

# Introduction to Game Theory

## Lecture 2: Strategic Game and Nash Equilibrium (cont.)

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## Best response functions: example

- In simple games we can examine each action profile in turn to see if it is a NE. In more complicated games it is better to use “best response functions”.
- Example:

		Player 2		
		L	M	R
Player 1	T	1, 1	1, 0	0, 1
	B	1, 0	0, 1	1, 0

- What are player 1's best response(s) when player 2 chooses L, M, or R?

## Best response functions: definition

- Notation:

$$B_i(a_{-i}) = \{a_i \text{ in } A_i : U_i(a_i, a_{-i}) \geq U_i(a'_i, a_{-i}) \text{ for all } a'_i \text{ in } A_i\}.$$

- I.e., any action in  $B_i(a_{-i})$  is **at least as good** for player  $i$  as every other action of player  $i$  when the other players' actions are given by  $a_{-i}$ .
- Example:

		Player 2		
		L	M	R
Player 1	T	1, 1	1, 0	0, 1
	B	1, 0	0, 1	1, 0

- $B_1(L) = \{T, B\}$ ,  $B_1(M) = \{T\}$ ,  $B_1(R) = \{B\}$

## Using best response functions to define Nash equilibrium

- Definition: the action/strategy profile  $a^*$  is a NE of a strategic game iff every player's action is a best response to the other players' actions:  $a_i^*$  is in  $B_i(a_{-i}^*)$  for every player  $i$ .
- If each player has a single best response to each list  $a_{-i}$  of the other players' actions, then  $a_i = b_i(a_{-i}^*)$  for every  $i$ .

## Using best response functions to find Nash equilibrium

- Method:
  - find the best response function of each player
  - find the action profile in which each player's action is a best response to the other player's action
- Example:

		Player 2		
		L	M	R
Player 1	T	1, 2	2, 1	1, 0
	M	2, 1	0, 1	0, 0
	B	0, 1	0, 0	1, 2

## One more example

- Example 39.1 Two people are involved in a synergistic relationship. If both devote more effort to the relationship, they are both better off. For any given effort of individual  $j$ , the return to individual  $i$ 's effort first increases, then decreases. Specifically, an effort level is a nonnegative number, and each individual  $i$ 's preferences are represented by the payoff function  $u_i = e_i(c + e_j - e_i)$ , where  $e_i$  is  $i$ 's effort level,  $e_j$  is the other individual's effort level, and  $c > 0$  is a constant.

## Solving the example

- $u_i = -e_i^2 + (c + e_j)e_i$ , a quadratic function (section 17.3);  
inverted U-shape
- $u_i = 0$  if  $e_i = 0$  or if  $e_i = c + e_j$ , so anything in between will  
give  $i$  a positive payoff
- Symmetry of quadratic functions means that  $b_i(e_j) = \frac{1}{2}(c + e_j)$
- Similarly,  $b_j(e_i) = \frac{1}{2}(c + e_i)$

## In you know a little calculus

- $U_i = e_i(c + e_j - e_i)$
- First order condition:  $\frac{\partial u_i}{\partial e_i} = c + e_j - 2e_i = 0 \Rightarrow$

$$e_i = \frac{c + e_j}{2} \quad (1)$$

- Similarly,

$$e_j = \frac{c + e_i}{2} \quad (2)$$

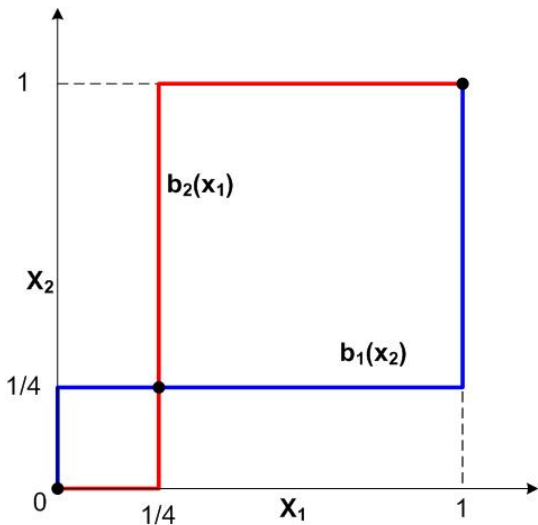
- Plugging (2) into (1), we have  $e_i^* = e_j^* = c$



## Another example about cooperation

- Exercise 42.2 (b) Two people are engaged in a joint project. If each person  $i$  puts in effort  $x_i$ , a nonnegative number equal to at most 1, which costs her  $x_i$ , each person will get a utility  $4x_1x_2$ . Find the NE of the game. Is there a pair of effort levels that yields higher payoffs for both players than do the NE effort levels?

## Best response functions in graph



## Strict Domination (强占优)

- Player  $i$ 's action  $a'_i$  **strictly dominates** action  $a''_i$  if

$$u_i(a'_i, a_{-i}) > u_i(a''_i, a_{-i})$$

for every list  $a_{-i}$  of the other players' actions. In this case the action  $a''_i$  is **strictly dominated**.

- In *Prisoner's Dilemma*, "confess" strictly dominates "silent".

		Suspect 2	
		silent	confess
Suspect 1	silent	0, 0	-2, 1
	confess	1, -2	-1, -1

- If player  $i$ 's action  $a'_i$  strictly dominates every other action of hers, then  $a'_i$  is  $i$ 's **strictly dominant action**.

## Elimination of strictly dominated action

- Not every game has a strictly dominated action. But if there is, it is not used in any Nash equilibrium and so can be eliminated.
- Any strictly dominated action in the following game? Any strictly dominant action?

		Player 2		
		L	C	R
Player 1	U	7, 3	0, 4	4, 4
	M	4, 6	1, 5	5, 3
	D	3, 8	0, 2	3, 0

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⇒ D is strictly dominated by M.

## Iterated elimination of strictly dominated action

- Sometimes we can repeat the procedure: eliminate all strictly dominated actions, and then continue to eliminate strategies that are now dominated in the simpler game.
- Are there more than one actions that can be eliminated from the following game?

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⇒ First B and then C can be eliminated.

## Weak Domination (弱占优)

- Player  $i$ 's action  $a'_i$  **weakly dominates** action  $a''_i$  if

$$u_i(a'_i, a_{-i}) \geq u_i(a''_i, a_{-i})$$

for every list  $a_{-i}$  of the other players' actions, and

$$u_i(a'_i, a_{-i}) > u_i(a''_i, a_{-i})$$

for some list  $a_{-i}$  of the other players' actions.

- Action  $a''_i$  is then **weakly dominated**.



## Weak Domination

- Any weakly dominated action in the following game?

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⇒ R weakly dominated by C; D weakly dominated by M.

- If player  $i$ 's action  $a'_i$  weakly dominates every other action of hers, then  $a'_i$  is  $i$ 's **weakly dominant action**.

## Example: Voting

There are two candidates A and B for an office, and  $N$  voters,  $N \geq 3$  and odd. A majority of voters prefer A to win.

- Is there a strictly dominated action? A weakly dominated action?
- What are the Nash equilibria of the game? Hint: Let  $N_A$  denote the number of voters that vote for A, and  $N_B$  the number of voters that vote for B,  $N_A + N_B = N$ , then
  - ▷ What if  $N_A = N_B + 1$  or  $N_B = N_A + 1$ , and some citizens who votes for the winner actually prefers the loser?
  - ▷ What if  $N_A = N_B + 1$  or  $N_B = N_A + 1$ , and nobody who votes for the winner actually prefers the loser?
  - ▷ Can it happen that  $N_A = N_B + 2$  or  $N_B = N_A + 2$ ?
  - ▷ What if  $N_A \geq N_B + 3$  or  $N_B \geq N_A + 3$ ?

## Solving the voting problem

- What if  $N_A = N_B + 1$  or  $N_B = N_A + 1$ , and some citizens who votes for the winner actually prefers the loser?  $\Rightarrow$  A citizen like that can unilaterally deviate and make her favorite candidate win. Not a NE.
- What if  $N_A = N_B + 1$  or  $N_B = N_A + 1$ , and nobody who votes for the winner actually prefers the loser?  $\Rightarrow$  The former is a NE, but the latter cannot occur (the supporters of B would be more than half).
- Can it happen that  $N_A = N_B + 2$  or  $N_B = N_A + 2$ ?  $\Rightarrow$  No, because  $N$  is odd.
- What if  $N_A \geq N_B + 3$  or  $N_B \geq N_A + 3$ ?  $\Rightarrow$  Yes, NE.

## Strategic voting

- There are three candidates, A, B, and C, and no voter is indifferent between any two of them.
- Voting for one's least preferred candidate is a weakly dominated action. What about voting for one's second preference? Not dominated.

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- Voting for one's least preferred candidate is a weakly dominated action. What about voting for one's second preference? Not dominated.
- Suppose you prefer A to B to C, and the other citizens' votes are tied between B and C, with A being a distant third. Then voting for B, your second preference, is your best choice! ⇒ **strategic voting**
- In two-candidate elections you are weakly better off by voting for your favorite candidate, but in three-candidate elections that is not necessarily the case. E.g, Nader supporters in 2000 US election.

## Hotelling/Downsian model

- A workhorse model of electoral competition. First proposed by Hotelling (1929) and popularized by Downs (1957).
- Setup:
  - ▷ Parties/candidates compete by choosing a policy on the line segment  $[0, 1]$ . The party with most votes wins; if there is a draw, each party has a 50% chance of winning.
  - ▷ Parties only care about winning, and will commit to the platforms they have chosen.
  - ▷ Each voter has a favorite policy on  $[0, 1]$ ; her utility decreases as the winner's position is further away from her favorite policy  $\Rightarrow$  **single-peaked** preference (单峰偏好)
  - ▷ Each voter will vote **sincerely**, choosing the party whose position is closest to her favorite policy.
  - ▷ There is a median voter position,  $m$ .



## Two parties

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- Suppose there are 2 parties,  $L$  and  $R$ . What is the Nash equilibrium for the parties' positions?
- The unique equilibrium is both parties choose position  $m$ .
  - ▷  $(m, m)$  is clearly a NE
  - ▷ any other action profile is not a NE
- This is the so-called "median voter theorem" (中位选举人定理).

## Three parties

- Suppose there is a continuum of voters, with favorite policies uniformly distributed on  $[0, 1]$ , and the number of parties is 3 (L, C, R). Do we still have the equilibrium that all parties choose  $m$ ?

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- Would the three parties positioning at 0.45, 0.55, 0.6 be a NE?
  - ⇒ Yes. L wins already; C and R cannot win by moving anywhere.

## Condorcet winner

- A **Condorcet winner** in an election is a position,  $x^*$ , such that for every other position  $y$  that is different from  $x^*$ , a majority of voters prefer  $x^*$  to  $y$
- The median voter position is a Condorcet winner.
- Not all election games have a Condorcet winner.
  - ▶ Condorcet paradox: A prefers X to Y to Z; B prefers Y to Z to X; C prefers Z to X to Y.
- Even if there is a Condorcet winner, it only has guaranteed victory in pairwise comparisons, not necessarily when there are three or more policy alternatives.

## The strategic model of the war of attrition

- Examples: animals fighting over prey; interest groups lobbying against each other; countries fighting each other to see who will give up first...
- Model setup
  - ▷ Two players,  $i$  and  $j$ , vying for an object, which is respectively worth  $v_i$  and  $v_j$  to the two players; a 50% chance of obtaining the object is worth  $\frac{v_i}{2}$  and  $\frac{v_j}{2}$ .
  - ▷ Time starts at 0 and runs indefinitely; each unit of time that passes before one of the parties concedes costs each player one unit of utility
  - ▷ So, a player  $i$ 's utility is

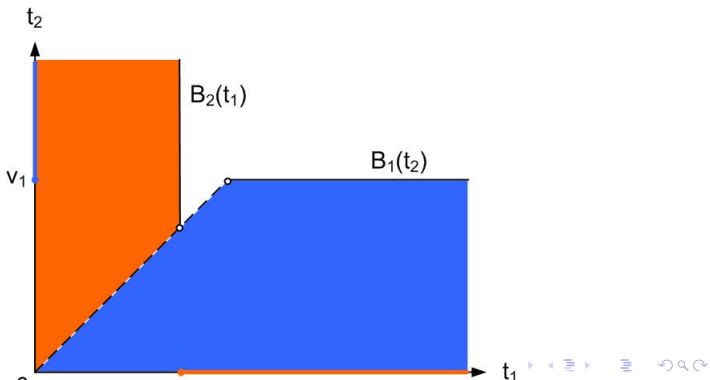
$$u_i(t_i, t_j) = \begin{cases} -t_i, & \text{if } t_i < t_j; \\ \frac{1}{2}v_i - t_j, & \text{if } t_i = t_j; \\ v_i - t_j, & \text{if } t_i > t_j. \end{cases}$$



## Best response function

- Player 2's best response function is (orange)

$$B_2(t_1) = \begin{cases} \{t_2 : t_2 > t_1\}, & \text{if } t_1 < v_2; \\ \{t_2 : t_2 = 0 \text{ or } t_2 > t_1\}, & \text{if } t_1 = v_2; \\ \{0\}, & \text{if } t_1 > v_2. \end{cases}$$



## NE in war of attrition

- $(t_1, t_2)$  is a NE iff  $t_1 = 0$  and  $t_2 \geq v_1$ , or  $t_2 = 0$  and  $t_1 \geq v_2$
- In equilibrium, either player may concede first, including the one who values the object more
- The equilibria is asymmetric, even when  $v_1 = v_2$  (i.e., when the game is symmetric)
  - ▷ A game is symmetric if  $u_1(a_1, a_2) = u_2(a_2, a_1)$  for every action pair  $(a_1, a_2)$  (if you and your opponent exchange actions, you also exchange your payoffs).

## A direct argument

- If  $t_i = t_j$ , then either player can increase her payoff by conceding slightly later and obtaining the object for sure;  $v_i - t_i - \epsilon > \frac{1}{2}v_i - t_i$  for a sufficiently small  $\epsilon$ .
- If  $0 < t_i < t_j$ , player  $i$  should rather choose  $t_i = 0$  to reduce the loss.
- If  $0 = t_i < t_j < v_i$ , player  $i$  can increase her payoff by conceding slightly after  $t_j$ , but before  $t_i = v_i$ .
- The remaining case is  $t_i = 0$  and  $t_j \geq v_i$ , which we can easily verify as a NE.

## Oligopolistic competition (寡头竞争): The Cournot model

- Two firms produce the same product. The unit cost of production is  $c$ . Let  $q_i$  be firm  $i$ 's output,  $Q = \sum_{i=1}^2 q_i$ , then the market price  $P$  is  $P(Q) = \alpha - Q$ , where  $\alpha$  is a constant.
- Firms choose their output simultaneously. What is the NE?
- Each firm wants to maximize profit. Firm 1's profit is

$$\begin{aligned}\pi_1 &= P(Q)q_1 - cq_1 \\ &= (\alpha - q_1 - q_2)q_1 - cq_1.\end{aligned}$$

- Differentiate  $\pi_1$  with respect to  $q_1$ , we know by the first order condition that firm 1's optimal output (best response) is

$$q_1 = b_1(q_2) \frac{\alpha - q_2 - c}{2} \quad (3)$$

## The Cournot model (cont.)

- Similarly (since the game is symmetric), firm 2's optimal output is

$$q_2 = b_2(q_1) = \frac{\alpha - q_1 - c}{2} \quad (4)$$

- Solving equations (3) and (4) together, we have

$$q_1^* = q_2^* = \frac{1}{3}(\alpha - c).$$

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- If the two firms can collude, they would maximize

$PQ - cQ = (\alpha - Q)Q - cQ$ . And the output

$Q = \frac{1}{2}(\alpha - c) < \frac{2}{3}(\alpha - c)$ , the market price would be

$\alpha - Q = \alpha - \frac{1}{2}(\alpha - c) > \alpha - \frac{2}{3}(\alpha - c)$ .

- Competition (instead of collusion) increases total output, and reduces market price.