

# Euclidean Algorithm

You may or may not have heard about the Euclidean Algorithm. Usually, it is introduced as a systematic method to find  $\text{GCD}(a, b)$  of two positive integers  $a$  and  $b$ . The Euclidean Algorithm relies on a simple result:

*Theorem. For two positive integers  $a > b$ , if  $a \bmod b = c$ , then  $\text{GCD}(a, b) = \text{GCD}(b, c)$ .*

Proof. We only need to prove that the set of (positive) common divisors of  $a$  and  $b$  is identical to the set of (positive) common divisors of  $b$  and  $c$ .

- Suppose  $p$  is a (positive) common divisor of  $a$  and  $b$ . Then  $a = mp$  and  $b = np$  for some positive integers  $m > n$ . Since  $a \bmod b = c$ , we know  $a = kb + c$  for some positive integer  $k$ . Therefore, we have  $c = a - kb = mp - knp = (m - kn)p$ , which implies  $p$  divides  $c$ .
- Suppose  $p$  is a (positive) common divisor of  $b$  and  $c$ . Then  $b = xp$  and  $c = yp$  for some positive integers  $x > y$ . (We know  $x > y$  because  $b > c$ .) Therefore, we have  $a = kb + c = kxp + yp = (kx + y)p$ , which implies  $p$  divides  $a$ .

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For a concrete example, suppose we were to find the greatest common divisor of **230** and **2024**:

$$\begin{aligned} 2024 &= 230 \times 8 + 184 && // \text{ therefore } \text{GCD}(230, 2024) = \text{GCD}(230, 184) \\ 230 &= 184 \times 1 + 46 && // \text{ therefore } \text{GCD}(230, 184) = \text{GCD}(46, 184) \\ 184 &= 46 \times 4 && // \text{ therefore } \text{GCD}(46, 184) = 46 \end{aligned}$$

Therefore, we have  $\text{GCD}(230, 2024) = 46$  (note that the comments on the right make a chain-of-equivalence).

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What is less obvious is the Euclidean Algorithm can also help find the multiplicative inverse in modulo arithmetic (if one exists). More specifically, if  $\text{GCD}(a, b) = 1$ , then the process of Euclidean Algorithm actually reveals the mystery number  $z$  such that  $a \times z \equiv 1 \pmod{b}$ . Look at this concrete example where we find the greatest common divisor of **230** and **7**, although we know in advance that it is **1** (because **7** is a prime and **230** is not a multiple of **7**):

$$\begin{aligned} 230 &= 7 \times 32 + 6 && // \text{ in other words, } 6 = 230 - 7 \times 32 \\ 7 &= 6 \times 1 + 1 && // \text{ in other words, } 1 = 7 - 6 \times 1 \\ 6 &= 1 \times 6 \end{aligned}$$

Now let's look at the notes on the right-hand side and combine the information there:

$$\begin{aligned} 1 &= 7 - 6 \times 1 \\ &= 7 - (230 - 7 \times 32) \times 1 \\ &= 7 \times 33 - 230 \\ &= 7 \times 33 + 230 \times (-1). \end{aligned}$$

This implies  $1 \equiv 230 \times (-1) \pmod{7}$ . If we don't want the mystery number to be negative, we can also conclude that  $1 \equiv 230 \times 6 \pmod{7}$ , since  $6 \equiv (-1) \pmod{7}$ .

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Although in the example above **7** is a prime, the algorithm works for any two coprime integers **a** and **b**. Therefore, it is more powerful (strictly speaking about finding multiplicative inverses) than Fermat's Little Theorem, because the latter only works when **b** is a prime.

Practice the Euclidean Algorithm in the next practice quiz.