

# Class starts after this song

*Yiruma – River Flows in You (2001)*  
*requested by Yiyang Shao (Backend head TA)*

I like playing basketball, volleyball and archery. I have a cat named Lizi(Chestnut).



Submitted at: January 23, 1:08 PM

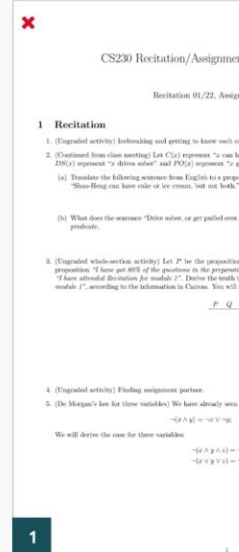
Select questions and pages to indicate where your responses are located. Use **esc** to deselect all.

Use the arrow keys on the keyboard.

## Question Outline

Select a question or a page.

Title	Points
1 1	0.0 pts
1.1 a	0.0 pts
1.2 b	0.0 pts
1.3 c	0.0 pts
2 2	0.0 pts
2.1 a	0.0 pts
2.2 b	0.0 pts
2.3 c	0.0 pts
3 3	0.0 pts
3.1 a	0.0 pts
3.2 b	0.0 pts
4 Recitation part (skip if attended recitation)	0.0 pts



## Unmatched Pages &amp; Questions

**i** You haven't matched all pages and questions.

**Page 1** doesn't have associated questions.

**Question 4** doesn't have associated pages.

You can still submit your assignment without this page associated, however we recommend matching all pages so that graders can easily find your work.

[Submit Assignment](#)
[Continue Matching](#)

- Double check you label the correct pages to each question
- If you don't submit recitation work, Gradescope will warn you that you haven't labeled pages for it (which is normal)

## CM1

● Ungraded

Student

Test Student

[View or edit group](#)

Total Points

- / 0 pts

Question 1

1

0 pts

1.1 a

0 pts

1.2 b

0 pts

1.3 c

0 pts

Question 2

2

0 pts

## Group Members

**i** Add or remove group members for this submission.

Your instructor has allowed you to submit as a group of up to **2 people**. You can change the group below. Students added or removed will be notified via email.

Student

Remove

Test Student

✕

Add Student

Close

Add

- Then add your teammate here; DO NOT SUBMIT SEPARATELY

“The limit of  $f(x)$  exists <sup>everywhere</sup> as  $x$  approaches  $a$ ”

$\forall a \in \mathbb{R}$

$$\exists L \in \mathbb{R} \quad \forall \varepsilon > 0 \quad \exists \delta > 0 \quad \left[ (0 < |x - a| < \delta) \rightarrow (|f(x) - L| < \varepsilon) \right]$$

- We made a mistake in CM1 recitation material here; we **forgot to quantify  $x$**
- The predicate with this part missing was not “false”; it is still true when the limit exists, but it’s this version with  $x$  quantified that semantically translates the limit definition
- In hindsight, this problem still relies on knowing what a limit is; don’t worry about that part and **we were/are just using it for practice on predicate logic**

# CS230 Spring 2024

## Module 02: Proof Methods

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# Why proof?

- You need to make arguments on how things work as a future computer scientist
  - Can you always base your arguments on evidence?
    - “The last 1,000 times I used it, this sorting algorithm actually gives me a sorted list”
  - If not, then you sometimes need to base them on logic
    - Proofs are nothing beyond that (but more formal, and more stylistic)
-

# What can be proved?

- Nothing can be proved *without assumptions/axioms/facts given*
    - Very very strictly speaking, a mathematical proof is incomplete without specifying what axioms are used
    - Practical compromise: state all nontrivial assumptions/axioms
  - Only propositions can be proved/disproved
    - This is how we “defined” propositions!
    - Not all propositions can be proved/disproved *within propositional logic*
-

# How are $\rightarrow$ and $\Rightarrow$ different?

- $\rightarrow$  (conditional operator):  $p \rightarrow q$  can be either true or false
- $\Rightarrow$  (implication):  $p \Rightarrow q$  means “ $p \rightarrow q$  is a tautology”



# Modus Ponens, revisited

$$\frac{p \rightarrow q \quad p}{\therefore q}$$

$$((p \rightarrow q) \wedge p) \Rightarrow q$$

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# Similarly:

- $\leftrightarrow$  (biconditional operator):  $p \leftrightarrow q$  means  $(p \rightarrow q) \wedge (q \rightarrow p)$
  - $\Leftrightarrow, =, \equiv$  (equivalence):  $p \Leftrightarrow q$  means  $(p \Rightarrow q) \wedge (q \Rightarrow p)$
  - When we prove  $p \leftrightarrow q$  is true, we proved  $p \Leftrightarrow q$   
(read:  $p$  is equivalent to  $q$ )
  - “We prove  $p$  if and only if  $q$  (is true)”
-

# PD: Proof Example 1

- Read this proof individually
- Then turn to your neighbor, discuss *where exactly is the proof wrong*



Proof:

Let  $a = 2n + 1$  and  $b = 2q + 1$  be two odd integers.

$n$  and  $q$  need to be integers

Then  $(2n + 1)(2q + 1) = 4nq + 2n + 1 = 2(2nq + n) + 1$ .

missing a  $2q$  (algebraic error)

Since  $(2nq + n)$  is an integer,  $2(2nq + n) + 1$  is an odd integer.

questionable choices of variables (not wrong)

## PD: Proof Example 2

- What about this proof?



states more than what is needed (not wrong, but bad practice)

If  $n$  is an integer,  $n^3 - n = (n - 1)n(n + 1)$  is the product of three consecutive integers  $(n - 1)$ ,  $n$ , and  $(n + 1)$ .

Among these three integers, at least one is a multiple of 2 and exactly one is a multiple of 3.

Since 2 and 3 are both primes,  $n^3 - n = (n - 1)n(n + 1)$  is a multiple of  $2 \times 3 = 6$ .

# Proof by construction

## *[Existential Generalization]*

- To prove a theorem of the form

$$\exists x[P(x)]$$

we merely need to find (“construct”) ONE  $c$  and show  $P(c)$ .

$$P(c) \Rightarrow \exists x[P(x)]$$

---

# Proof by construction

## *[Existential Generalization]*

- Theorem: There exist integers  $a, b, c$  such that  $a^2 + b^2 = c^2$ .
- Proof: let  $a = 3, b = 4, c = 5$ .

# Proof by (complicated) construction

- Theorem: Given two strings  $x_1x_2 \cdots x_n$  and  $y_1y_2 \cdots y_m$ , there is an algorithm that runs in  $O(mn)$ -time that finds the length of their longest common substring, i.e., the largest  $k$  for which there exist indices  $i, j$  with  $x_ix_{i+1} \cdots x_{i+k-1} = y_jy_{j+1} \cdots y_{j+k-1}$ .
  - Proof:
    - Describe the algorithm
    - Prove the algorithm runs in  $O(mn)$ -time
    - Prove the algorithm finds the largest such  $k$  for all possible input strings
-

# Proof by construction

## [Existential Generalization]

- To prove a theorem of the form

$$\exists x[P(x)]$$

we merely need to find (“construct”) ONE  $c$  and show  $P(c)$ .

$$P(c) \Rightarrow \exists x[P(x)]$$

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# Disproof by counterexample

- A disproof by counterexample is itself *a proof by construction of the negation of the initial proposition*.
  - Disproving  $\forall x[P(x)]$  is just proving  $\neg(\forall x[P(x)])$ , which is equivalent to  $\exists x[\neg P(x)]$ .
  - The counterexample is some  $c$  that makes  $\neg P(c)$  true.
-

# Class starts after this song

***Paramore – Proof (2013)***

***requested by Luke Lorentzatos (TA-of-CM2)***

I'm half Greek and can speak a little bit.  
I am a huge Houston Astros fan.



# Visiting Class Today: TAs-of-CM2



***Anirudh Jain***

I'm always down for poker and table tennis. And I respond to at least 8 pronunciations of my name.

Hi, I like taekwondo and reading. I also like skiing and baking but those are harder to do at Duke.

***Jessica Chen***



# Reminders on Gradescope assignment

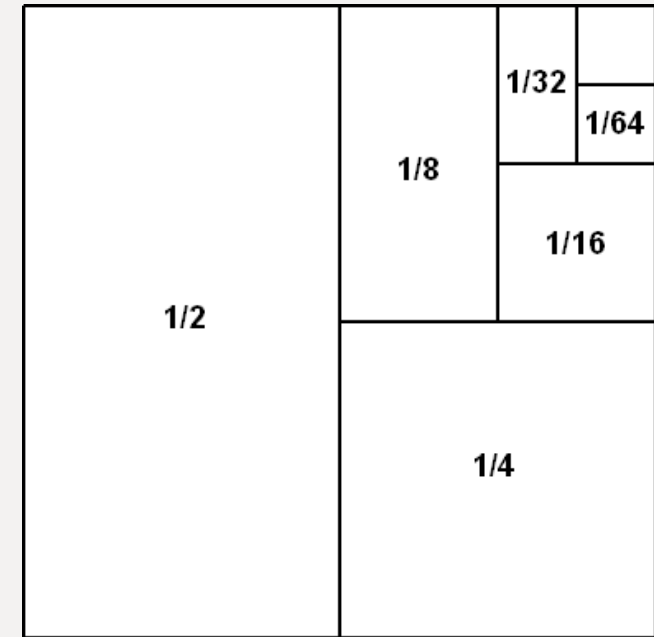
- Type your work; no handwriting accepted
  - The best way to use the provided LaTeX source code is to **directly modify it** (as it's meant to be the skeleton/template)
    - Put `macros.tex` in your project folder as a separate file
    - Use `\mathbb{R}` to typeset blackboard bold  $\mathbb{R}$
  - Label pages after uploading PDF
  - Submit just once per group, then add your teammate
  - No names necessary in PDF
-

# Prove a proposition/theorem/statement

- Direct proof
  - Proof by contrapositive
  - Proof by contradiction
  - Proof by cases
  - Proof by construction
  - Proof by induction [*big topic itself* - CM5]
  - What else?
-

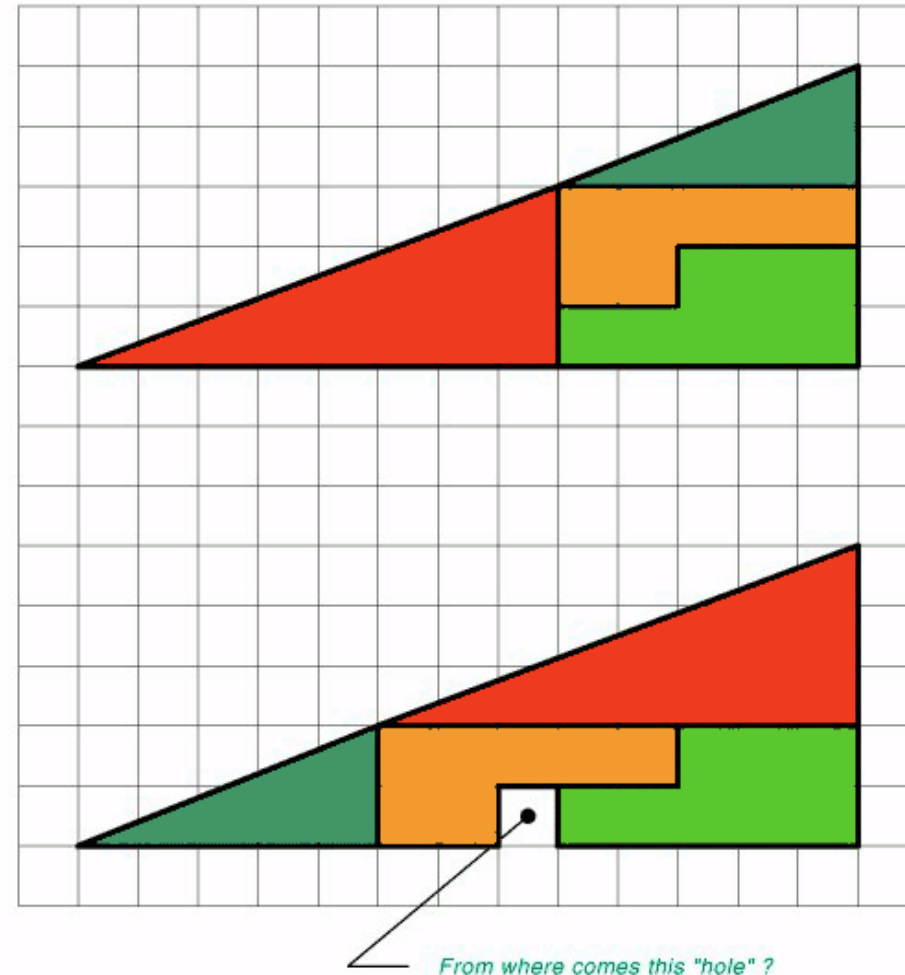
# Proof by picture...?

- $\sum_{i=1}^{\infty} \frac{1}{2^i} = 1$
- A picture alone is not a proof



# Proof by picture

HOW CAN THIS BE TRUE ?



*Below the four  
parts are  
moved around*

*The partitions  
are exactly the  
same, as those  
used above*

*From where comes this "hole" ?*

# PI: Biconditional





# An abstracted version of the question

- Want to prove:  $p \leftrightarrow q \iff (p \rightarrow q) \wedge (q \rightarrow p)$

a)  $p \rightarrow q$

b)  $q \rightarrow p$

c)  $\neg q \rightarrow \neg p$

d)  $\neg p \rightarrow \neg q$

# In case you did not notice

- The theorem is false
  - If exactly two of  $a, b, c$  are even and the other is odd,  $a^2 + b^2 + c^2$  is also odd (check this yourself)
  - That should not matter
    - We were identifying what conditional statements or their contrapositives, that we should prove
    - We were not actually trying to prove them (we would fail)
-

# Proving more than what is needed

- is technically correct
  - is usually unnecessary
  - sometimes makes the proof easier
    - more examples about this in CM5 (Inductions)
-

# Write good proofs, not just correct proofs

- Correctness of the proof is the *first* priority, not the *only* priority
  - Correct proofs are just correct. Good proofs can be *understood*.
    - You want your proofs to be correct and understood
    - What is obvious to yourself, may not be obvious to others
-

# Peer Review (Proof-by-cases)

- Navigate to the Canvas quiz
- Complete two simple proof-by-cases
  - Don't share accounts; complete the proofs on your own
  - Don't discuss with anyone else



# Peer Review (Proof-by-cases)

- Navigate to Canvas quiz again
- You are now assigned two anonymous assignments by your classmates
- Read their work, then give feedback
  - Evaluate the technical correctness, readability, and conciseness of the proofs via the rubric provided



**Proof.** “For every positive integer  $n$ , if  $\sqrt{n}$  is rational, then  $\sqrt{n}$  is an integer.”.

- We use proof by contradiction: assume  $\neg(\forall n \in \mathbb{N} [\sqrt{n} \in \mathbb{Q} \rightarrow \sqrt{n} \in \mathbb{Z}])$
- $\therefore \exists n \in \mathbb{N} [\neg(\sqrt{n} \in \mathbb{Q} \rightarrow \sqrt{n} \in \mathbb{Z})]$
- $\therefore \exists n \in \mathbb{N} [(\sqrt{n} \in \mathbb{Q}) \wedge (\sqrt{n} \notin \mathbb{Z})]$
- $\therefore \exists a \in \mathbb{N} \exists b \in \mathbb{N} [(\sqrt{n} = \frac{a}{b}) \wedge (b \nmid a)]$
- Let  $c \geq 1 = \text{GCD}(a, b)$ .
- $\therefore \exists x \in \mathbb{N} \exists y \in \mathbb{N} [(a = cx) \wedge (b = cy) \wedge (y \nmid x)]$
- $\therefore n = \frac{a^2}{b^2} = \frac{c^2 x^2}{c^2 y^2} = \frac{x^2}{y^2}$ .
- Since  $(y \nmid x) \rightarrow (y^2 \nmid x^2)$ ,  $n = \frac{x^2}{y^2} \notin \mathbb{Z}$ .

**Excessive  
symbolism**

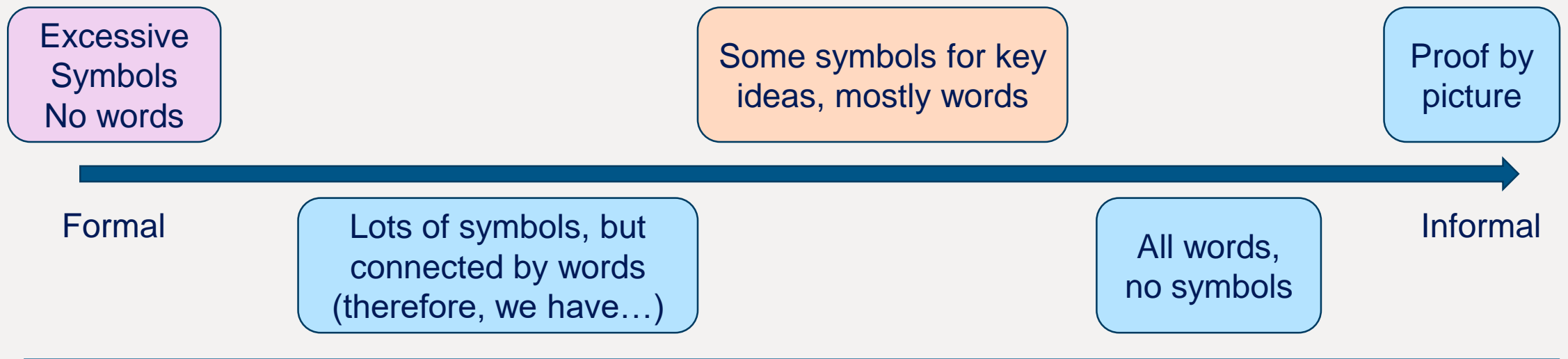
**Proof.** “For every positive integer  $n$ , if  $\sqrt{n}$  is rational, then  $\sqrt{n}$  is an integer.” 32

- We use proof by contradiction: assume the theorem is not true.
- Then there exists a positive integer  $n$  such that  $\sqrt{n}$  is rational but not an integer.
- Then there exist positive integers  $a, b$  such that  $\sqrt{n} = \frac{a}{b}$  and  $b \nmid a$  (i.e.,  $b$  does not divide  $a$ ).
- Let  $c \geq 1$  be the greatest common divisor of  $a$  and  $b$ .
- Then there exist positive integers  $x, y$  such that  $a = cx$ ,  $b = cy$ , and  $y \nmid x$ . (otherwise  $b = cy$  would divide  $a = cx$ ).
- Then  $n = \frac{a^2}{b^2} = \frac{c^2 x^2}{c^2 y^2} = \frac{x^2}{y^2}$ .
- Since  $y \nmid x$  implies  $y^2 \nmid x^2$ ,  $n = \frac{x^2}{y^2}$  is not an integer.
- Since we reached a contradiction, the theorem is proved.

Less  
symbolistic



# Level of formality



# How to come up with the proof steps?

- Reading proofs (especially simple and beautiful ones) make you feel proofs come “naturally” as if all steps just magically fall into place
    - For most of sufficiently complicated theorems, this is usually not the case
    - A lot of trial-and-error (you don’t have a compiler to tell you there’s an error)
    - Messy thoughts, dead-ends, useful but out-of-order ideas...
  - No one expects you to be perfect on the first try
    - Like no one expects you to write 100 lines of code that work immediately
-