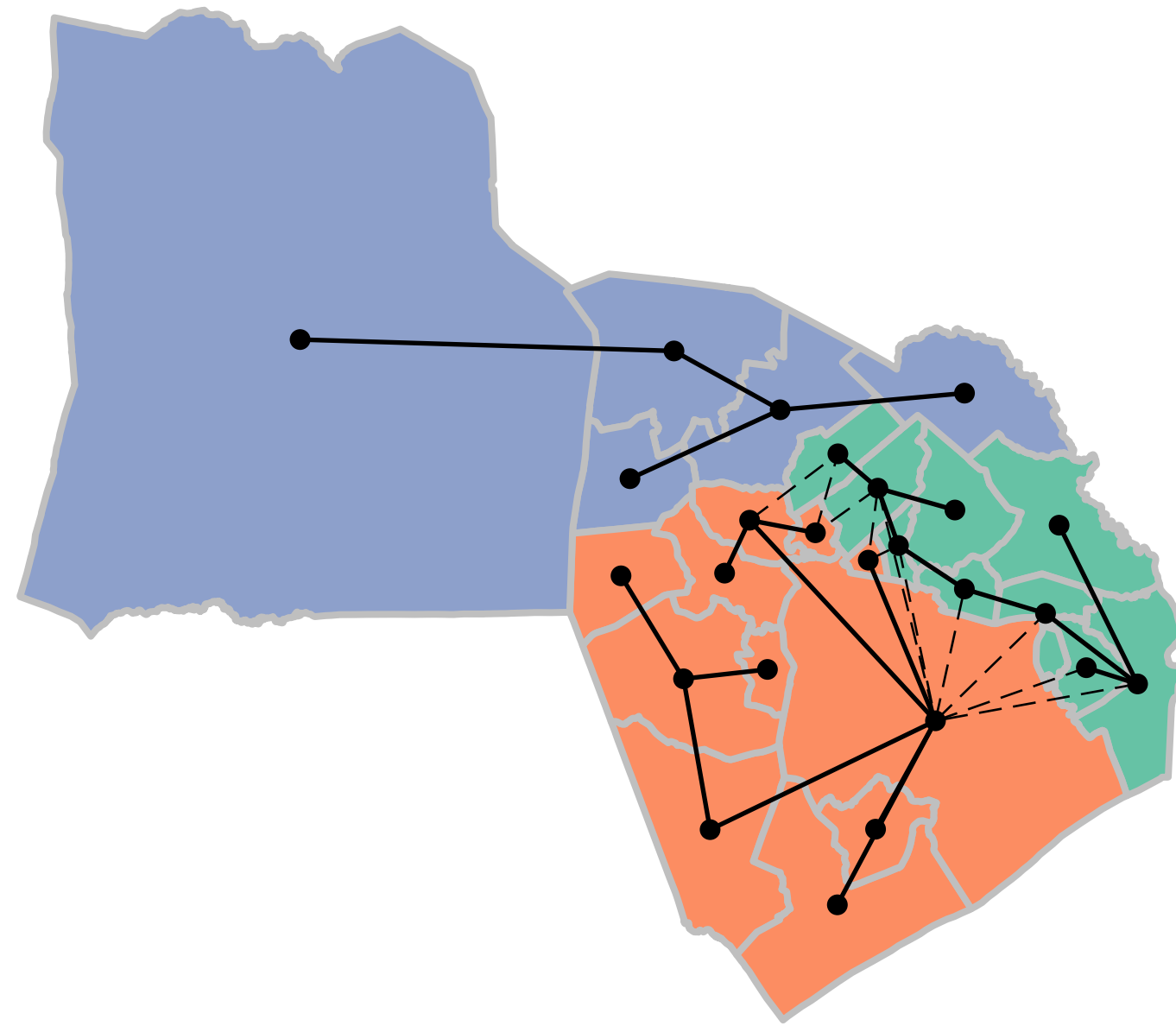
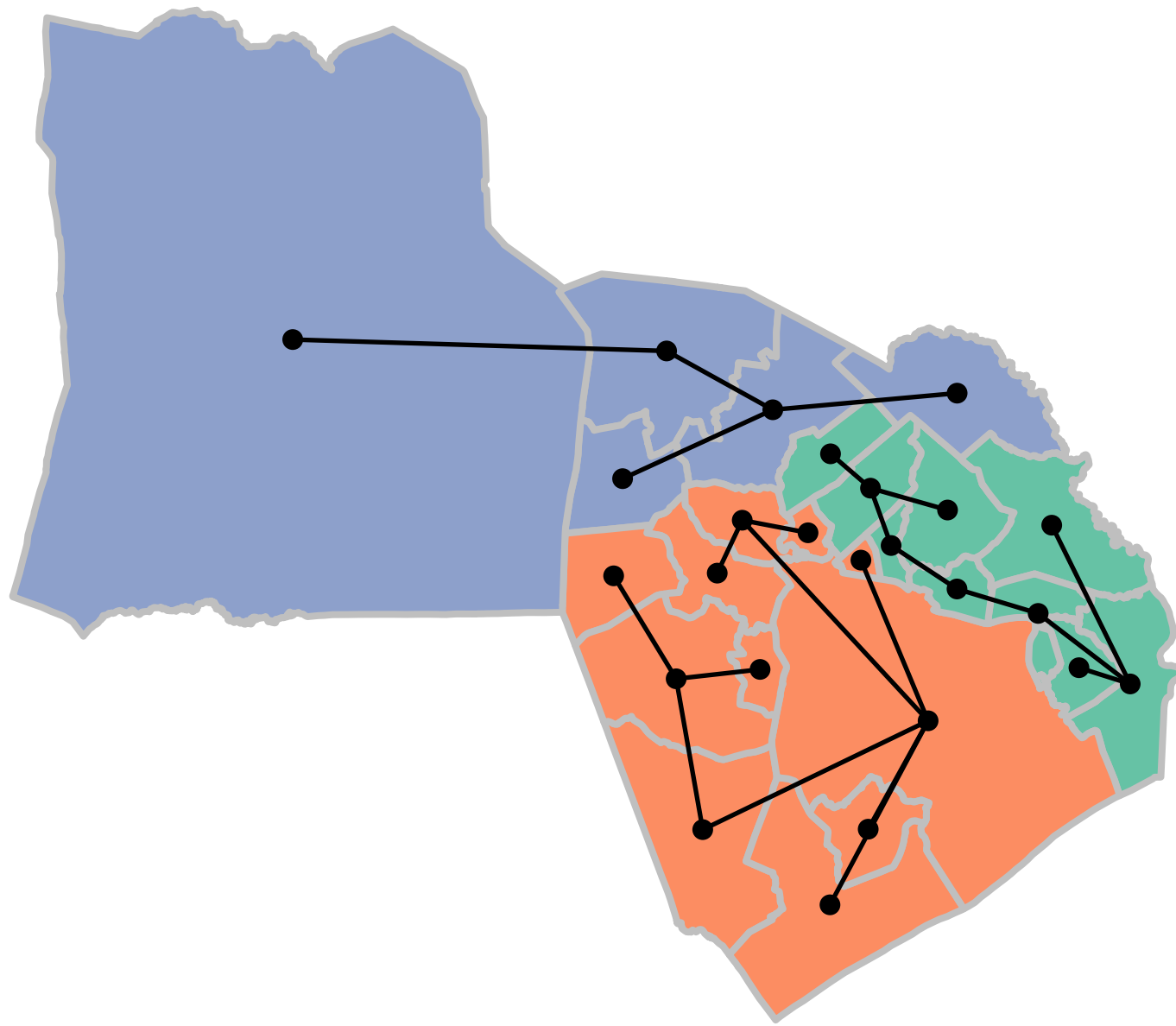
# Adding Persistent Connections



**Keep one of these edges in the state; only need to compute one cut set**



# Adding Persistent Connections



**Merge-Split**

**Reversible Recom**

**Persistent Edge Merge-Split**

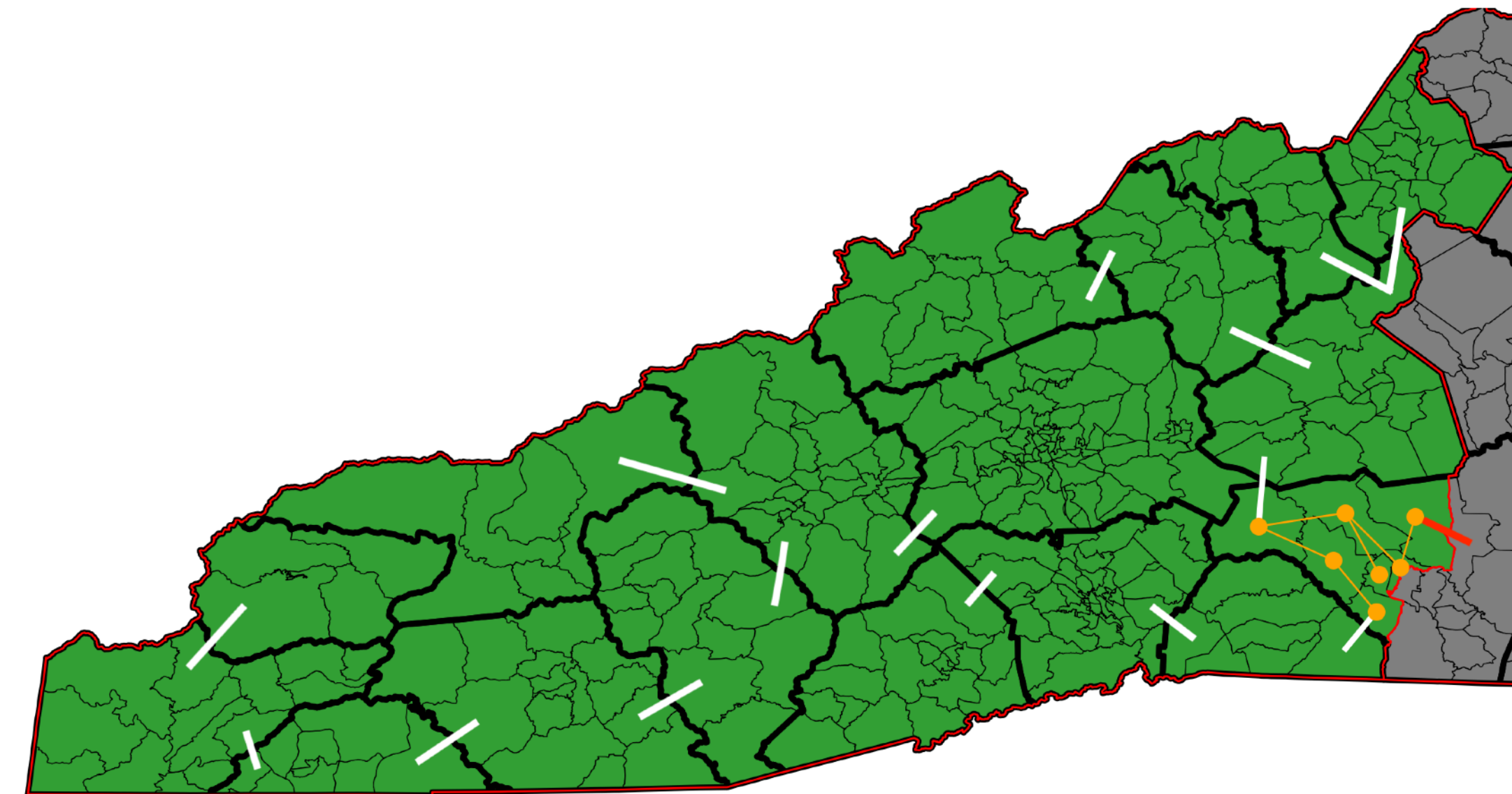
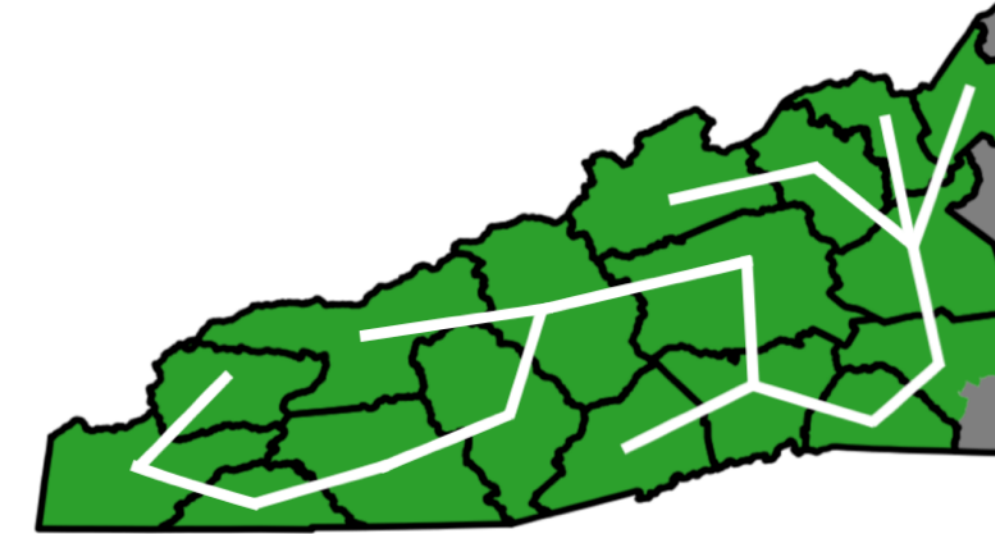
$$\sum_{e \in A \cup B} P_{tree}(T_A, T_B, e) P_{cut}(e) = \sum_{e \in A \cup B} \frac{1}{\tau(A \cup B)} \frac{1}{|E_{cut}|(T_A, T_B, e)}$$

$$\sum_{T_A \in ST(A)} \sum_{T_B \in ST(B)} \sum_{e \in A \cup B} P_{tree}(T_A, T_B, e) P_{cut}(e) = \sum_{T_A \in ST(A)} \sum_{T_B \in ST(B)} \sum_{e \in A \cup B} \frac{1}{\tau(A \cup B)} \frac{1}{m}$$

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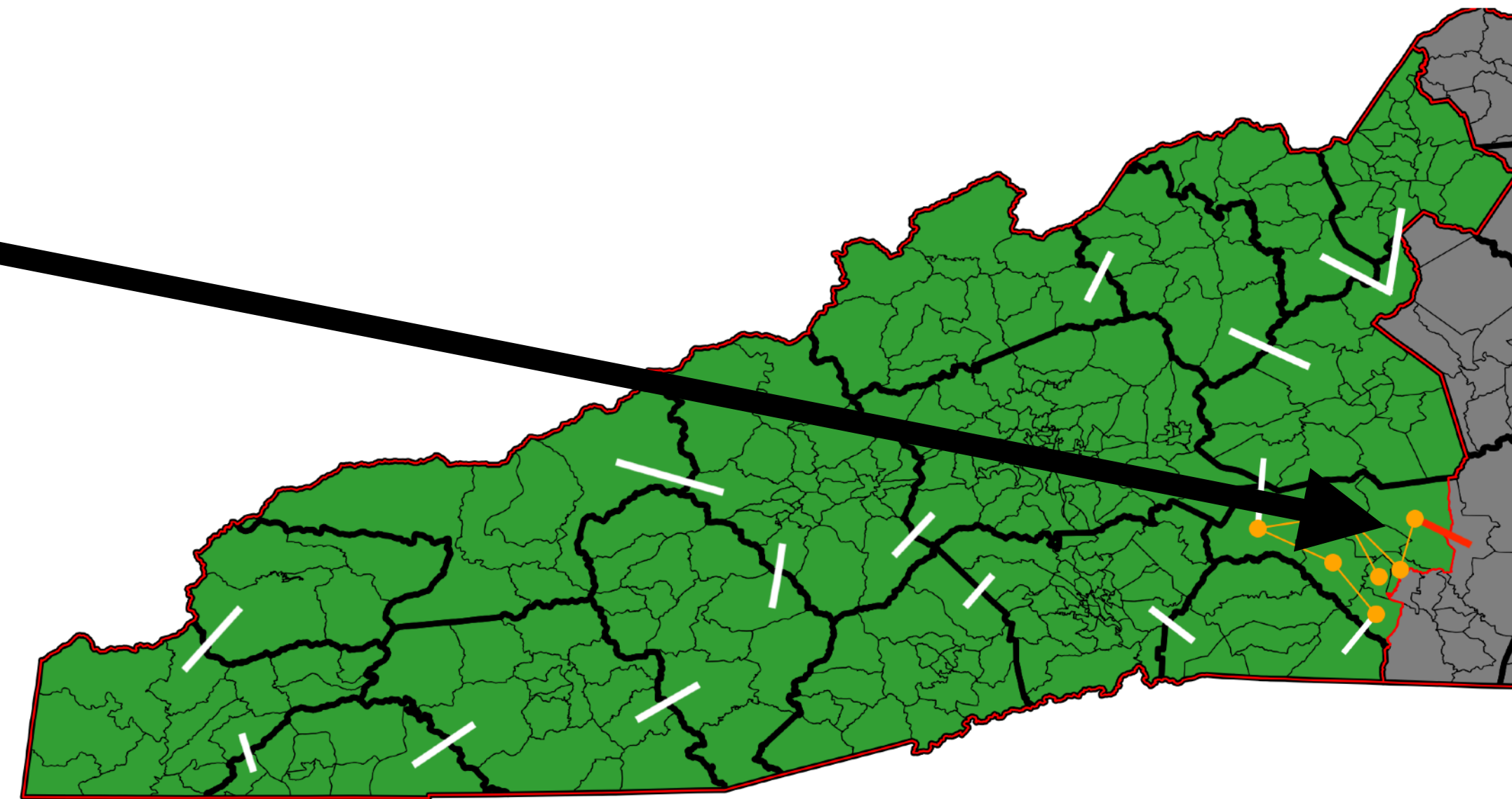
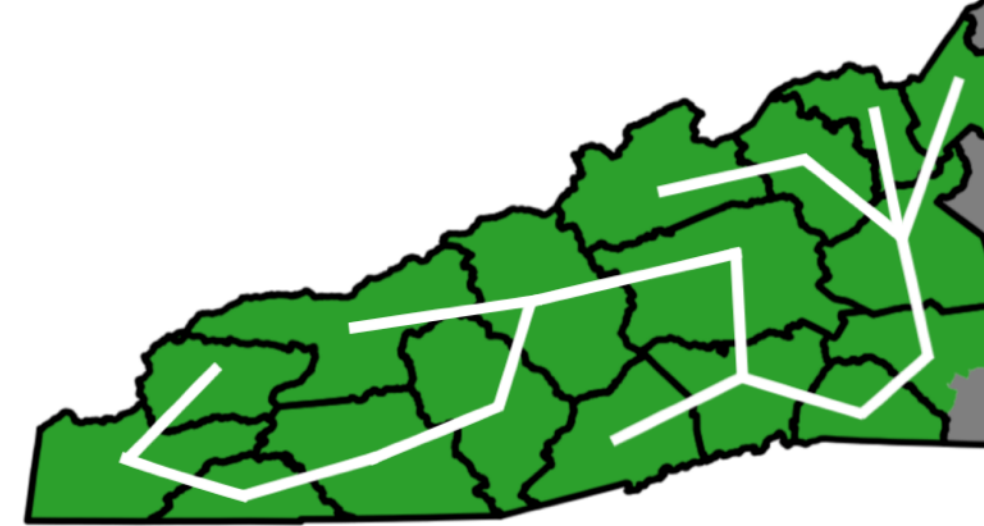
# The State Space

- The state space is a forest; each tree is a district
- Each district is a tree on the finest scale
- We keep edges between certain districts; call them persistent edges



# The State Space

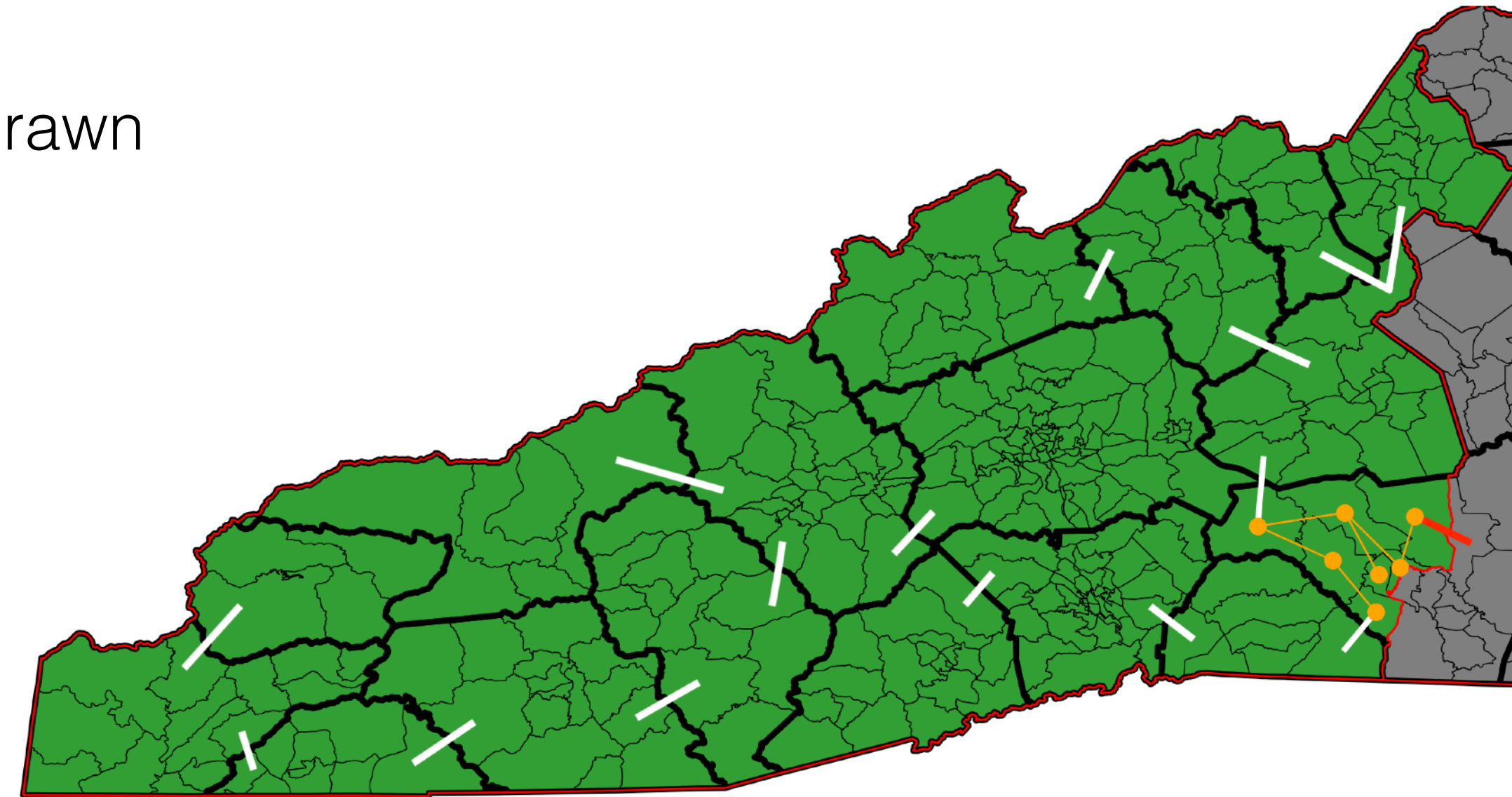
- The state space is a forest; each tree is a district
- Each district is a tree on the finest scale
- There are edges between certain districts





# The number of states associated with a plan

- For each district:
  - A product of the number of trees that can be drawn on each level of the hierarchy
- For each persistent edge:
  - The number of other persistent edges that could have been drawn





# The number of states associated with a plan

We could sample with a uniform measure over the extended state space:

$$\pi(\vec{T}, E_p) \propto 1$$

Or adapt by modding out by the number of similar plans with different persistent edges and trees

$$\pi(\vec{T}, E_p) \propto \left[ \left[ \prod_{d \in D} \frac{1}{\tau(g_C(d))} \prod_{c \in C} \frac{1}{\tau(g_p(c, d))} \cdots \right] \prod_{e \in E_p} \frac{1}{|\mathcal{P}(e, \vec{T})|} \right]^\gamma$$

$\vec{T}$ : the forest       $D$ : district set       $g_C(d)$ : county graph restricted to district

$E_p$ : persistent edges       $C$ : county set       $g_p(c, d)$ : the precinct graph of a county and district

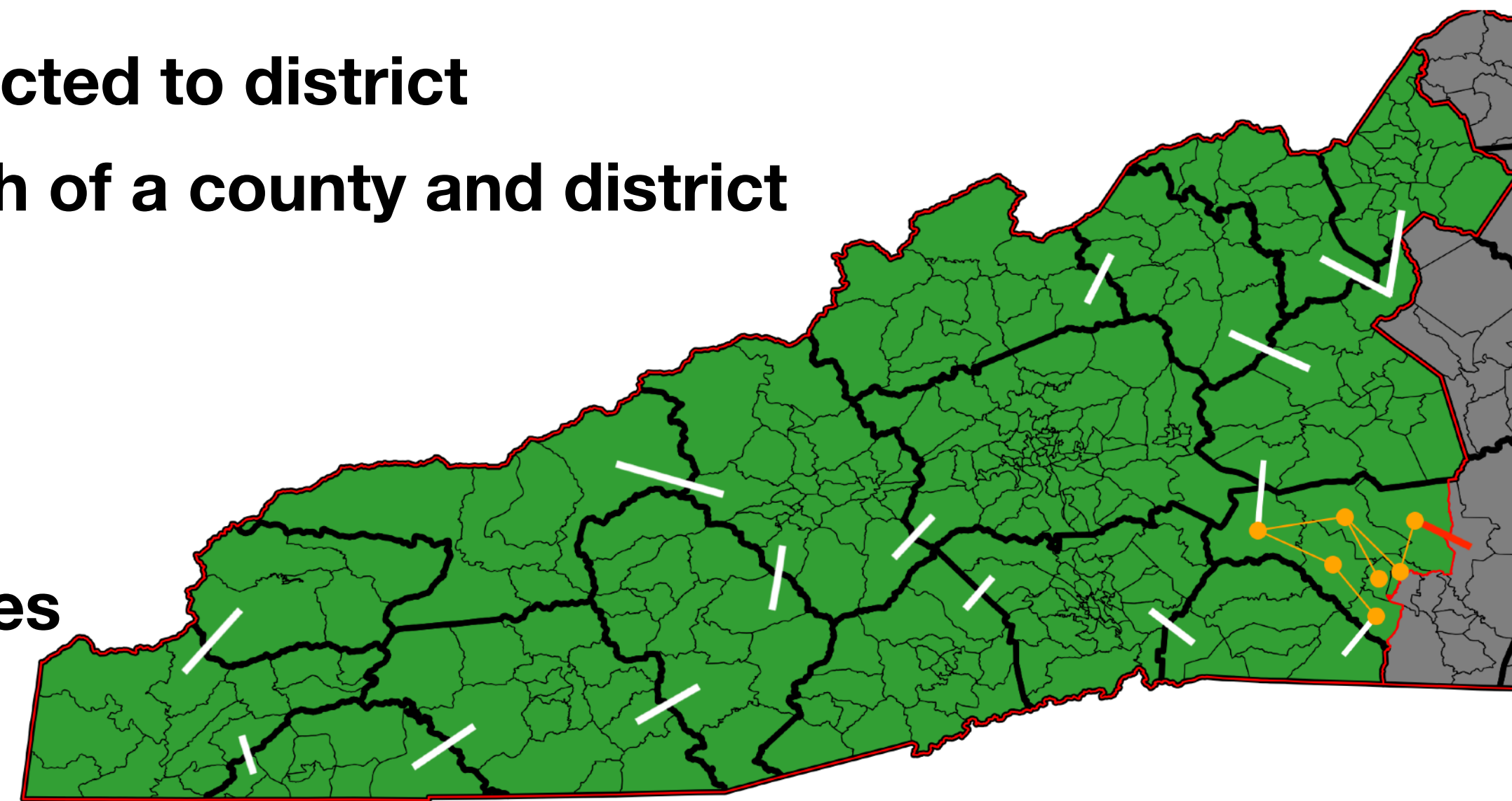
$\tau$ : number of spanning trees

$\mathcal{P}$ : set of possible persistent edges

$\gamma \in [0, 1]$

$\gamma = 0$ : Uniform over product space of trees and persistent edges

$\gamma = 1$ : Uniform over partitions

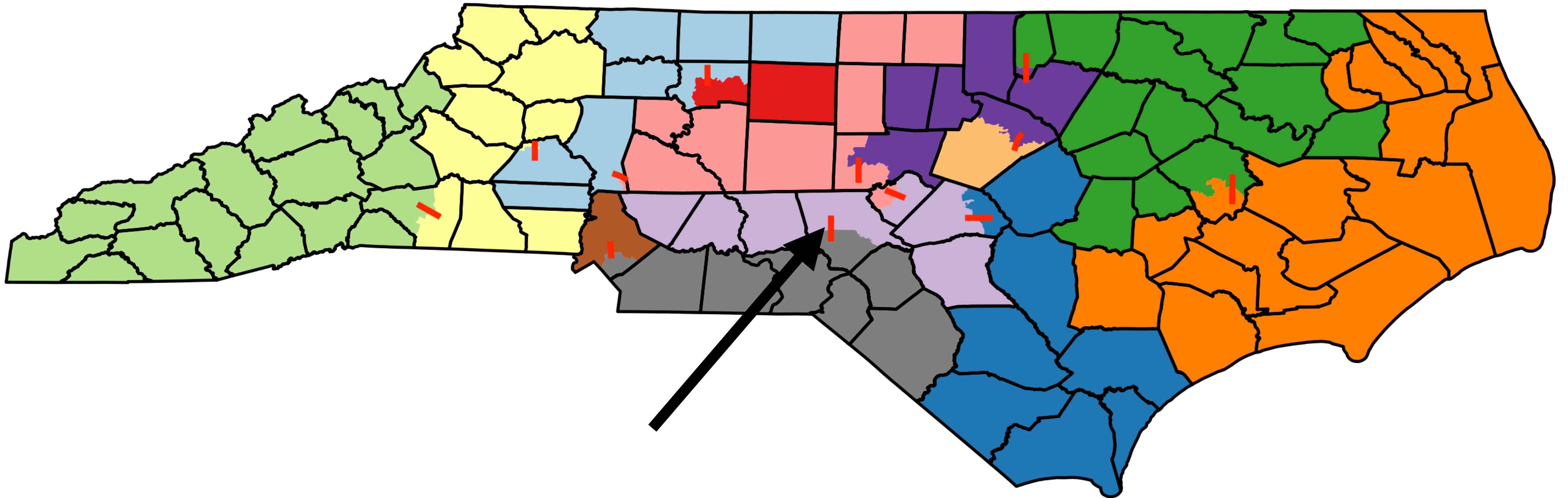


# The proposal

1. Choose a persistent edge  $\propto \frac{1}{|E_p|}$
2. Merge the district county graphs and draw a uniform tree on the resulting multigraph  $\propto \frac{1}{\tau(g_c(d1 \cup d2))}$
3. Find edges and nodes that can be cut
4. On each cuttable node, draw a new uniform tree on the next level down (specify coarse edges where needed) and repeat steps 3 and 4  $\propto \frac{1}{\tau(g_p(c, d1 \cup d2))}$
5. Aggregate all edges that can be cut across all levels.  
Pick one uniformly; this is the new persistent edge.  $\propto \frac{1}{|E_{cut}|}$

# 1. Choose a Persistent Edge

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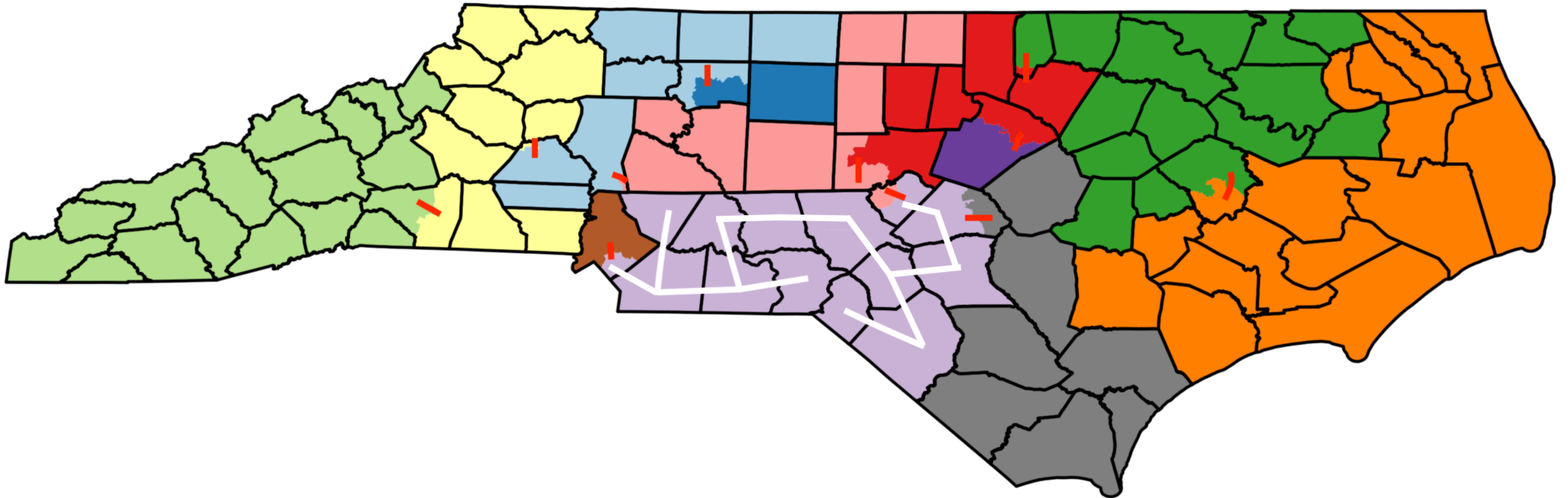




## 2. Merge districts; draw tree at the County level

2. Merge the district county graphs and draw a uniform tree

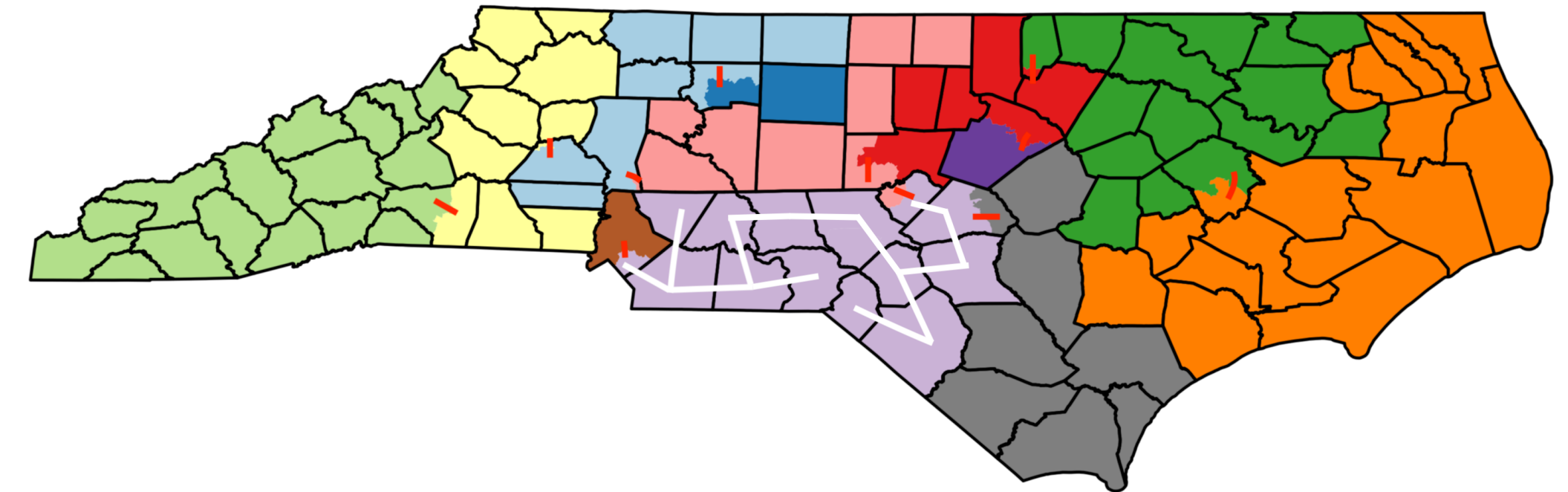
on the resulting multigraph  $\propto \frac{1}{\tau(g_c(d1 \cup d2))}$





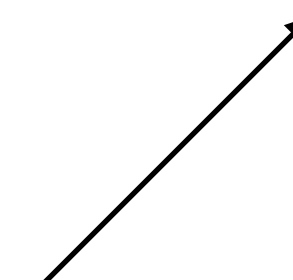
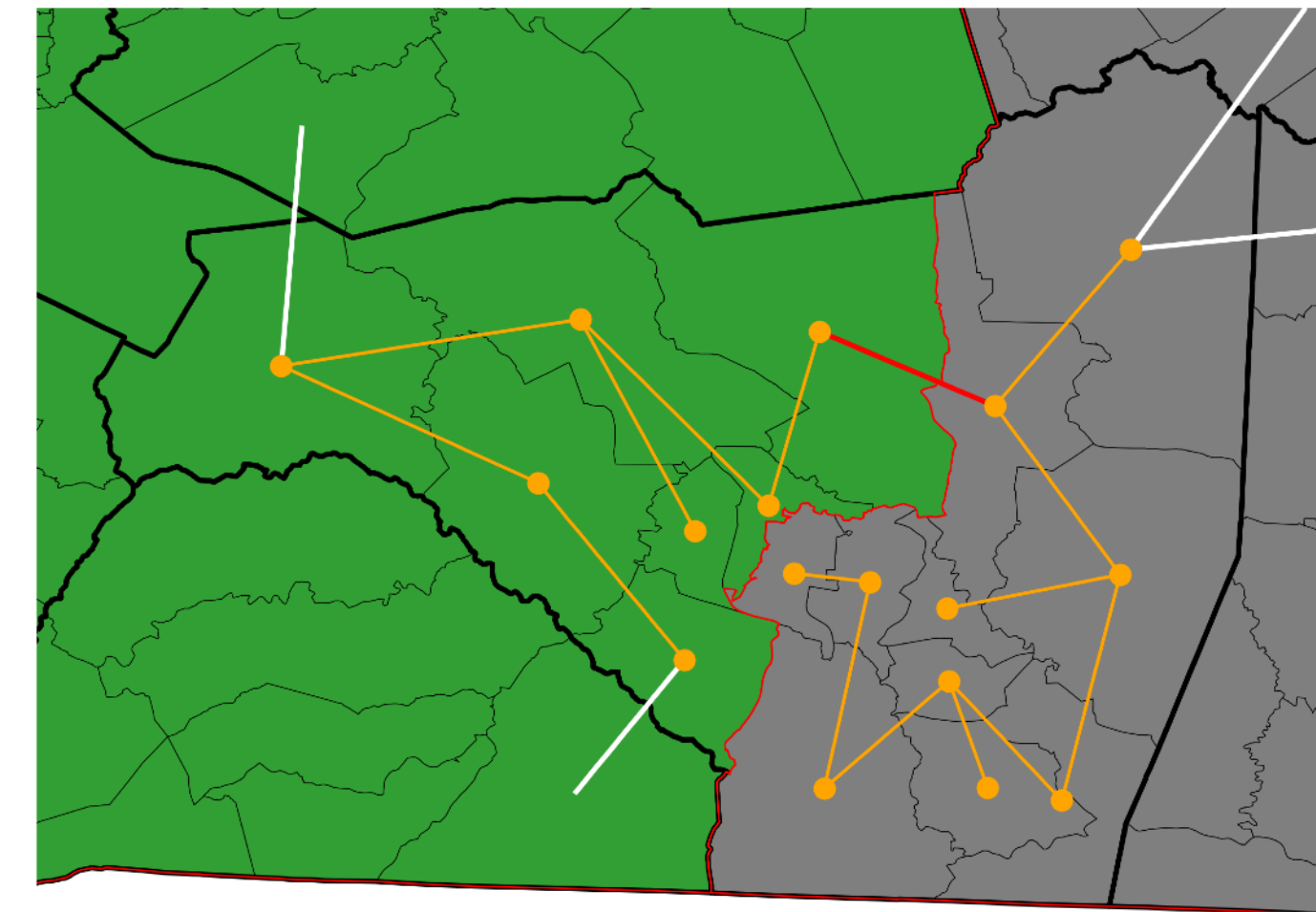
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# The proposal (computational acceleration)

- Even though we have only explicitly drawn trees on the cuttable, nodes, we have implicitly drawn them everywhere.
- The persistent edge gives a tree on the merged graph. We no longer need to iterate over all other trees!

$$\frac{Q(T_1, T_2, e)}{Q(T'_1, T'_2, e')} = \frac{|E'_{cut}|(T'_1, T'_2, e')}{|E_{cut}|(T_1, T_2, e)}$$

**With persistent edges**

$$\frac{Q(T_1, T_2, e)}{Q(T'_1, T'_2, e')} = \frac{\sum_{e \in \mathcal{P}(T_1, T_2)} \frac{1}{|E_{cut}|(T_1, T_2, e)}}{\sum_{e \in \mathcal{P}(T'_1, T'_2)} \frac{1}{|E_{cut}|(T'_1, T'_2, e)}}$$

**Without persistent edges**



# Probability Ratio

$$\pi(\vec{T}, E_p) \propto \left[ \left[ \prod_{d \in D} \frac{1}{\tau(g_c(d))} \prod_{c \in C} \frac{1}{\tau(g_p(c, d))} \cdots \right] \prod_{e \in E_p} \frac{1}{|\mathcal{P}(e, \vec{T})|} \right]^\gamma$$

$$\frac{Q(T_1, T_2, e)}{Q(T'_1, T'_2, e')} = \frac{|E'_{cut}|(T'_1, T'_2, e')}{|E_{cut}|(T_1, T_2, e)}$$

$$\frac{\pi(\vec{T}, E_p)}{\pi(\vec{T}', E'_p)} = \frac{\tau(g_c(T'_1))\tau(g_c(T'_2))}{\tau(g_c(T_1))\tau(g_c(T_2))} \frac{\tau(g_p(c_{n'}, d'_1))\tau(g_p(c_{n'}, d'_2))}{\tau(g_p(c_{n'}))} \frac{\tau(g_p(c_n))}{\tau(g_p(c_n, d_1))\tau(g_p(c_n, d_2))} \frac{|\mathcal{P}(e')|}{|\mathcal{P}(e)|} \prod_{e_i \in E'_p \setminus e} \frac{|\mathcal{P}(e_i, \vec{T}')|}{|\mathcal{P}(e_i, \vec{T})|}$$

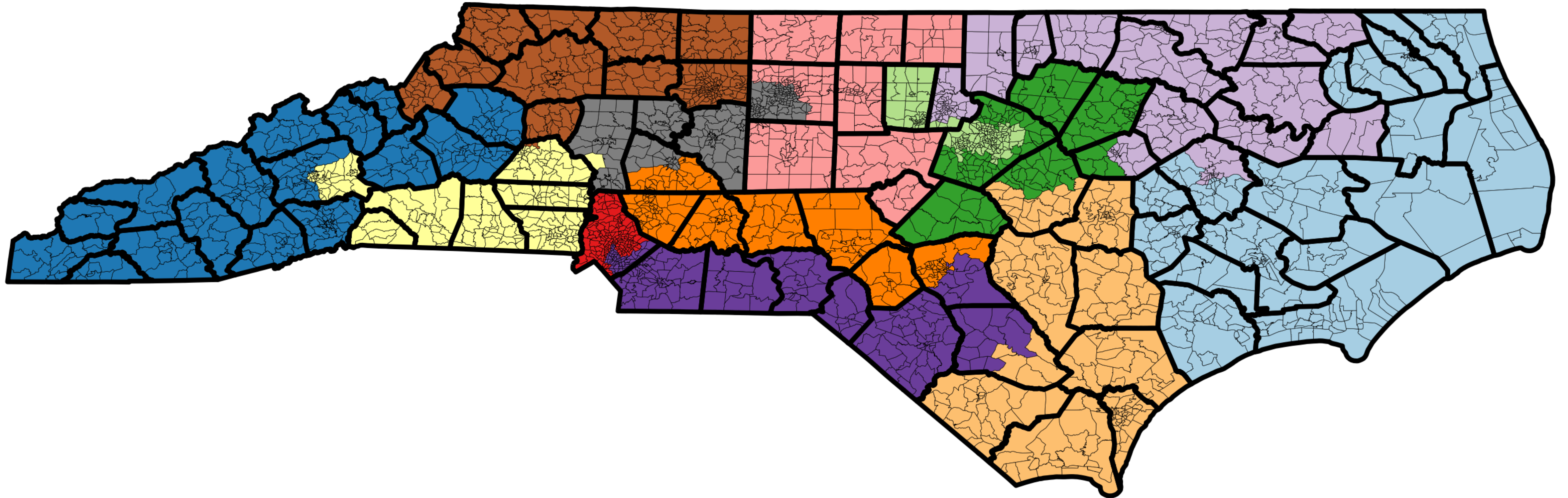
$c'_n, c_n$       **Cut counties on new and old districts, resp.**

$e', e$       **Persistent edge for new and old districts, resp.**

# Remarks

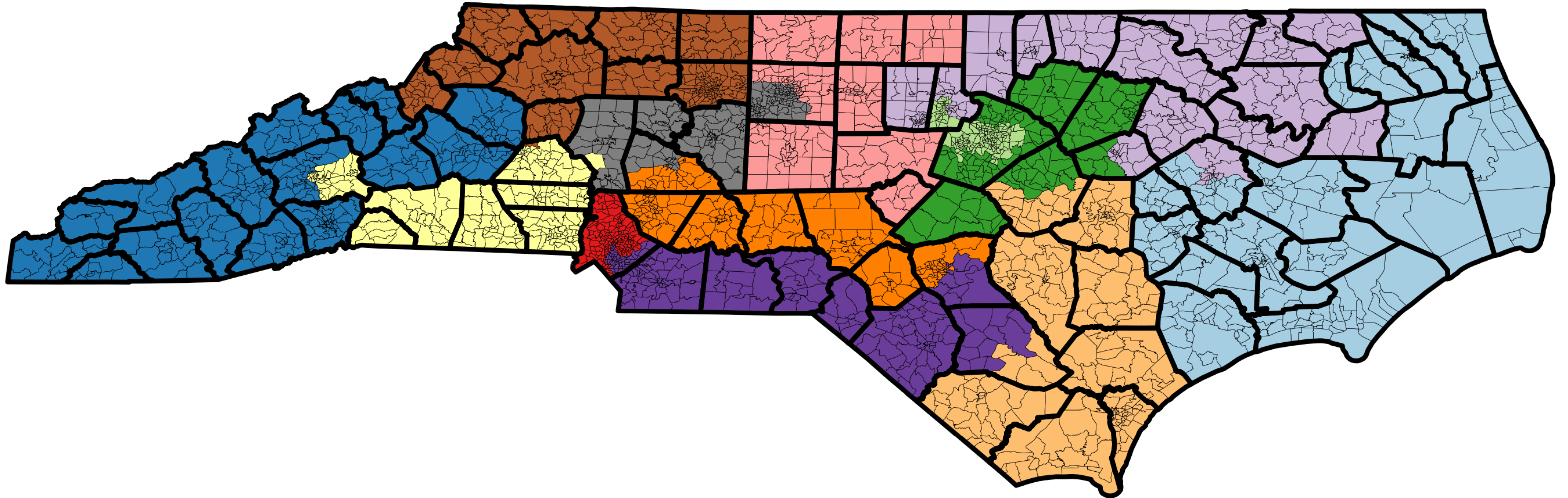
- Must ensure there are not two persistent edges linking the same two districts.
- The number of cut counties is bounded, from above, by the number of persistent edges.
- In our implementation, we do not allow nodes to be cut into three districts.

# Results



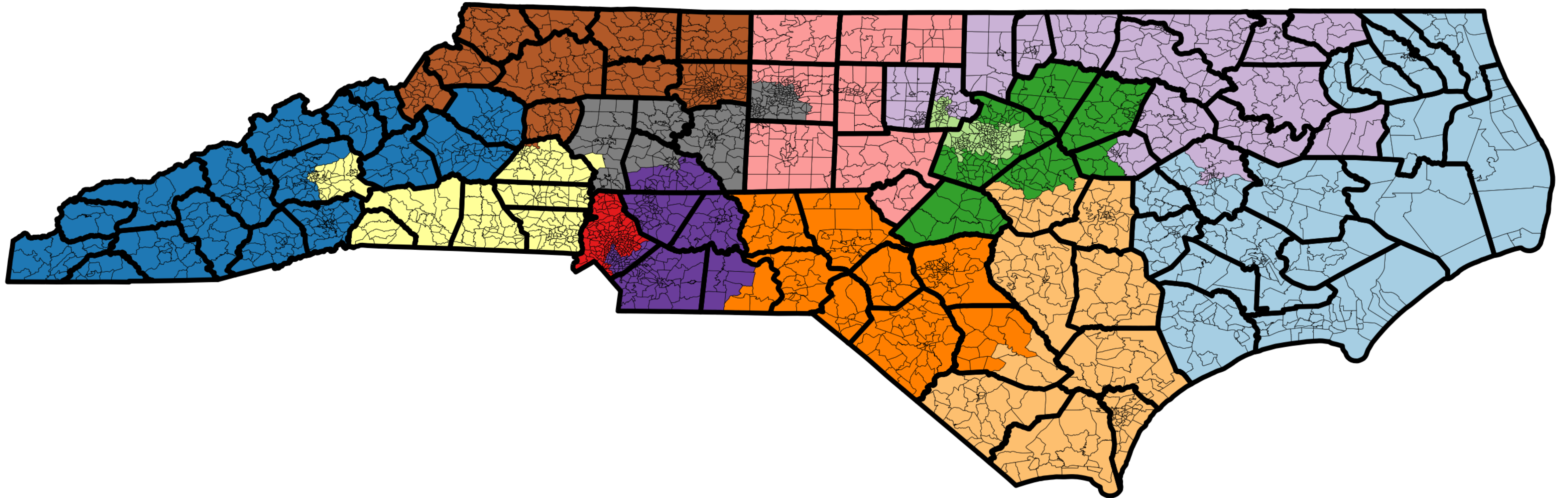


# Results



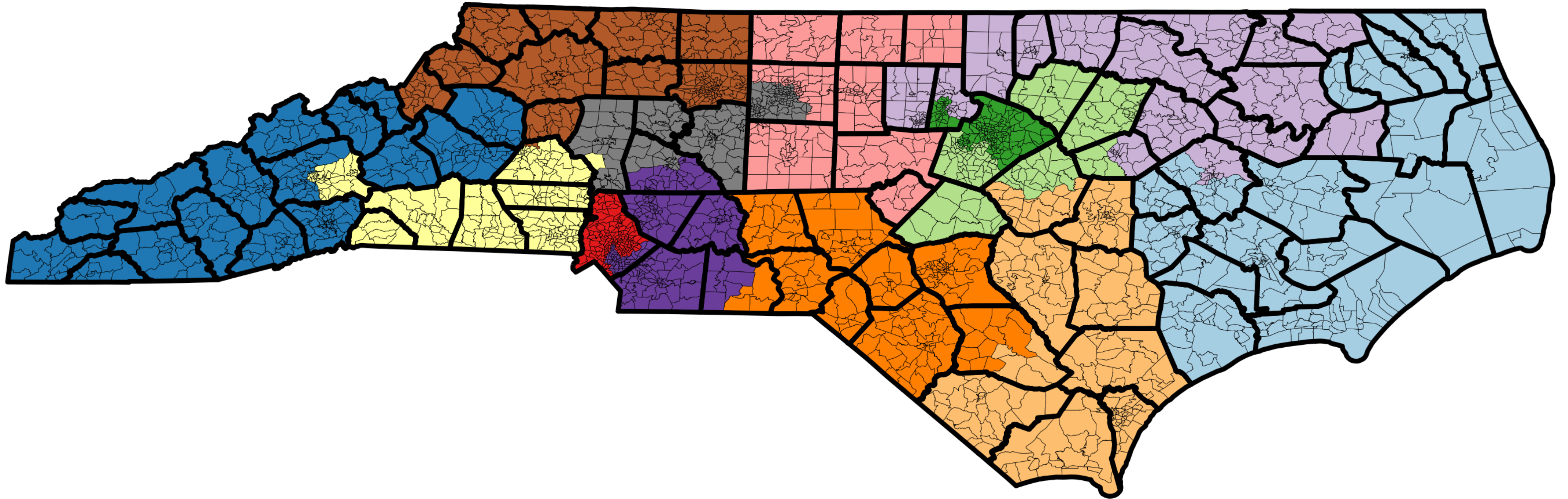


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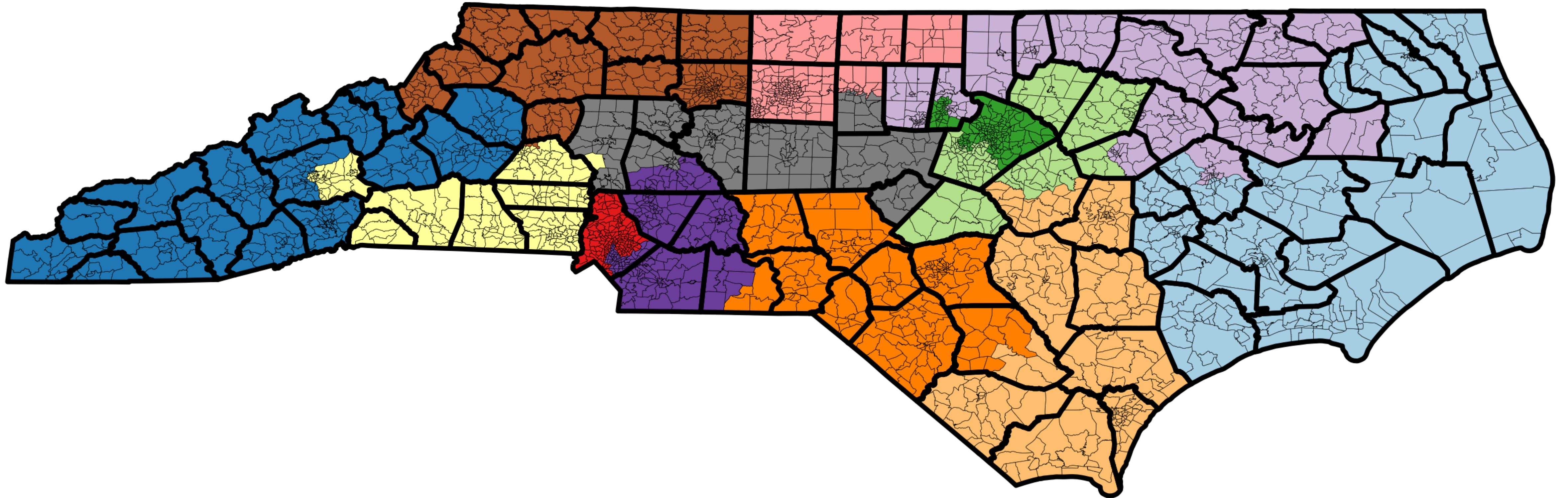


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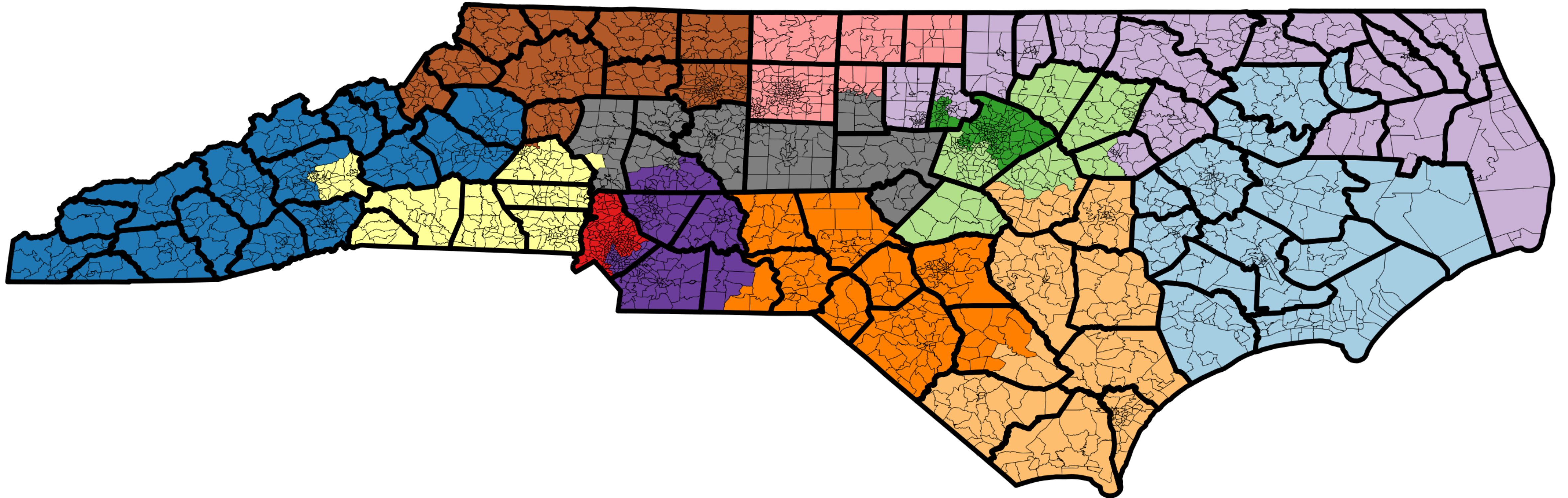


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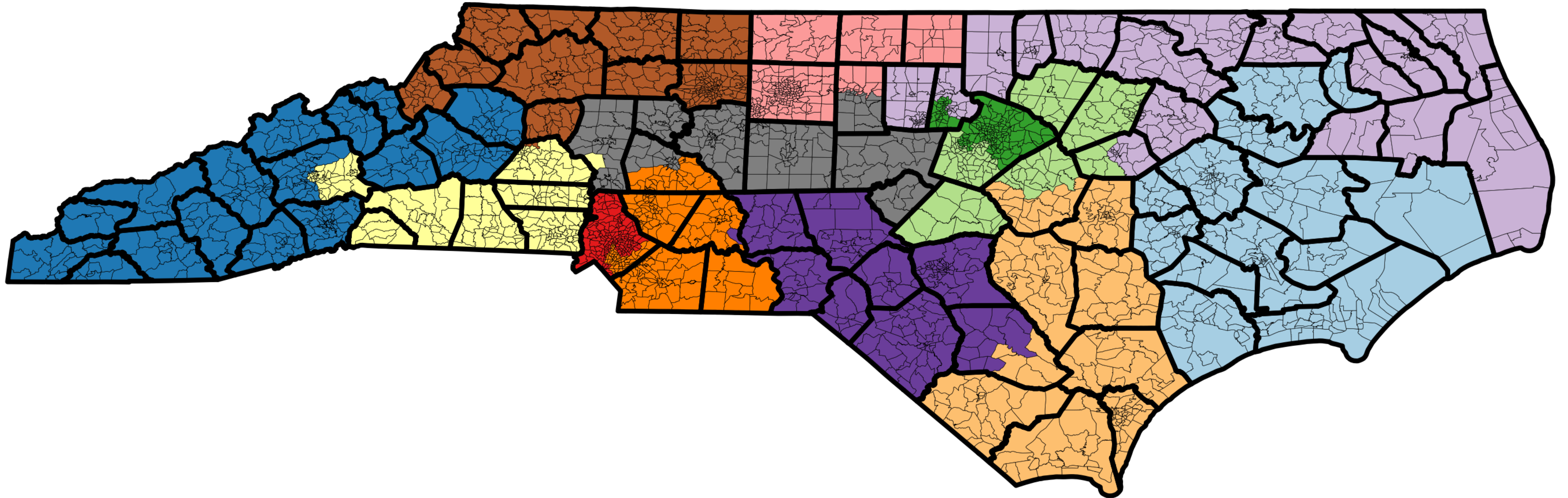


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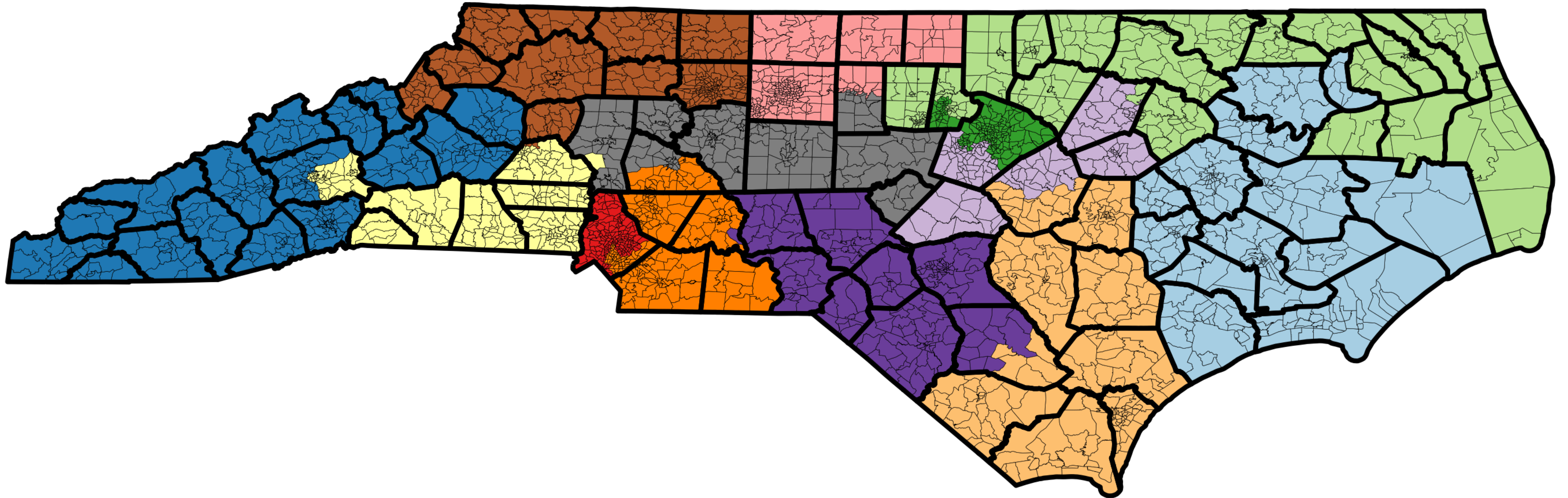


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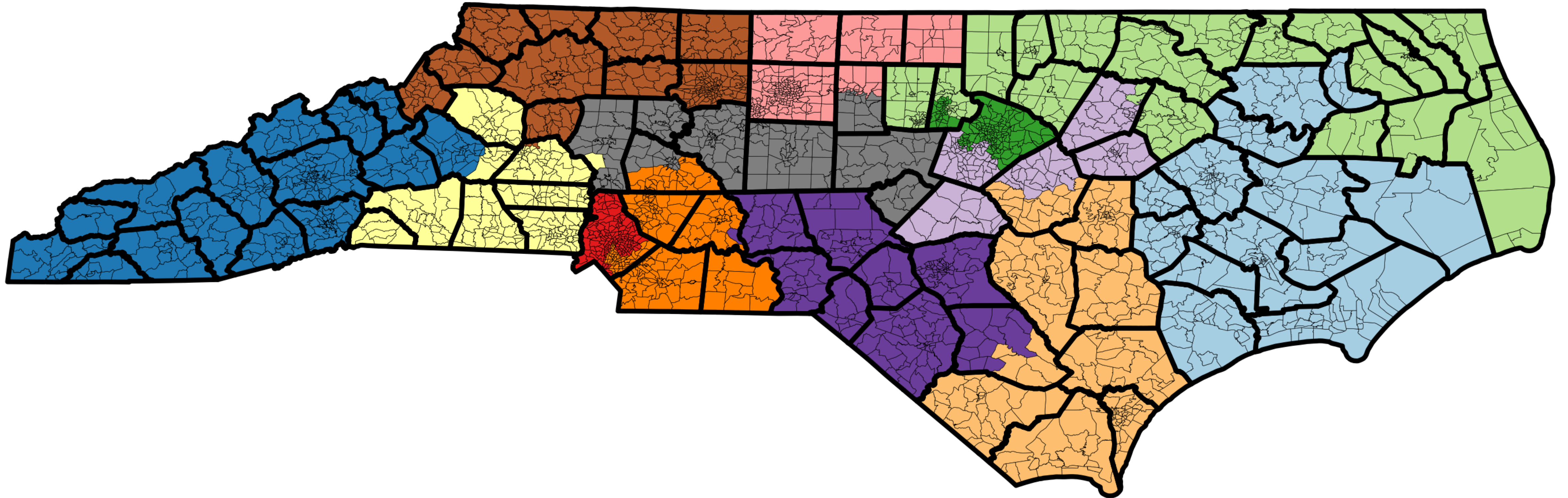


# Results





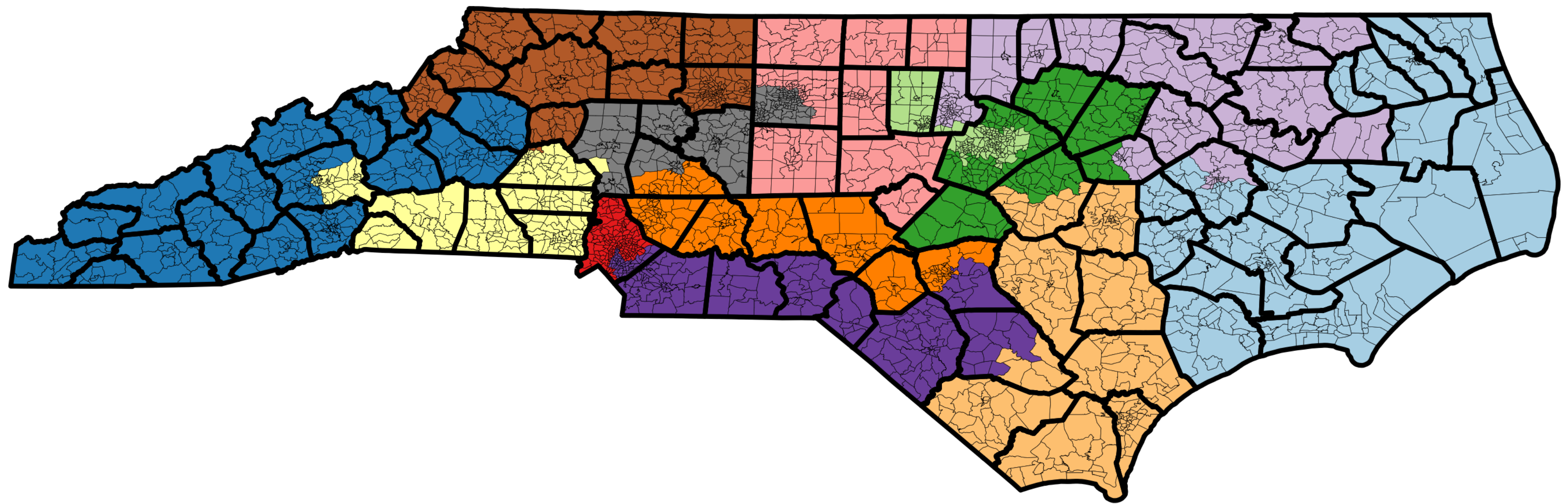
# Results



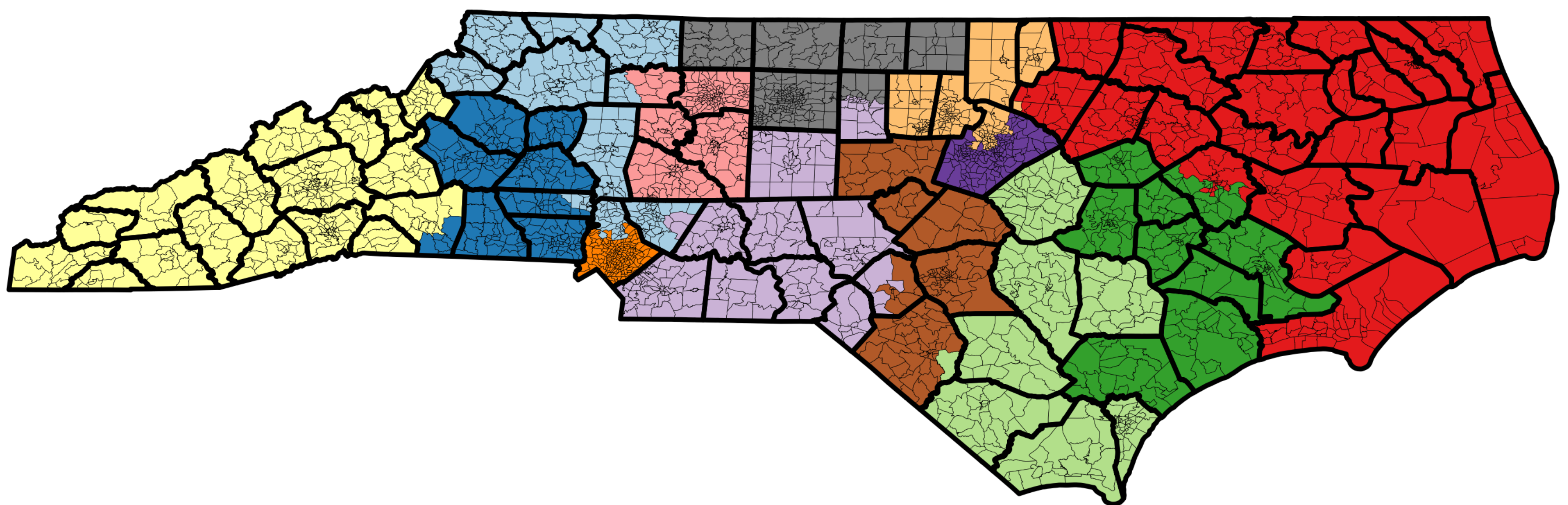


# Results

Initial



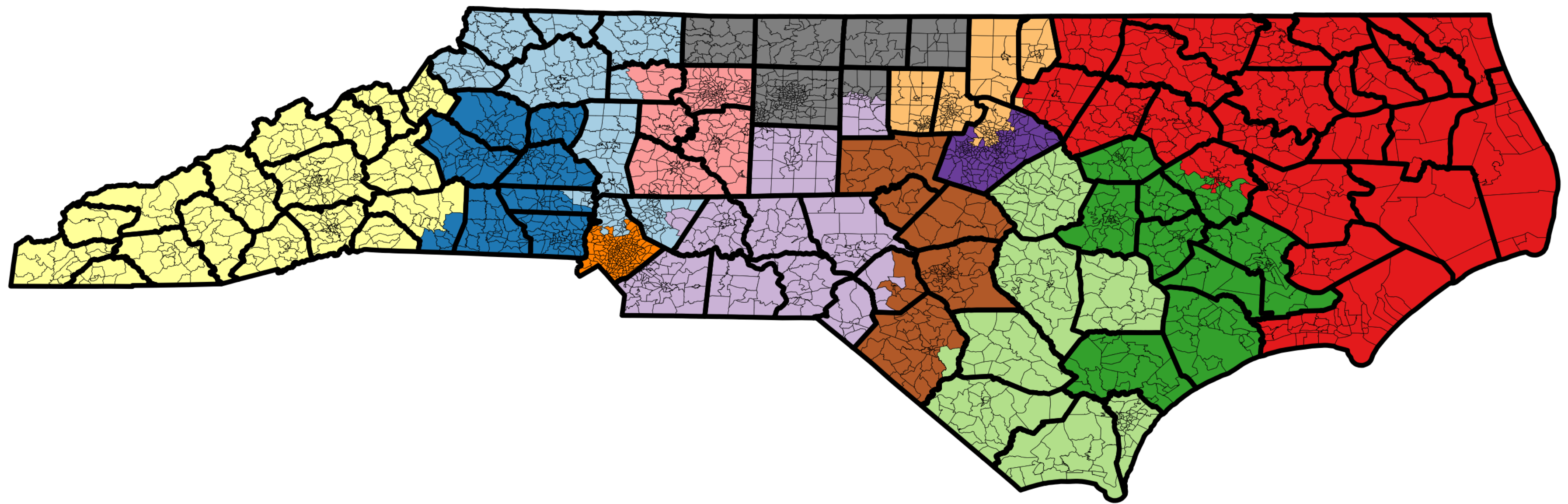
10K Proposals



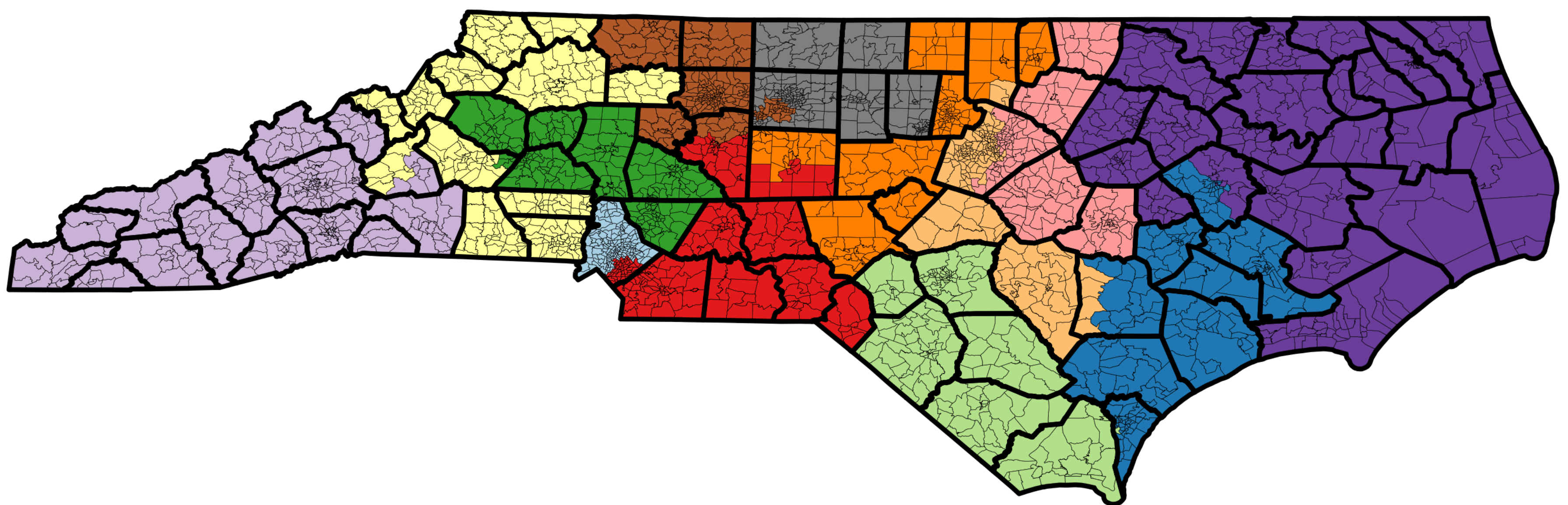


# Results

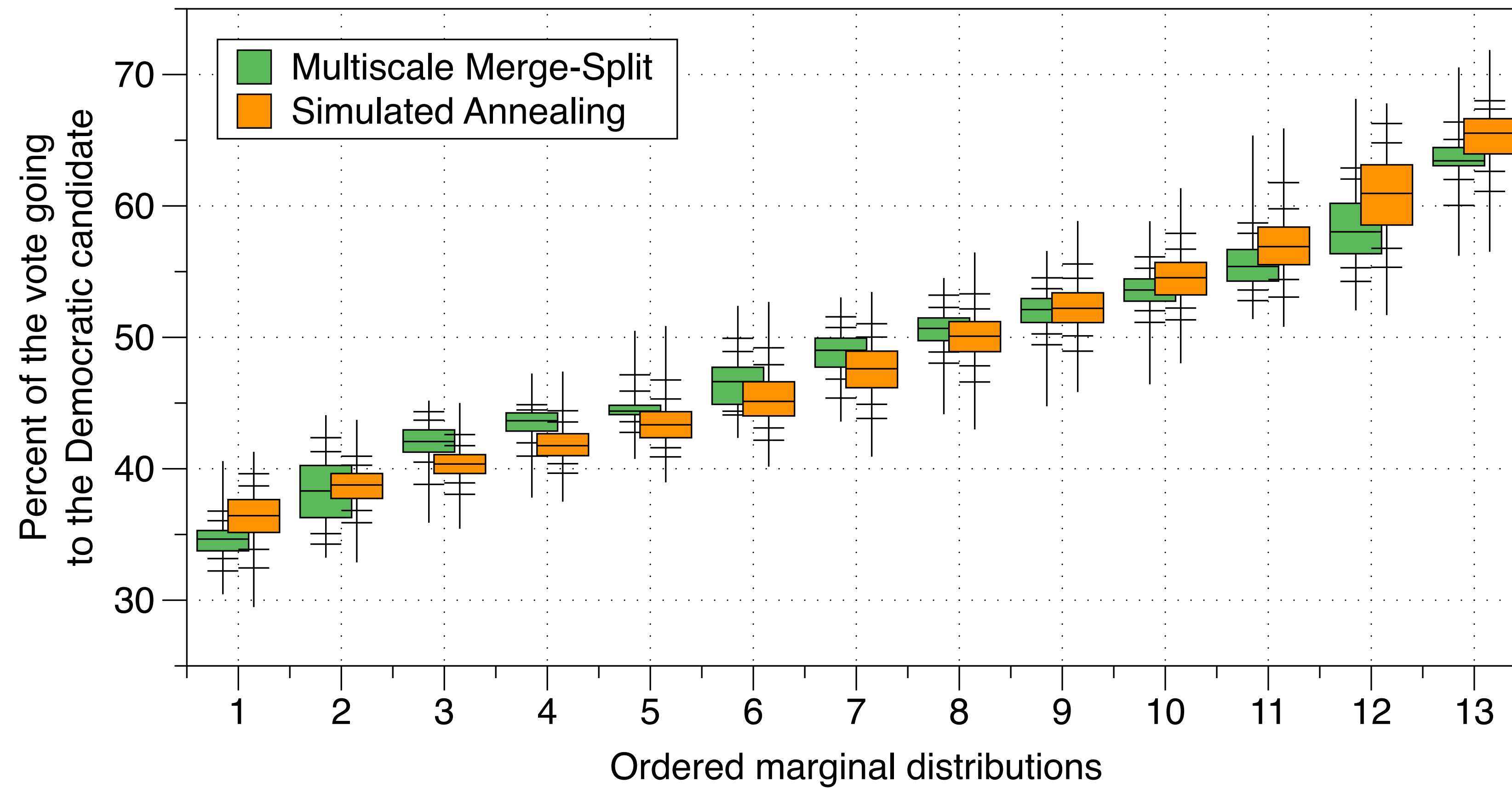
10K Proposals



15K Proposals



# Results

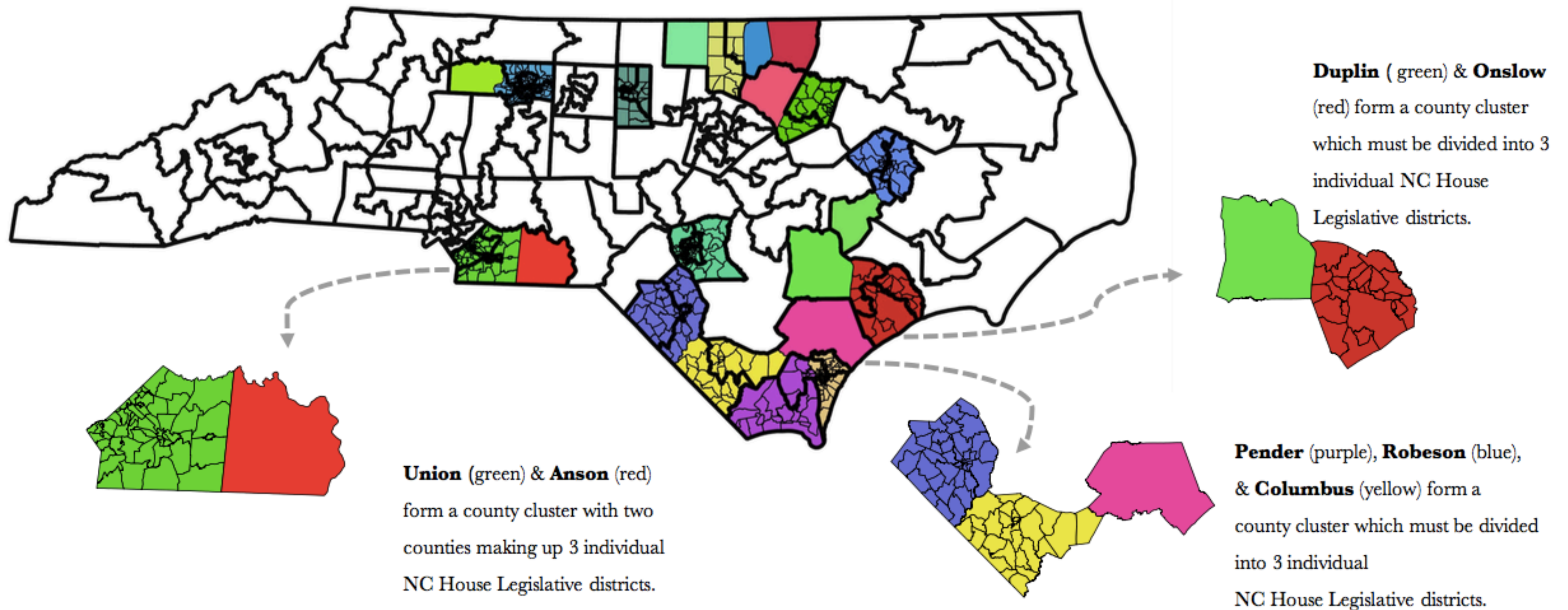




# County Clusters

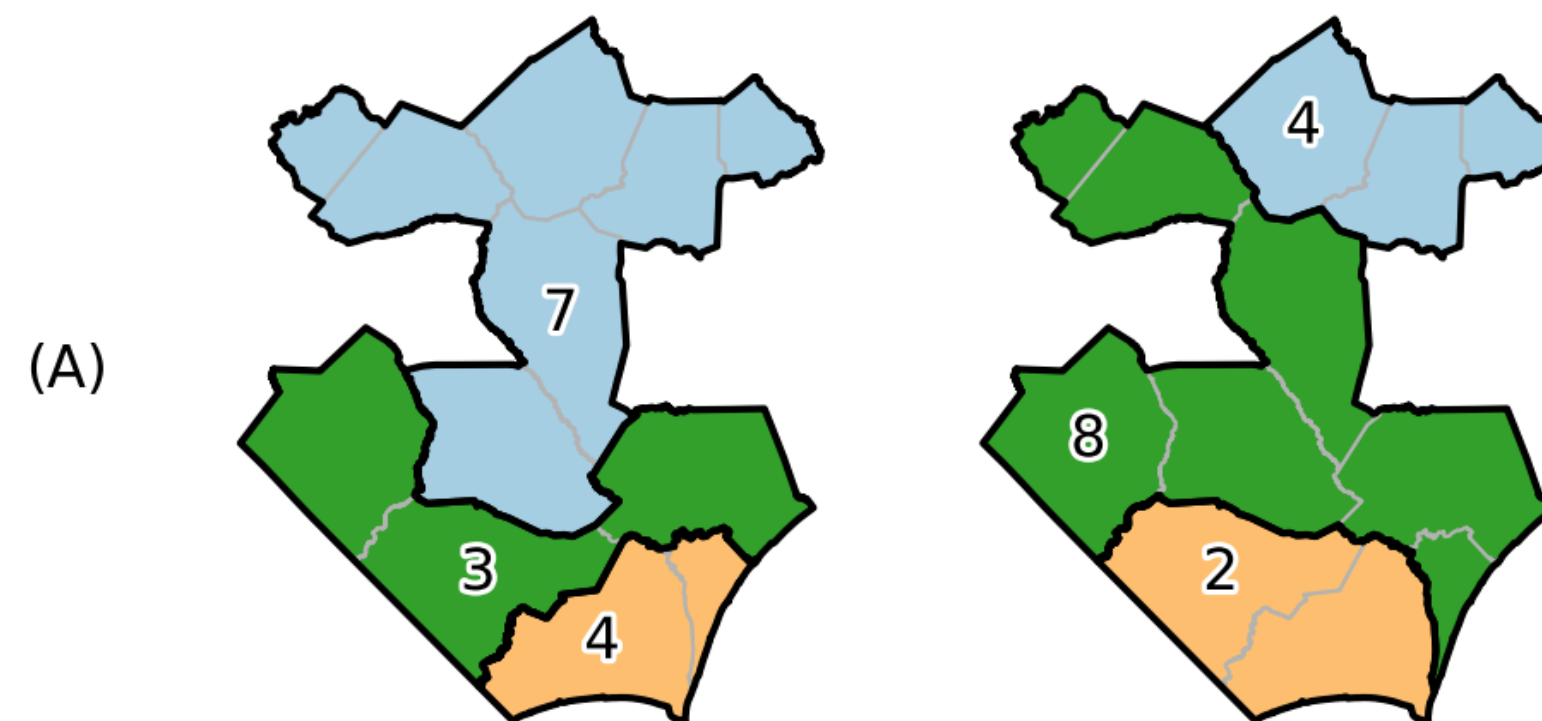
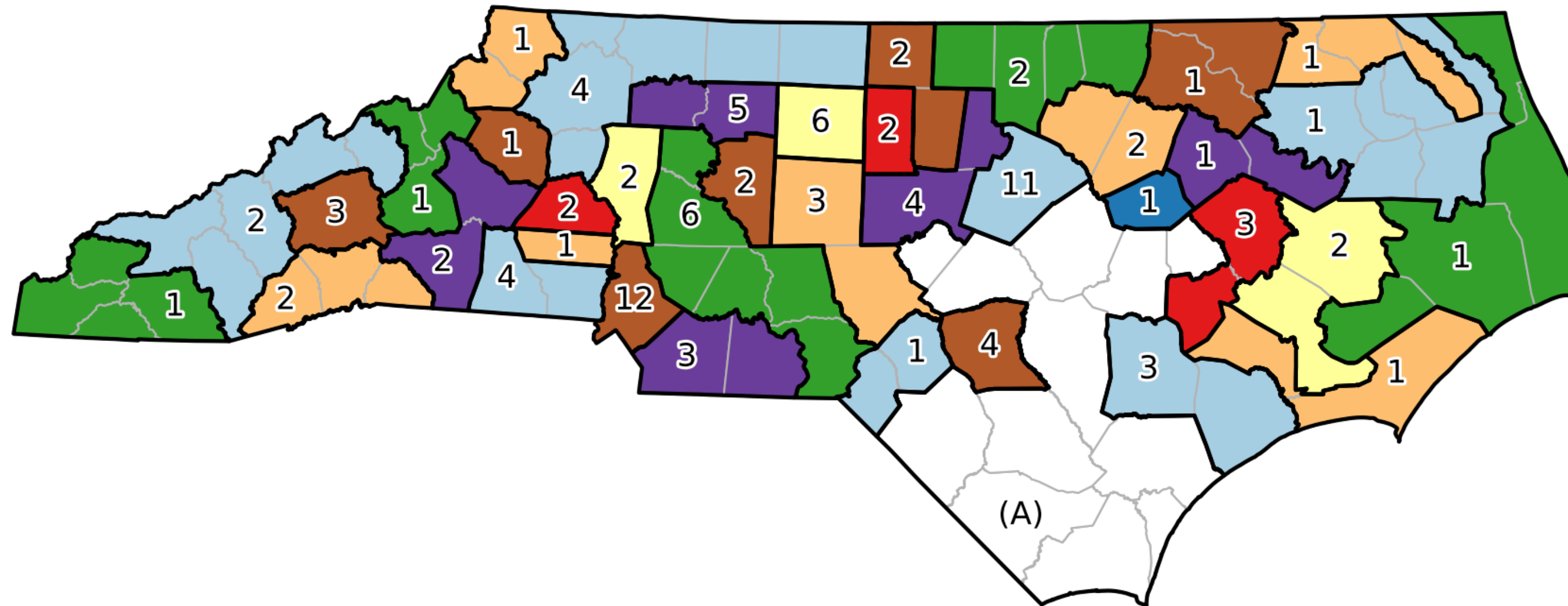


# NC Legislature



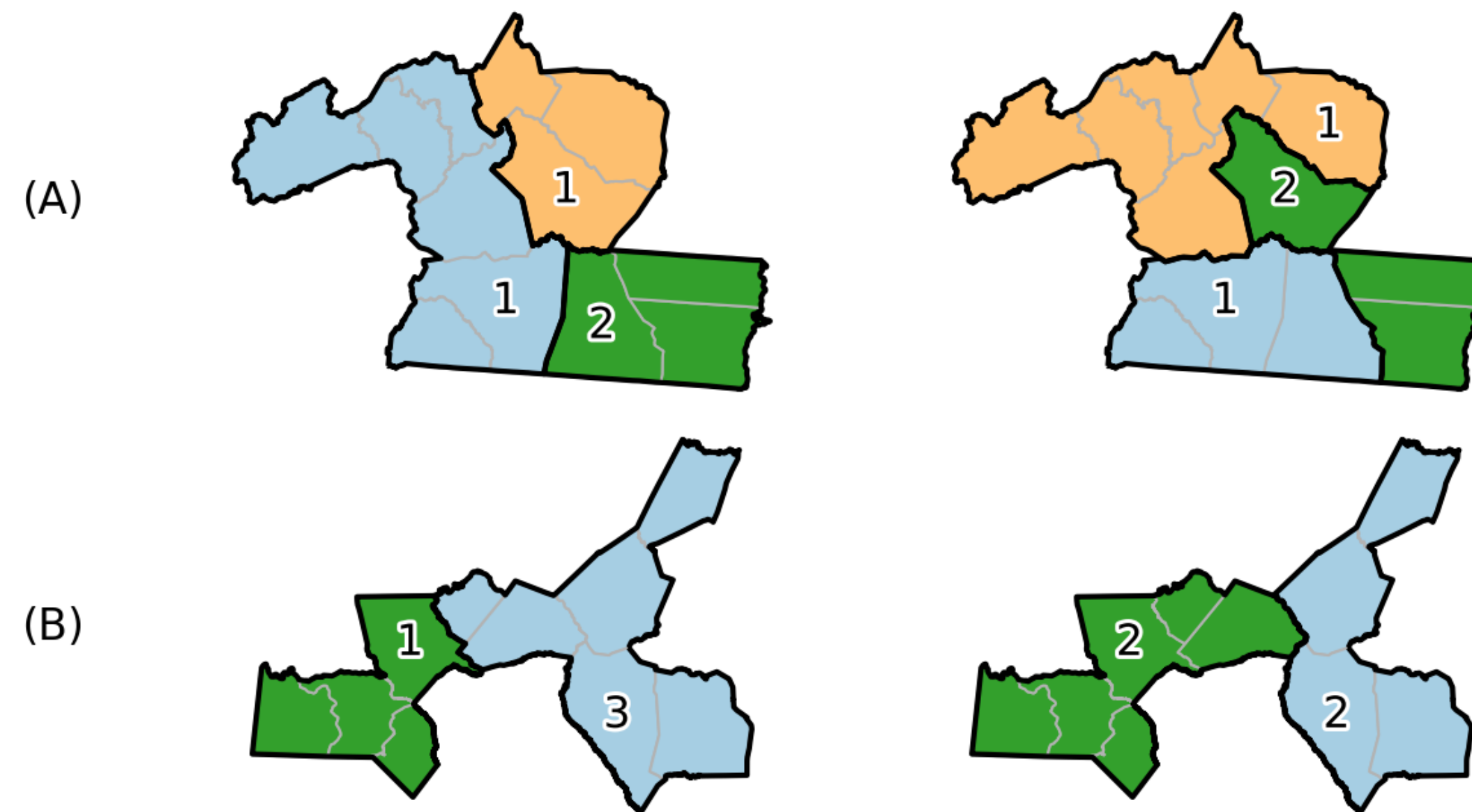
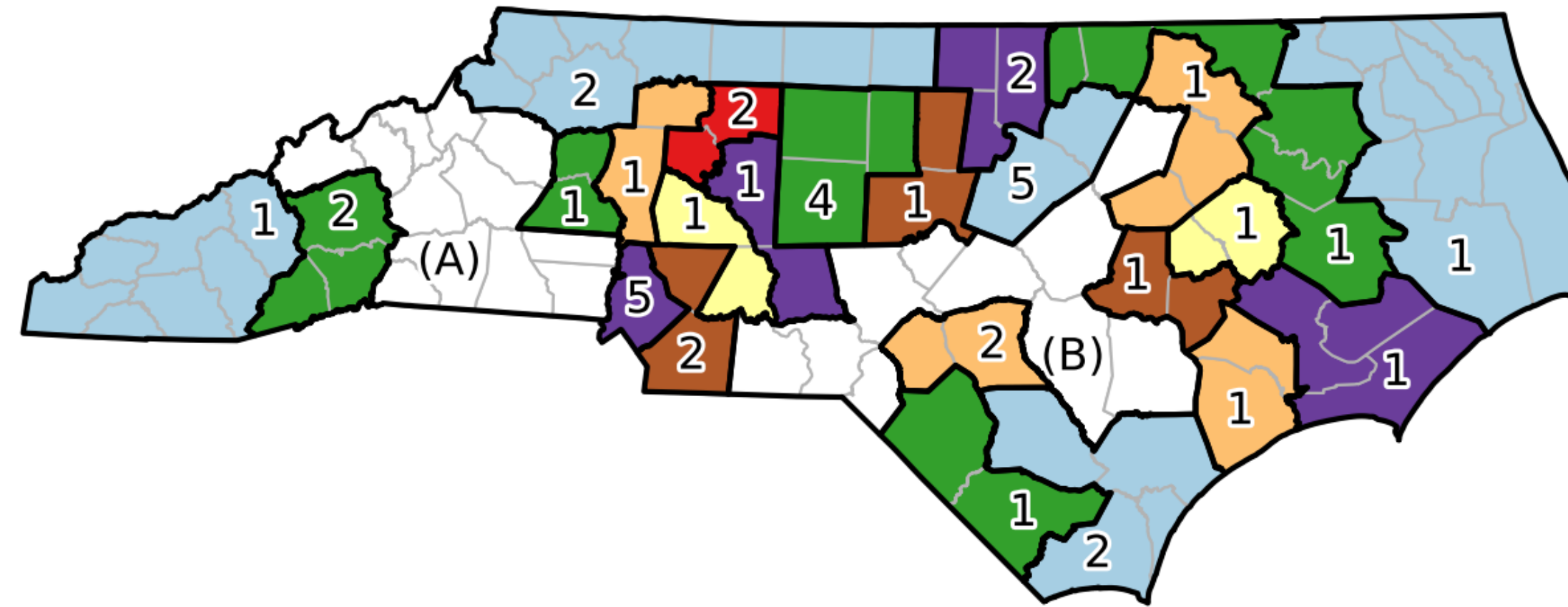
**Court order: Draw most single-county clusters, then the most two-county clusters, and so on. Maximize the number of clusters.**

# Optimal Clusters



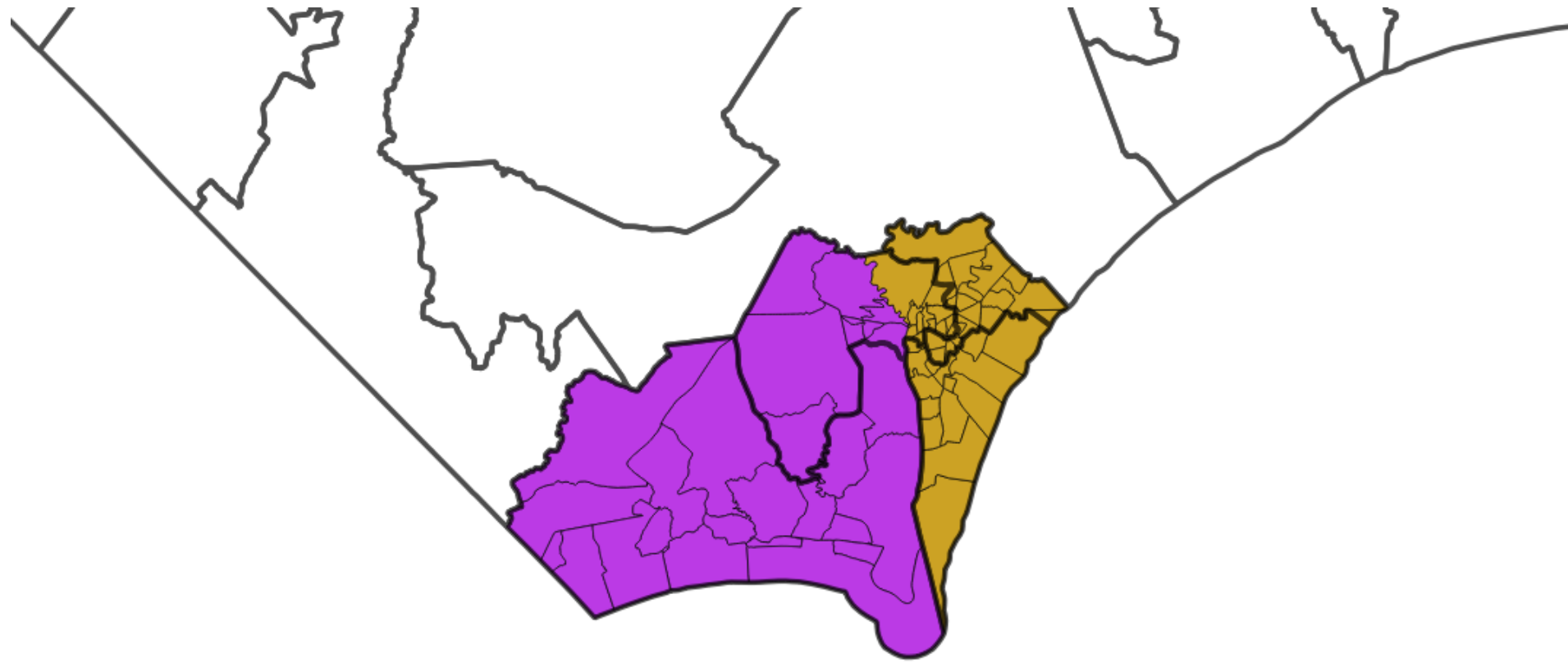


# Optimal Clusters

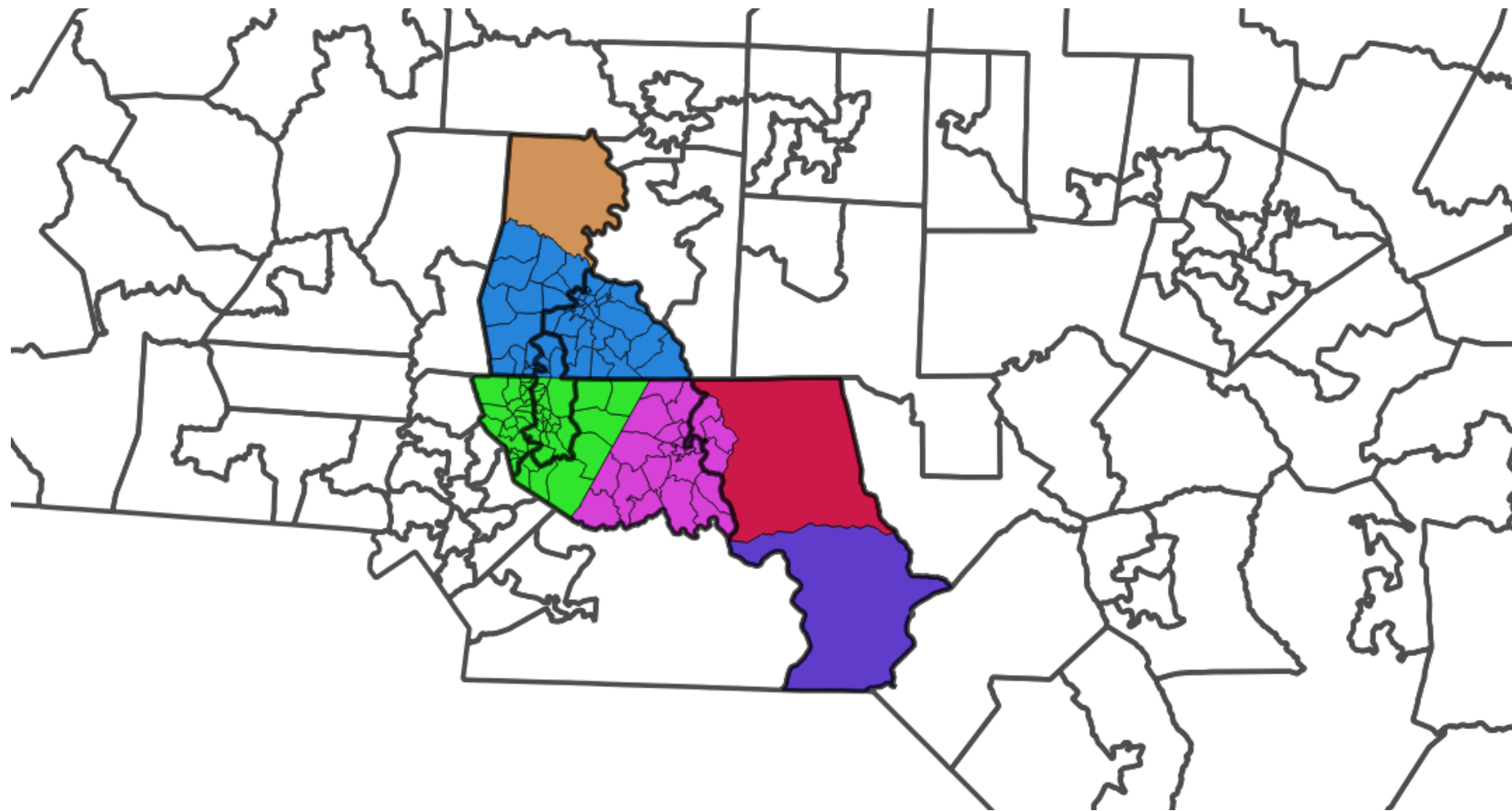




# Implicit Rules



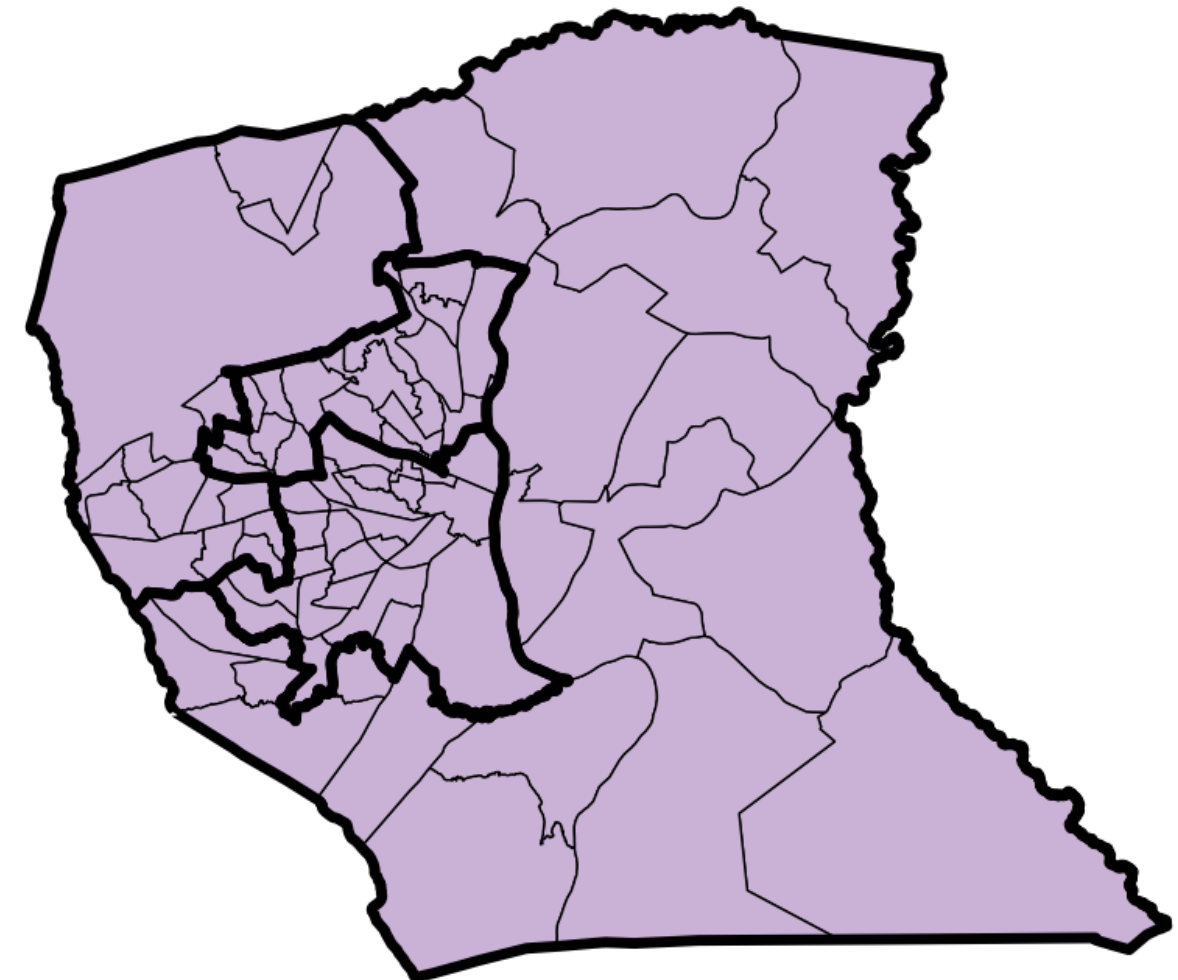
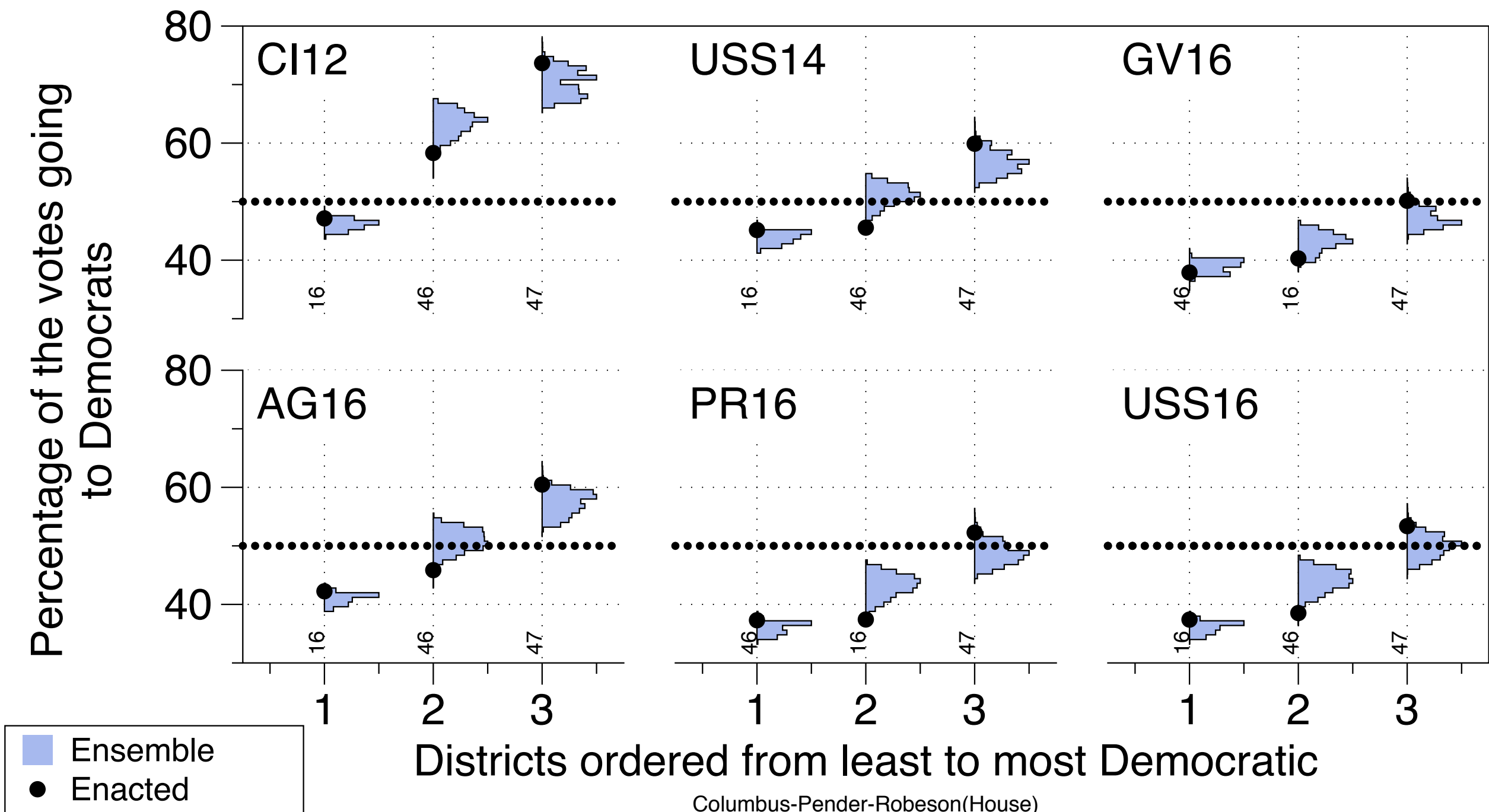
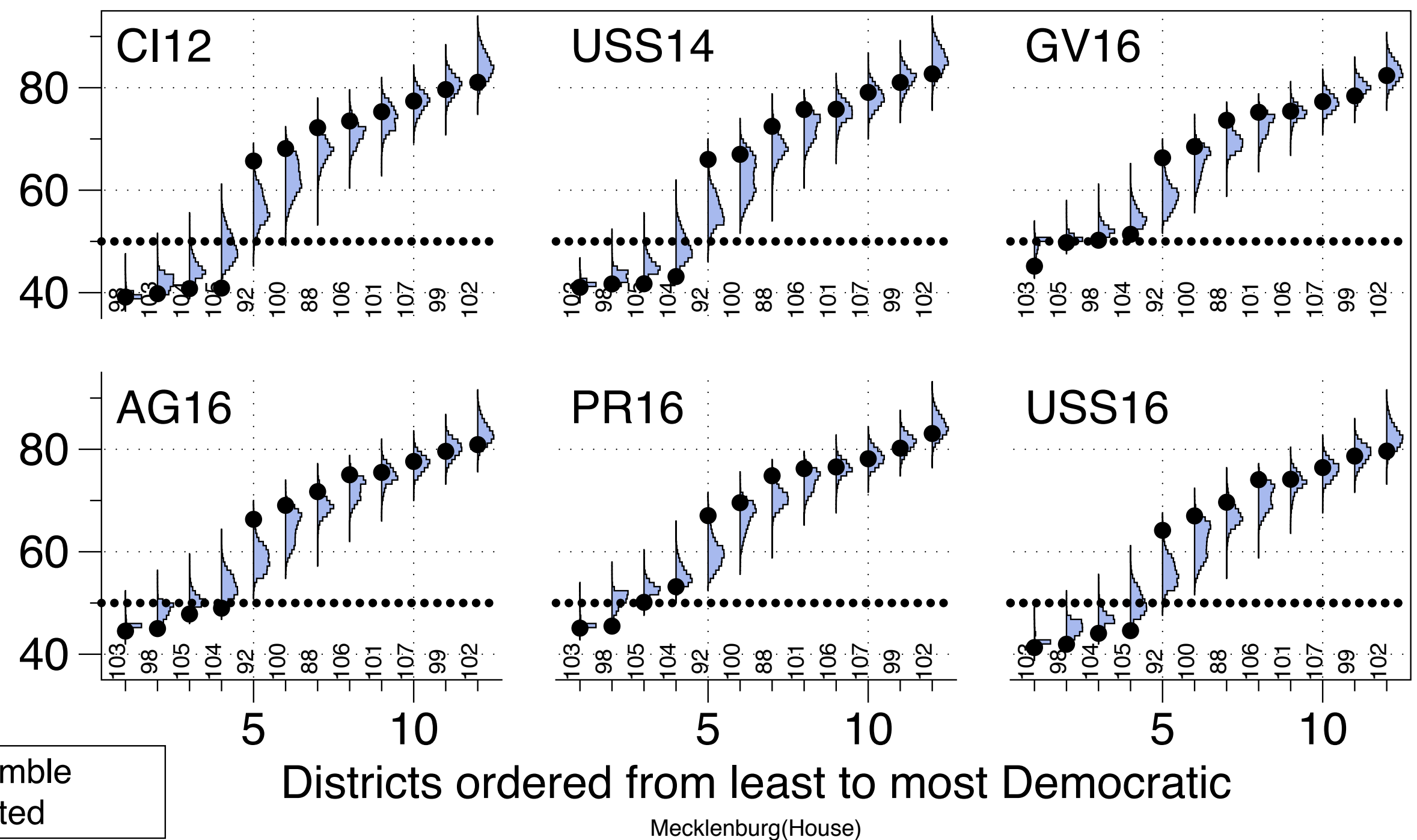
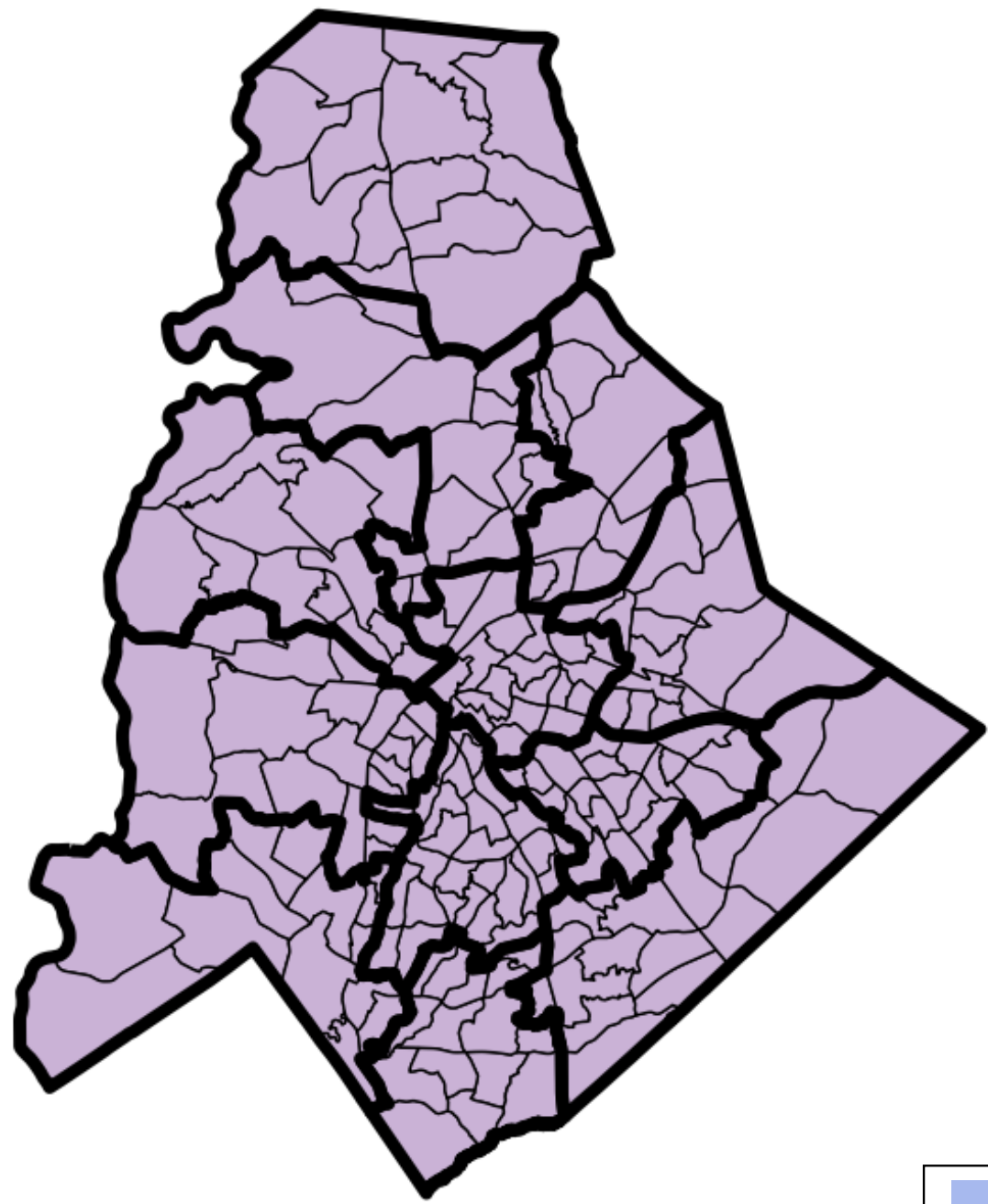
**Counties boundaries are only crossed by a single district**



**When possible districts only span two counties**

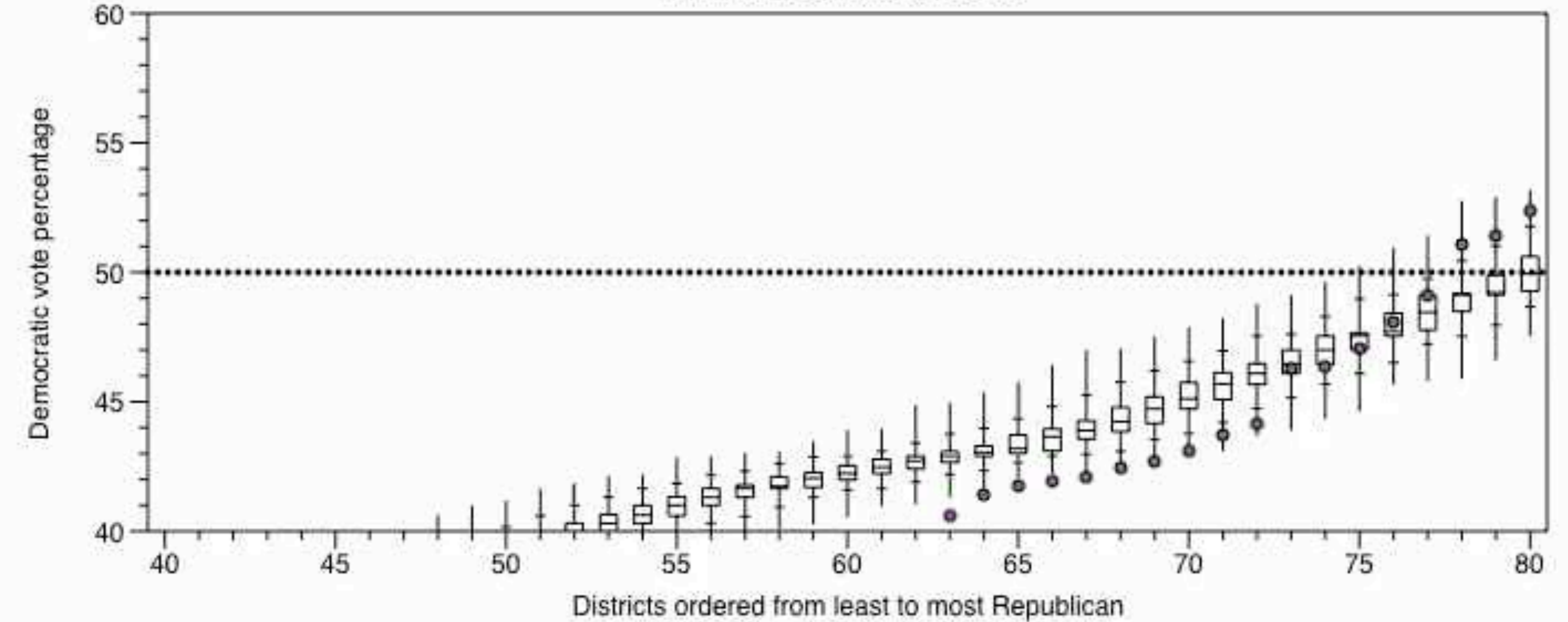
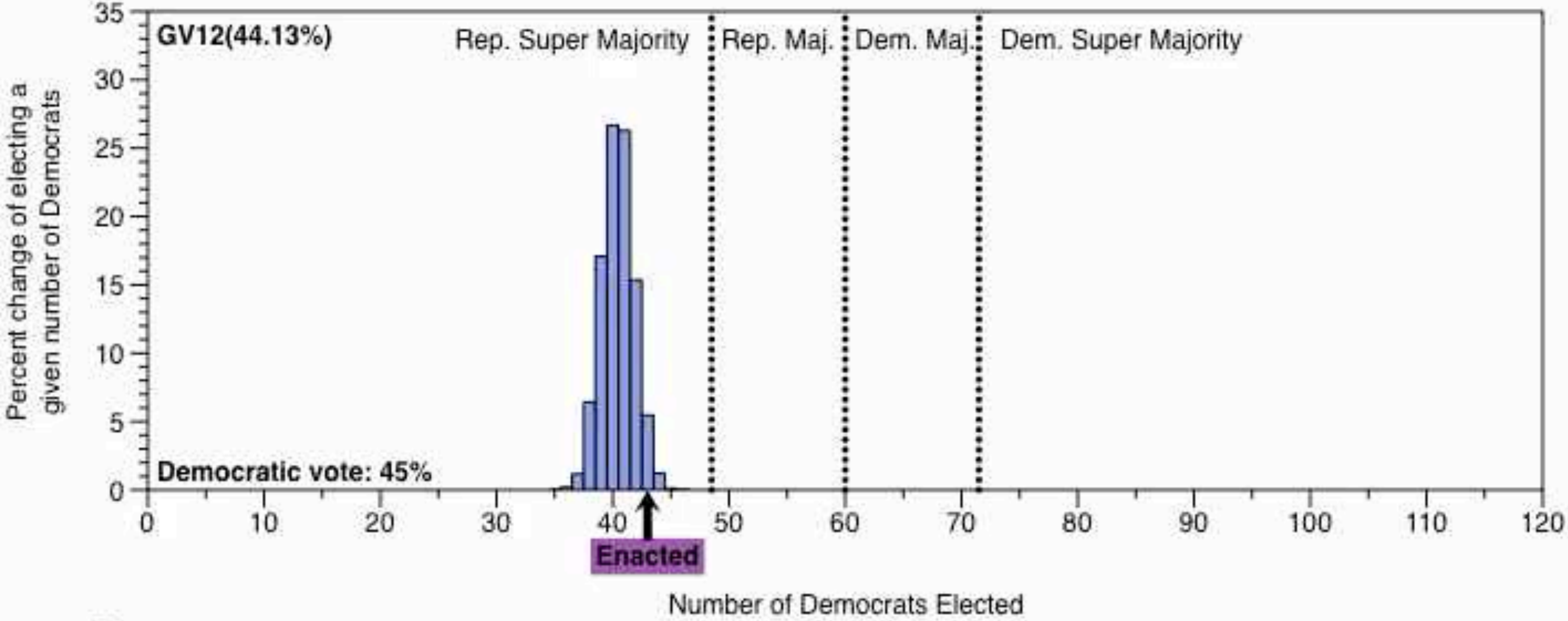
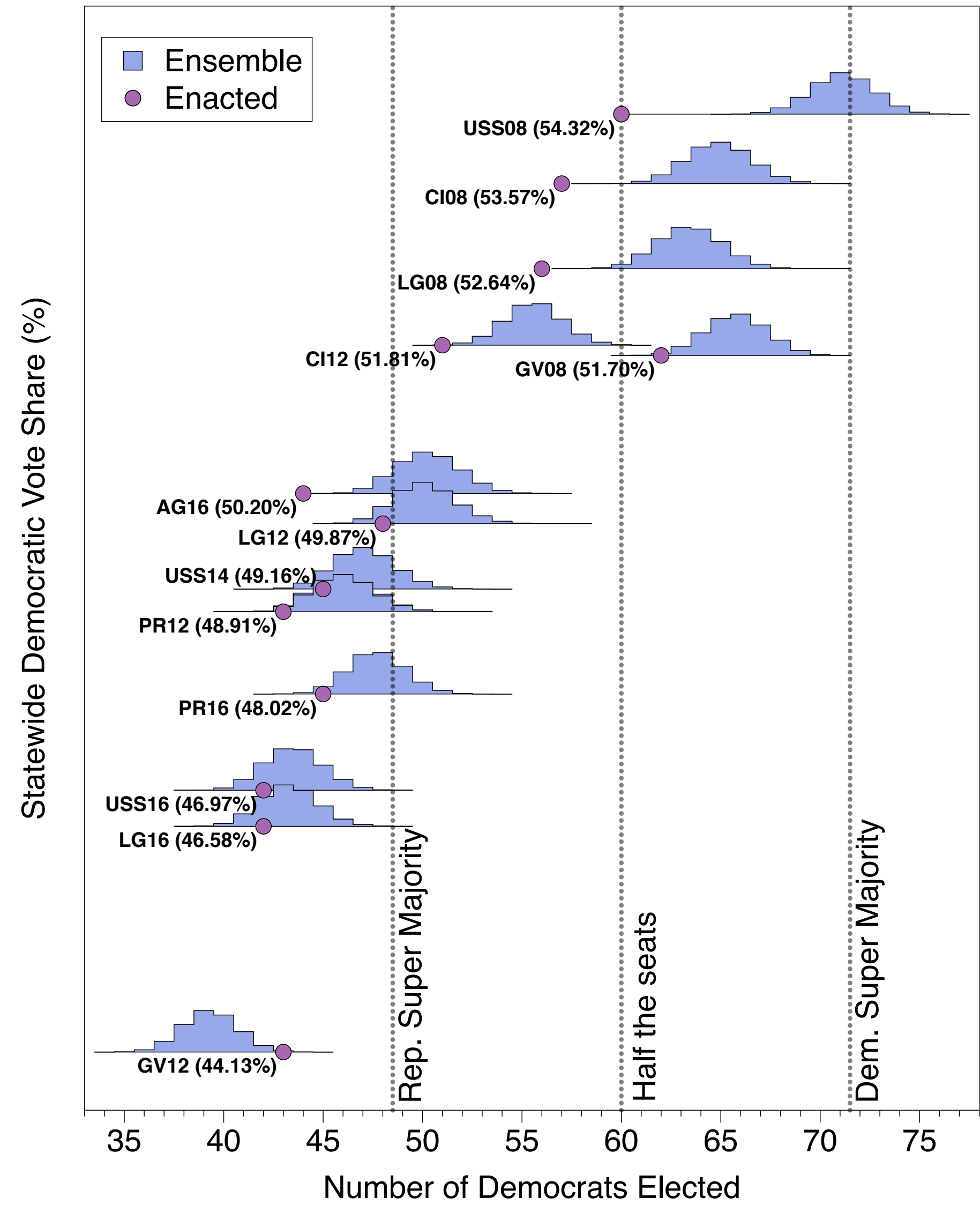


**Counties are kept intact to the extent possible even when split**





# Across many elections



# The Team is Constantly Growing

Christy Vaughn Graves (UG; 2013-2016)

Sachet Bangia (UG; 2016-2017)

Sophie Guo (UG; 2016)

Bridget Duo (UG; 2016)

Hansung Kang (UG; 2016-2017)

Justin Luo (UG; 2016-2017)

Michael Kepler (MS; 2018)

Sam Eure (UG; 2018-2019)

Mike Bell (GS; 2017-2019)

Rahul Ramesh (UG; 2018-Present)

Lisa Lebovich (MS; 2018-Present)

Robert Ravier (GS; 2016-Present)

Andrew Chin (2018-Present)

**Zach Hunter** (2019)

**Daniel Carter** (2019)

Matthias Sachs (2019-Present)

**Eric Autrey** (2019-Present)

**Jonathan Mattingly** (2013-Present)

**Gregory Herschlag** (2016-Present)

## Collaborators

Guy-Uriel Charles

Janice McCarthy

Lydia Kwee

Andrew Chin

Colin Rundel

Jason Parsely

Adam Graham-Squire

Stephen Schechter

Wes Pegden

## Bass Connections Class (year long UG class)

Claire Weibe

Ella Van Engen

Jay Patel

Gillian Samios

Mitra Kiciman

Isaac Nicchitta

Nima Mohammadi

Yashas Manjunatha

Rayan Tofique

Samuel Eure

Tiffany Mei

Luke Farrell

Samuel Eure

Tiffany Mei

Luke Farrell

Jake Shulman

Vinay Kshirsagar

Rahul Ramesh

Haley Sink

Jacob Rubin

Chris Welland

Lynn Fan

Jake Shulman

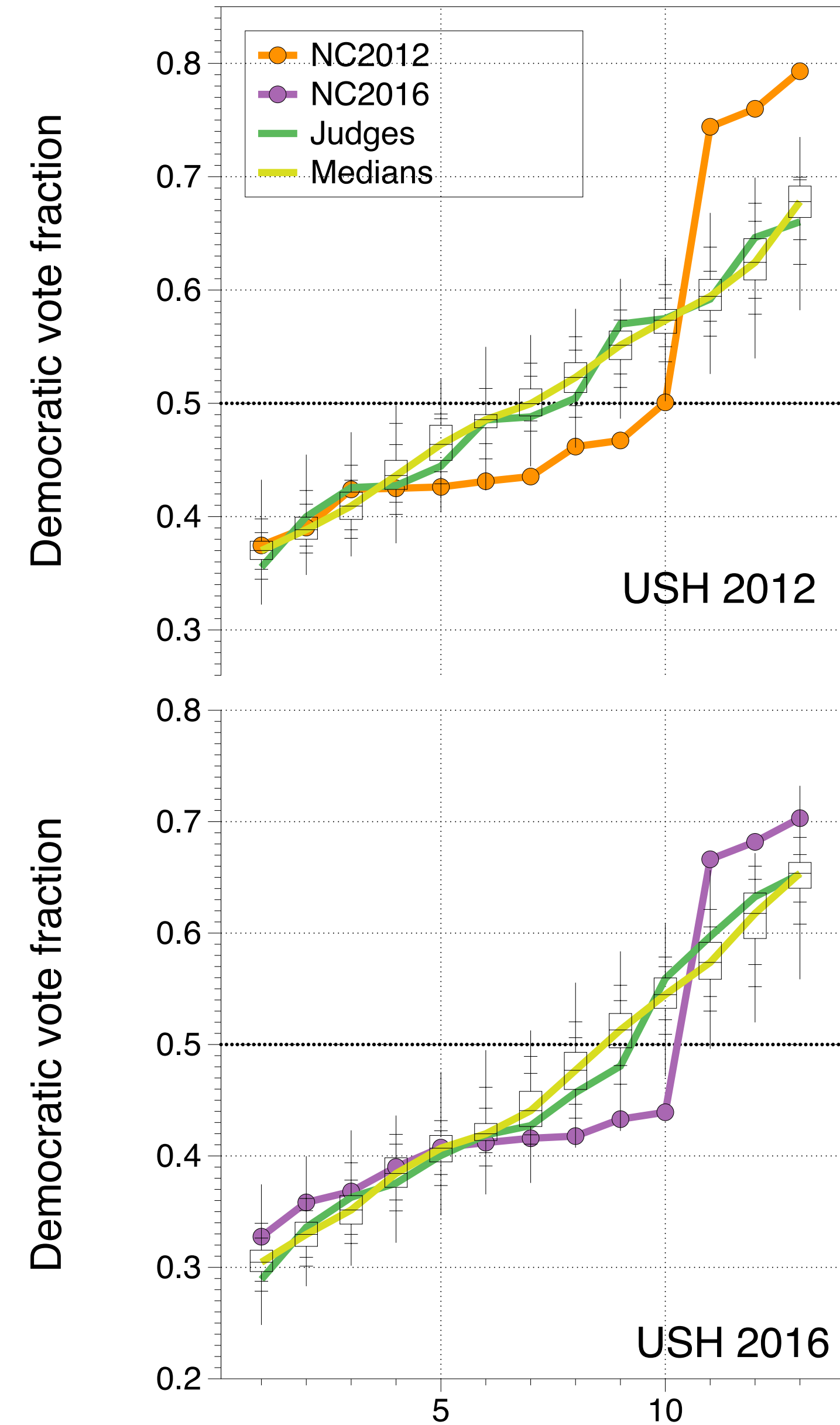
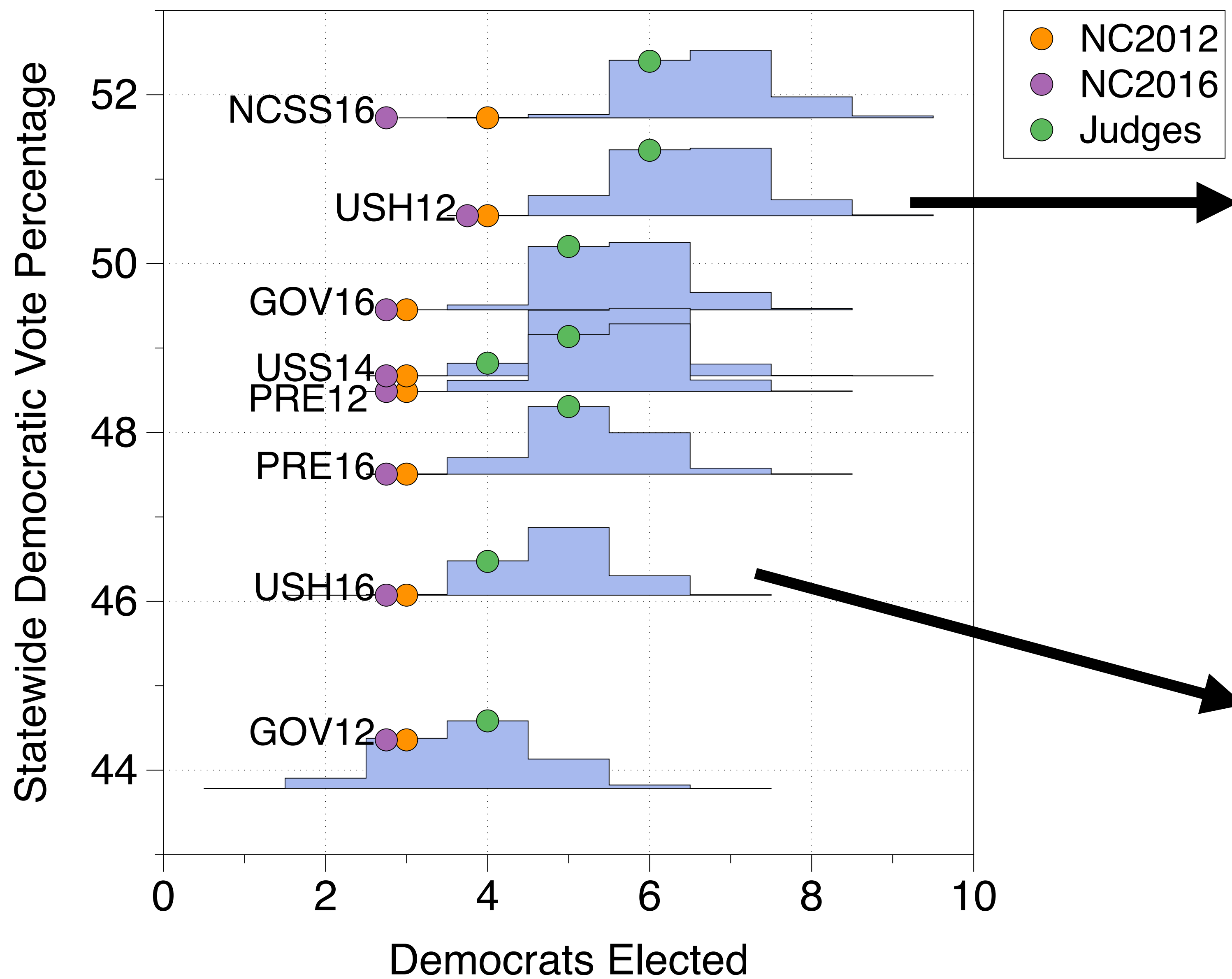
Vinay Kshirsagar



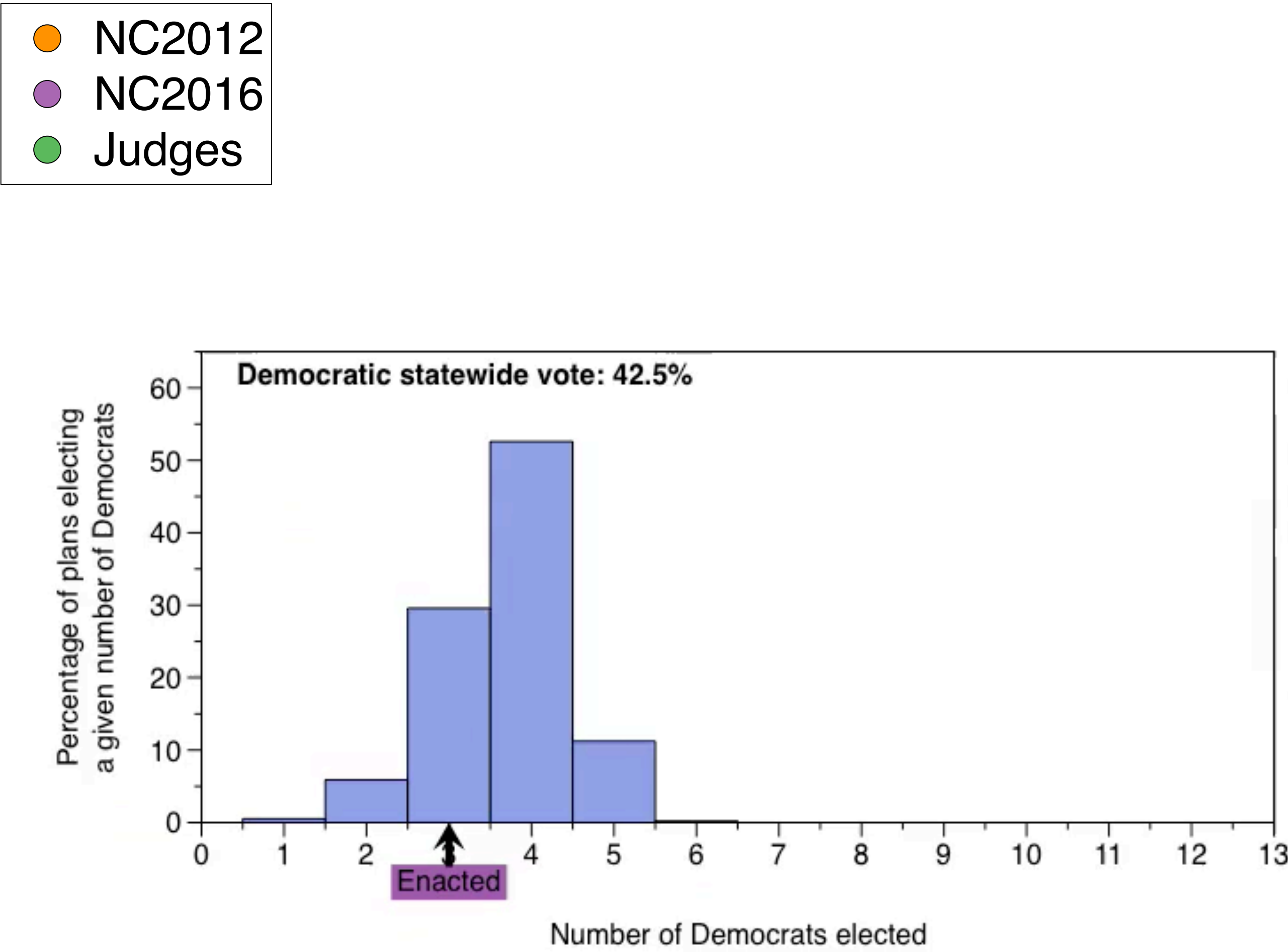
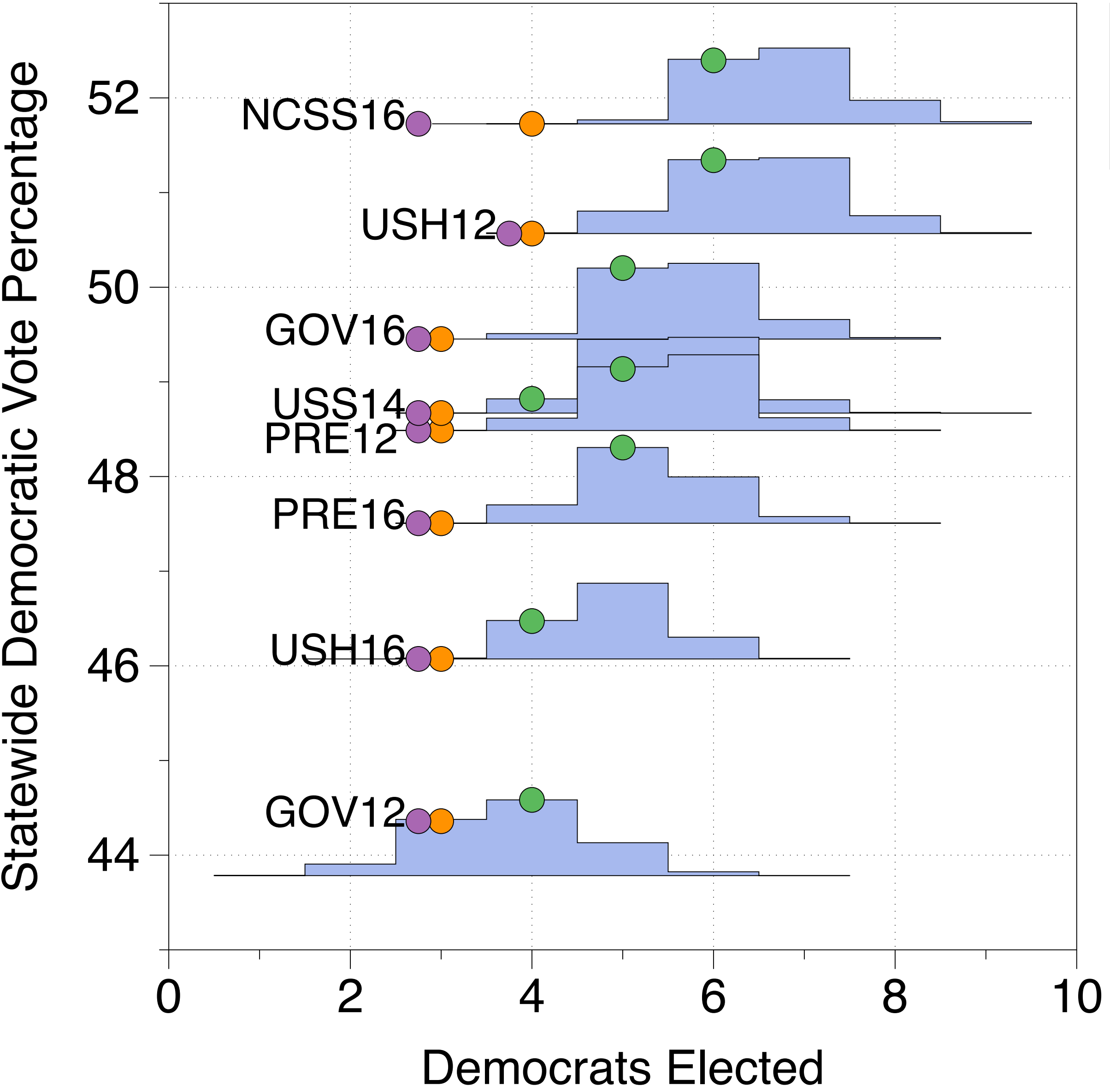
**Blog: <https://sites.duke.edu/quantifyinggerrymandering/>**



# Stagnating election results



# Across many elections





# Partition v. Spanning Forest

**Initial Partition:**  $\xi = (\xi_1, \dots, \xi_i, \dots, \xi_j, \dots, \xi_n)$

**Target Partition:**  $\xi' = (\xi_1, \dots, \xi'_i, \dots, \xi'_j, \dots, \xi_n)$

**ReCom:**

$$A(\xi, \xi') = e^{-\beta[J(\xi') - J(\xi)]} \frac{Q(\xi', \xi)}{Q(\xi, \xi')} \frac{\tau(\xi)^{\gamma-1}}{\tau(\xi')^{\gamma-1}}$$

$Q(\xi, \xi')$  &  $Q(\xi', \xi)$  **expensive**

**Merge-Split:**

$$A(T, T') = e^{-\beta[J(\xi') - J(\xi)]} \frac{Q(T', T)}{Q(T, T')} \frac{\tau(\xi)^\gamma}{\tau(\xi')^\gamma}$$

$Q(\xi, \xi')$  &  $Q(\xi', \xi)$  **cheaper**

$Q(\xi, \xi') : \{\xi_i, \xi_j\} \xrightarrow[\text{deterministic}]{\text{many-to-one}} \xi_{ij} \xrightarrow[\text{random}]{\text{one-to-many}} T'_{ij} \xrightarrow[\text{random}]{\text{one-to-a-few}} \{T'_i, T'_j\} \xrightarrow[\text{random}]{\text{many-to-one}} \{\xi'_i, \xi'_j\}$

$Q(T, T') : \{T_i, T_j\} \xrightarrow[\text{deterministic}]{\text{many-to-one}} \xi_{ij} \xrightarrow[\text{random}]{\text{one-to-many}} T'_{ij} \xrightarrow[\text{random}]{\text{one-to-a-few}} \{T'_i, T'_j\}$